

# Gödel's System T

15-814: Types and Programming Languages

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September 11, 2024

This handout contains System T with a call-by-value (or eager) dynamic semantics.

## 1 Syntax

<b>Typ</b>	$\tau ::= \text{nat}$	number
	$\tau_1 \rightarrow \tau_2$	function
<b>Exp</b>	$e ::= x$	variable
	$z$	zero
	$s(e)$	successor
	$\text{rec}\{e_0; x.y.e_1\}(e)$	recursion
	$\lambda(x : \tau)e$	abstraction
	$e_1(e_2)$	application

## 2 Statics

$$\frac{}{\Gamma, x : \tau \vdash x : \tau} T_{\text{var}} \quad \frac{}{\Gamma \vdash z : \text{nat}} T_z \quad \frac{\Gamma \vdash e : \text{nat}}{\Gamma \vdash s(e) : \text{nat}} T_s$$

$$\frac{\Gamma \vdash e : \text{nat} \quad \Gamma \vdash e_0 : \tau \quad \Gamma, x : \text{nat}, y : \tau \vdash e_1 : \tau}{\Gamma \vdash \text{rec}\{e_0; x.y.e_1\}(e) : \tau} T_{\text{rec}}$$

$$\frac{\Gamma, x : \tau_1 \vdash e_2 : \tau_2}{\Gamma \vdash \lambda(x : \tau_1)e_2 : \tau_1 \rightarrow \tau_2} T_{\text{lam}} \quad \frac{\Gamma \vdash e_1 : \tau_1 \rightarrow \tau_2 \quad \Gamma \vdash e_2 : \tau_1}{\Gamma \vdash e_1(e_2) : \tau_2} T_{\text{ap}}$$

## 3 Call-by-Value Dynamics

### 3.1 Values

$$\frac{}{\lambda(x : \tau)e \text{ val}} V_{\text{lam}} \quad \frac{}{z \text{ val}} V_z \quad \frac{e \text{ val}}{s(e) \text{ val}} V_s$$

### 3.2 Transitions

$$\frac{e \mapsto e'}{s(e) \mapsto s(e')} E_s \quad \frac{e \mapsto e'}{\text{rec}\{e_0; x.y.e_1\}(e) \mapsto \text{rec}\{e_0; x.y.e_1\}(e')} E_{\text{rec1}}$$

$$\frac{}{\mathbf{rec}\{e_0; x.y.e_1\}(z) \mapsto e_0} E_{\text{rec}2} \quad \frac{s(e) \text{ val}}{\mathbf{rec}\{e_0; x.y.e_1\}(s(e)) \mapsto [e, \mathbf{rec}\{e_0; x.y.e_1\}(e)/x, y]e_1} E_{\text{rec}3}$$

$$\frac{e_1 \mapsto e'_1}{e_1(e_2) \mapsto e'_1(e_2)} E_{\text{ap}1} \quad \frac{e_1 \text{ val} \quad e_2 \mapsto e'_2}{e_1(e_2) \mapsto e_1(e'_2)} E_{\text{ap}2} \quad \frac{e_2 \text{ val}}{(\lambda(x : \tau)e)(e_2) \mapsto [e_2/x]e} E_{\text{ap}3}$$