

# Designing extensible, domain-specific languages for mathematical diagrams

Katherine Ye

with Keenan Crane Jonathan Aldrich Josh Sunshine

## (A proposal for)



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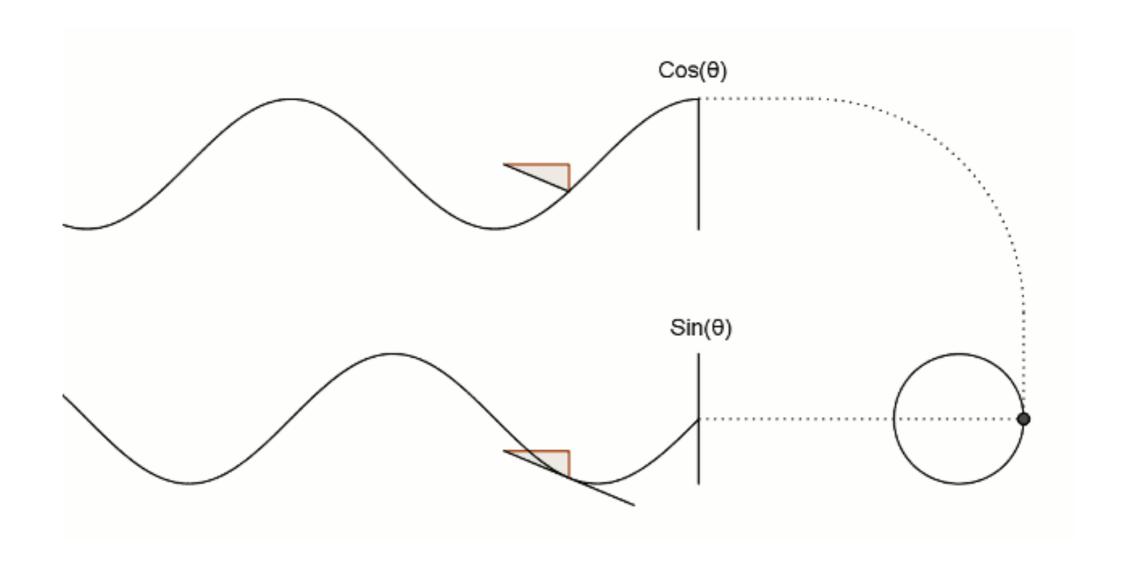
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# Motivation

$$\sin \theta = \frac{opposite}{hypotenuse}$$

$$\cos \theta = \frac{adjacent}{hypotenuse}$$



#### MEAN VALUE THEOREMS

5.7 **Definition** Let f be a real function defined on a metric space X. We say that f has a *local maximum* at a point  $p \in X$  if there exists  $\delta > 0$  such that  $f(q) \le f(p)$  for all  $q \in X$  with  $d(p, q) < \delta$ .

Local minima are defined likewise.

Our next theorem is the basis of many applications of differentiation.

5.8 Theorem Let f be defined on [a, b]; if f has a local maximum at a point  $x \in (a, b)$ , and if f'(x) exists, then f'(x) = 0.

The analogous statement for local minima is of course also true.

**Proof** Choose  $\delta$  in accordance with Definition 5.7, so that

$$a < x - \delta < x < x + \delta < b$$
.

If  $x - \delta < t < x$ , then

$$\frac{f(t) - f(x)}{t - x} \ge 0.$$

Letting  $t \to x$ , we see that  $f'(x) \ge 0$ .

If  $x < t < x + \delta$ , then

$$\frac{f(t) - f(x)}{t - x} \le 0,$$

which shows that  $f'(x) \le 0$ . Hence f'(x) = 0.

5.9 Theorem If f and g are continuous real functions on [a, b] which are differentiable in (a, b), then there is a point  $x \in (a, b)$  at which

$$[f(b) - f(a)]g'(x) = [g(b) - g(a)]f'(x).$$

Note that differentiability is not required at the endpoints.

Proof Put

$$h(t) = [f(b) - f(a)]g(t) - [g(b) - g(a)]f(t) \qquad (a \le t \le b).$$

Then h is continuous on [a, b], h is differentiable in (a, b), and

(12) 
$$h(a) = f(b)g(a) - f(a)g(b) = h(b).$$

To prove the theorem, we have to show that h'(x) = 0 for some  $x \in (a, b)$ . If h is constant, this holds for every  $x \in (a, b)$ . If h(t) > h(a) for some  $t \in (a, b)$ , let x be a point on [a, b] at which h attains its maximum (Theorem 4.16). By (12),  $x \in (a, b)$ , and Theorem 5.8 shows that h'(x) = 0. If h(t) < h(a) for some  $t \in (a, b)$ , the same argument applies if we choose for x a point on [a, b] where h attains its minimum.

This theorem is often called a generalized mean value theorem; the following special case is usually referred to as "the" mean value theorem:

**5.10 Theorem** If f is a real continuous function on [a, b] which is differentiable in (a, b), then there is a point  $x \in (a, b)$  at which

$$f(b) - f(a) = (b - a)f'(x).$$

**Proof** Take g(x) = x in Theorem 5.9.

- **5.11 Theorem** Suppose f is differentiable in (a, b).
  - (a) If  $f'(x) \ge 0$  for all  $x \in (a, b)$ , then f is monotonically increasing.
  - (b) If f'(x) = 0 for all  $x \in (a, b)$ , then f is constant.
  - (c) If  $f'(x) \le 0$  for all  $x \in (a, b)$ , then f is monotonically decreasing.

Proof All conclusions can be read off from the equation

$$f(x_2) - f(x_1) = (x_2 - x_1)f'(x),$$

which is valid, for each pair of numbers  $x_1$ ,  $x_2$  in (a, b), for some x between  $x_1$  and  $x_2$ .

#### THE CONTINUITY OF DERIVATIVES

We have already seen [Example 5.6(b)] that a function f may have a derivative f' which exists at every point, but is discontinuous at some point. However, not every function is a derivative. In particular, derivatives which exist at every point of an interval have one important property in common with functions which are continuous on an interval: Intermediate values are assumed (compare Theorem 4.23). The precise statement follows.

**5.12 Theorem** Suppose f is a real differentiable function on [a, b] and suppose  $f'(a) < \lambda < f'(b)$ . Then there is a point  $x \in (a, b)$  such that  $f'(x) = \lambda$ .

A similar result holds of course if f'(a) > f'(b).

**Proof** Put  $g(t) = f(t) - \lambda t$ . Then g'(a) < 0, so that  $g(t_1) < g(a)$  for some  $t_1 \in (a, b)$ , and g'(b) > 0, so that  $g(t_2) < g(b)$  for some  $t_2 \in (a, b)$ . Hence g attains its minimum on [a, b] (Theorem 4.16) at some point x such that a < x < b. By Theorem 5.8, g'(x) = 0. Hence  $f'(x) = \lambda$ .

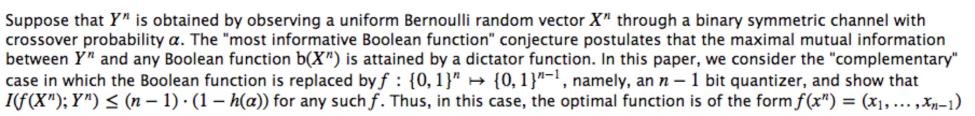
#### [1] arXiv:1701.03119 [pdf, ps, other]

#### How to Quantize n Outputs of a Binary Symmetric Channel to n-1 Bits?

Wasim Huleihel, Or Ordentlich

Comments: 5 pages, submitted to ISIT 2017

Subjects: Information Theory (cs.IT)





#### [2] arXiv:1701.03120 [pdf, ps, other]

#### The fourth moment theorem on the Poisson space

Christian Döbler, Giovanni Peccati

Comments: 31 pages

Subjects: Probability (math.PR)

We prove an exact fourth moment bound for the normal approximation of random variables belonging to the Wiener chaos of a general Poisson random measure. Such a result -- that has been elusive for several years -- shows that the so-called `fourth moment phenomenon', first discovered by Nualart and Peccati (2005) in the context of Gaussian fields, also systematically emerges in a Poisson framework. Our main findings are based on Stein's method, Malliavin calculus and Mecke-type formulae, as well as on a methodological breakthrough, consisting in the use of carr\'e-du-champ operators on the Poisson space for controlling residual terms associated with add-one cost operators. Our approach can be regarded as a successful application of Markov generator techniques to probabilistic approximations in a non-diffusive framework: as such, it represents a significant extension of the seminal contributions by Ledoux (2012) and Azmoodeh, Campese and Poly (2014). To demonstrate the flexibility of our results, we also provide some novel bounds for the Gamma approximation of non-linear functionals of a Poisson measure.



#### [3] arXiv:1701.03124 [pdf, ps, other]

#### Totaro's Question for Adjoint Groups of Types A\_{1} and A\_{2n}

Reed Gordon-Sarney Comments: 7 pages

Subjects: Algebraic Geometry (math.AG)

Let G be a smooth connected linear algebraic group over a field k, and let X be a G-torsor. Totaro asked: if X admits a zerocycle of degree  $d \ge 1$ , then does X have a closed \'etale point of degree dividing d? We give an affirmative answer for absolutely simple classical adjoint groups of types  $A_1$  and  $A_{2n}$  over fields of characteristic  $\neq 2$ .

We provide an integral estimate for a non-divergence (non-variational) form second order elliptic equation  $a_{ii}u_{ii}=u^p$ ,  $u\geq 0$ , n C [0, 1) with bounded discontinuous coefficients a boying small PMO norm. We consider the simplest discontinuity of the

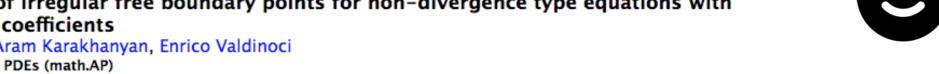


#### [4] arXiv:1701.03131 [pdf, other]

#### Classification of irregular free boundary points for non-divergence type equations with discontinuous coefficients

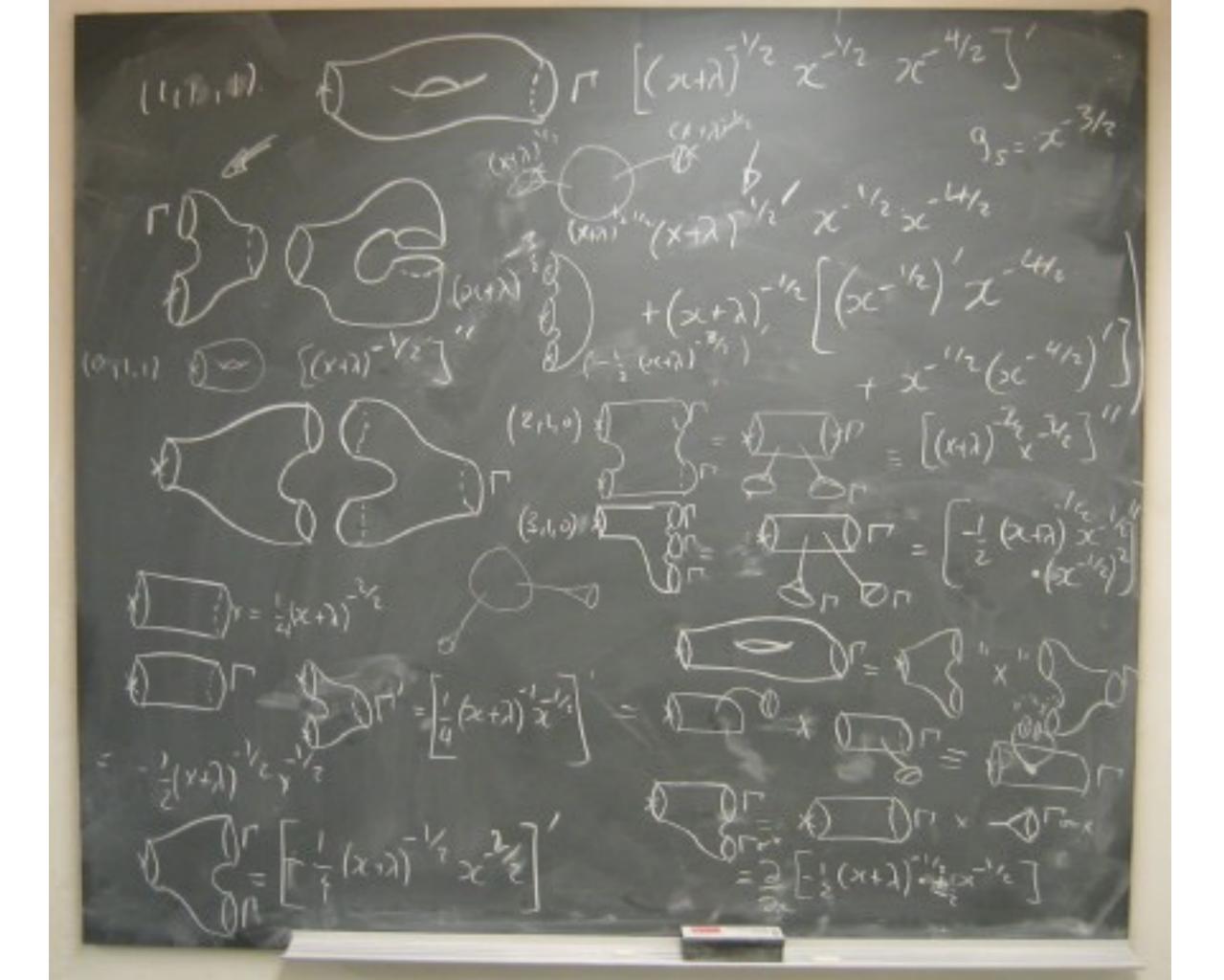
Serena Dipierro, Aram Karakhanyan, Enrico Valdinoci

Subjects: Analysis of PDEs (math.AP)

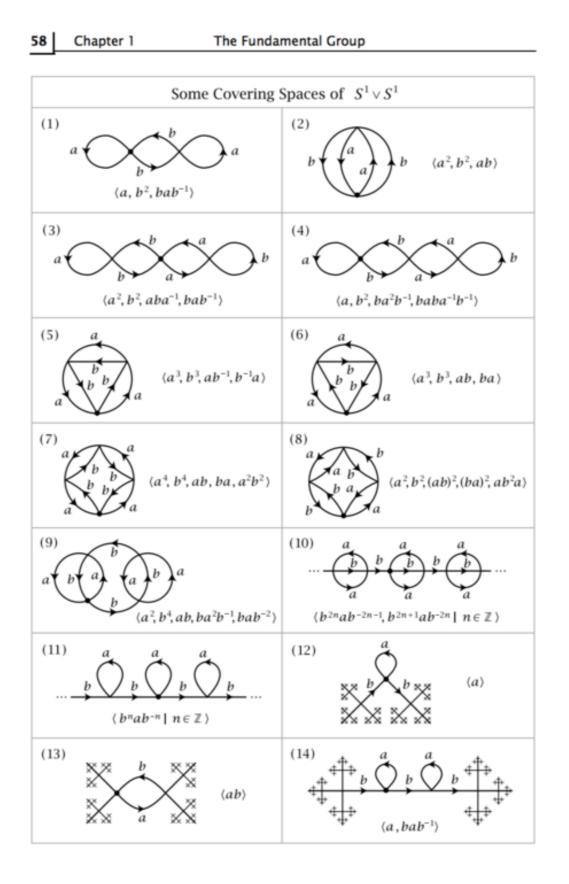




Sometimes people do illustrate mathematics beautifully.



## illustrating groups (with their notations)



Source: Algebraic Topology Hatcher

## composing different types of diagrams

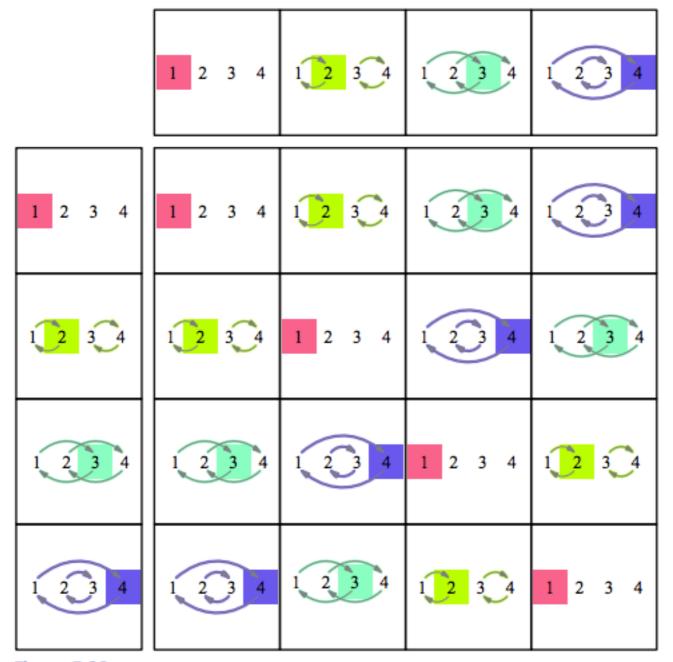


Figure 5.32. A multiplication table made up of the permutations created in Figure 5.31. Each cell highlights the destination to which the permutation sends 1, using the corresponding color from Figure 5.31, emphasizing what the colors of the arrows already showed: The two tables contain the same pattern.

## Problem:

Drawing good diagrams requires a tremendous amount of effort and expertise.

People have very powerful facilities for taking in information visually... and thinking spatially.

William Thurston Fields medalist People have very powerful facilities for taking in information visually... and thinking spatially.

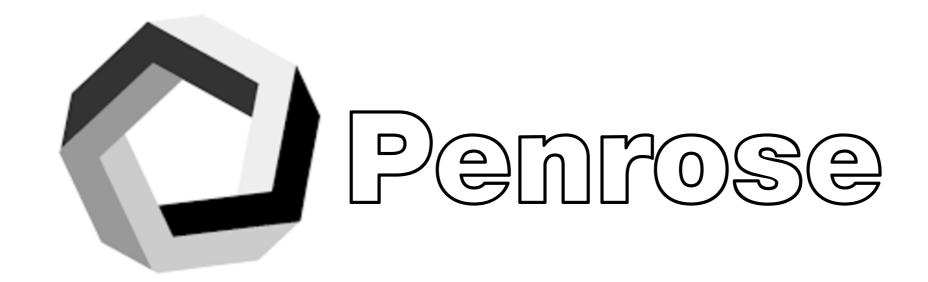
On the other hand, they do not have a very good built-in facility for inverse vision, that is, turning an internal spatial understanding back into a two-dimensional image.

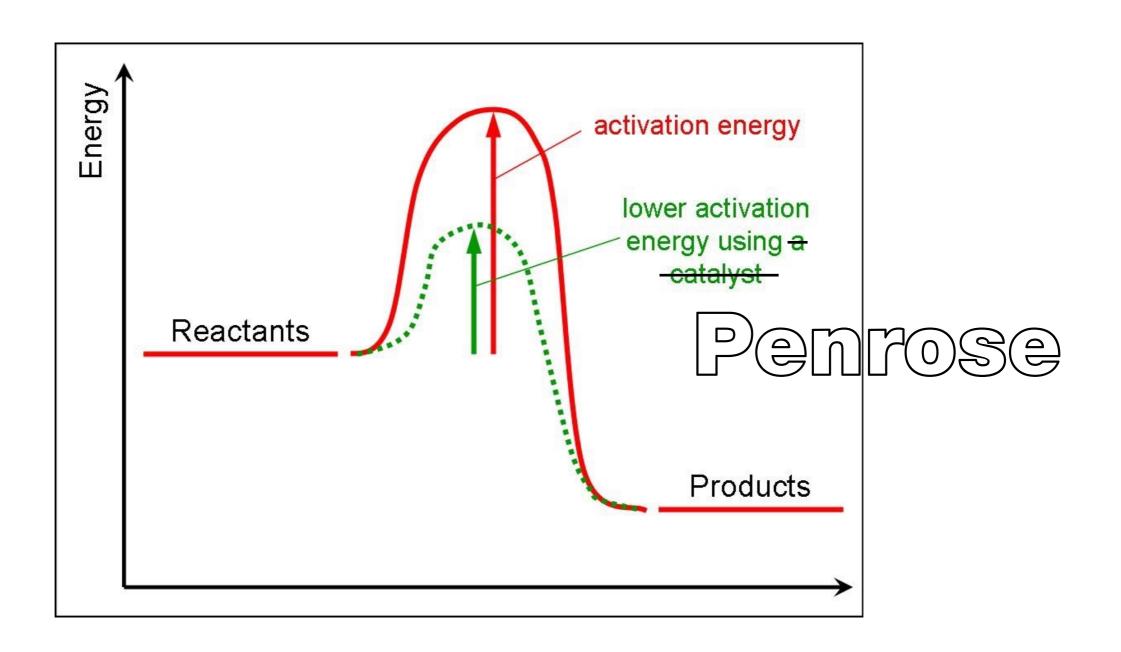
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On the other hand, they do not have a very good built-in facility for inverse vision, that is, turning an internal spatial understanding back into a two-dimensional image.

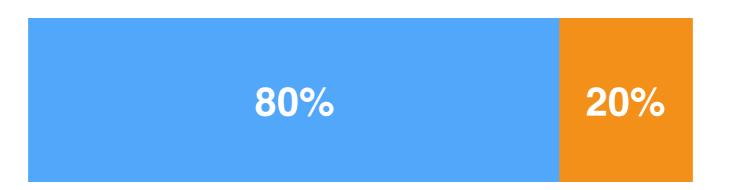
Consequently, mathematicians usually have fewer and poorer figures in their papers and books than in their heads.

William Thurston Fields medalist What if mathematicians had *more* and *richer* figures in their papers and books than in their heads?





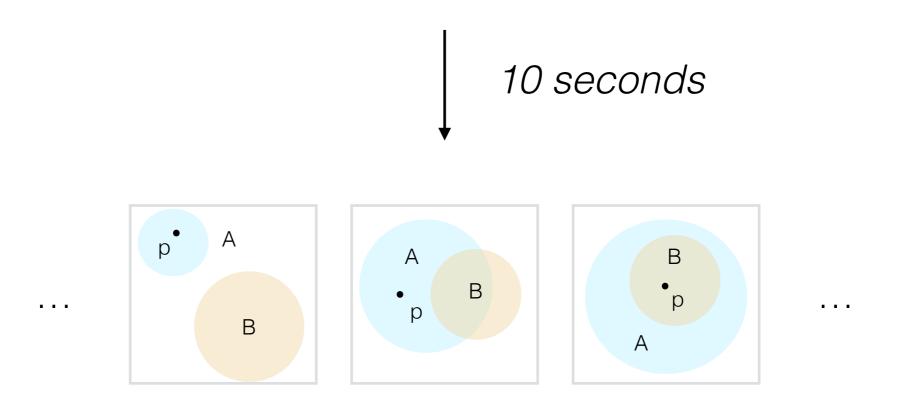
## Casual users



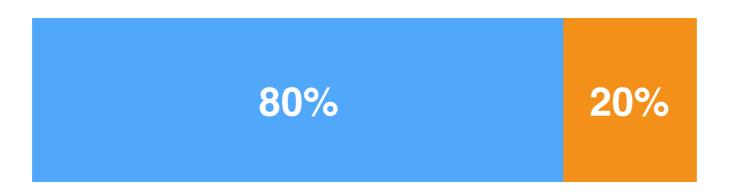
"Let A and B be sets, and p be a point in A."

10 seconds

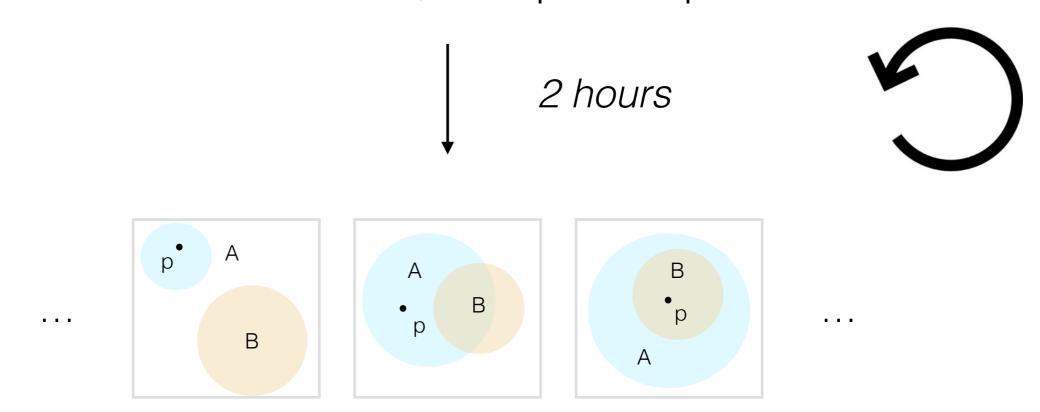
"Let A and B be sets, and p be a point in A."



### Power users



"Let A and B be sets, and p be a point in A."

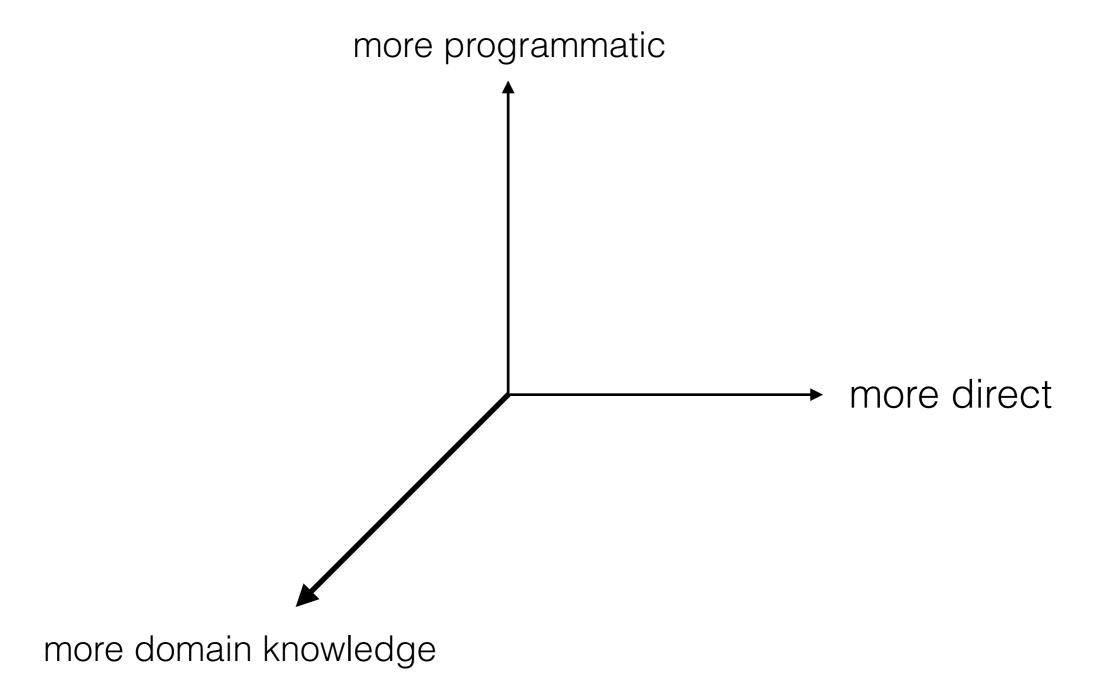


# Existing solutions for diagramming

# Factors in choosing a tool:

Programmatic, direct manipulation, domain knowledge

## The design space



## Writing a program for a diagram

$$K_4 = \langle a, b \mid a^2 = b^2 = (ab)^2 = e \rangle$$

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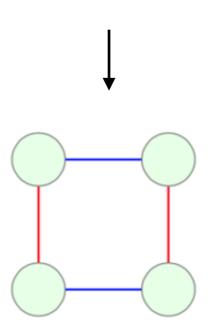


Figure 1: Cayley diagram of the Klein 4-group

Figure 1: Cayley diagram of the Klein 4-group

$$K_4 = \langle a, b \mid a^2 = b^2 = (ab)^2 = e \rangle$$

```
% node style
circ/.style={circle, draw=black!60, fill=green!10, minimum size=7mm}],
% nodes
\node[circ] (gen_bl) {};
\node[circ] (gen_tl) [above=of gen_bl] {};
\node[circ] (gen_br) [right=of gen_bl] {};
\node[circ] (gen_tr) [above=of gen_br, right=of gen_tl] {};
% \node[circ] (gen_bl);
% \node[circ] (gen_tl) [above=of gen_bl]; % {2} label
% \node[circ] (gen_tr) [above=of gen_br, right=of gen_tl];
% \node[circ] (gen br) [right=of gen bl];
% draw lines between nodes
\draw[-][red, thick] (gen tl.south) -- (gen bl.north);
\draw[-][blue, thick] (gen_tl.east) -- (gen_tr.west);
\draw[-][blue, thick] (gen_bl.east) -- (gen_br.west);
\draw[-][red, thick] (gen_tr.south) -- (gen_br.north);
```

*TikZ* 

Figure 1: Cayley diagram of the Klein 4-group

node style

$$K_4 = \langle a, b \mid a^2 = b^2 = (ab)^2 = e \rangle$$

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% node style
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\draw[-][red, thick] (gen_tr.south) -- (gen_br.north);
```

*TikZ* 

node style node positions

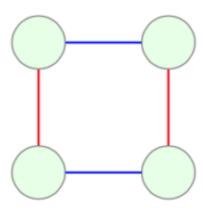


Figure 1: Cayley diagram of the Klein 4-group

$$K_4 = \langle a, b \mid a^2 = b^2 = (ab)^2 = e \rangle$$

TikZ



node style node positions

line style + positions

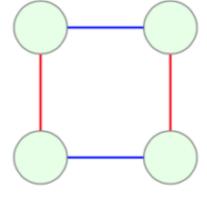
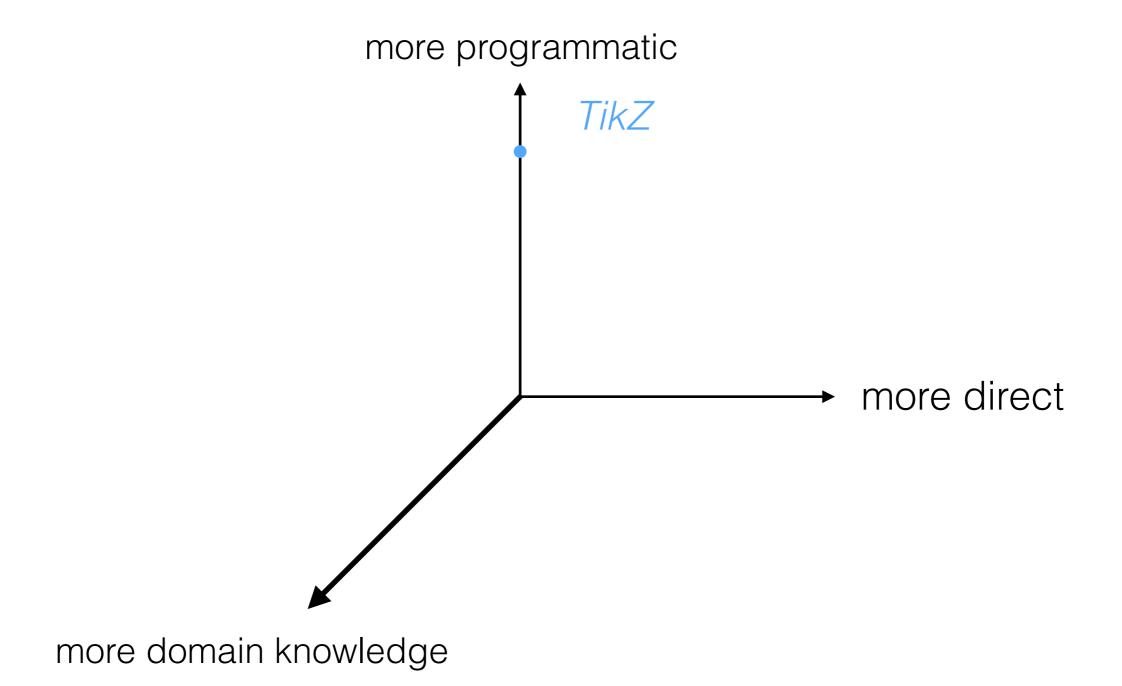
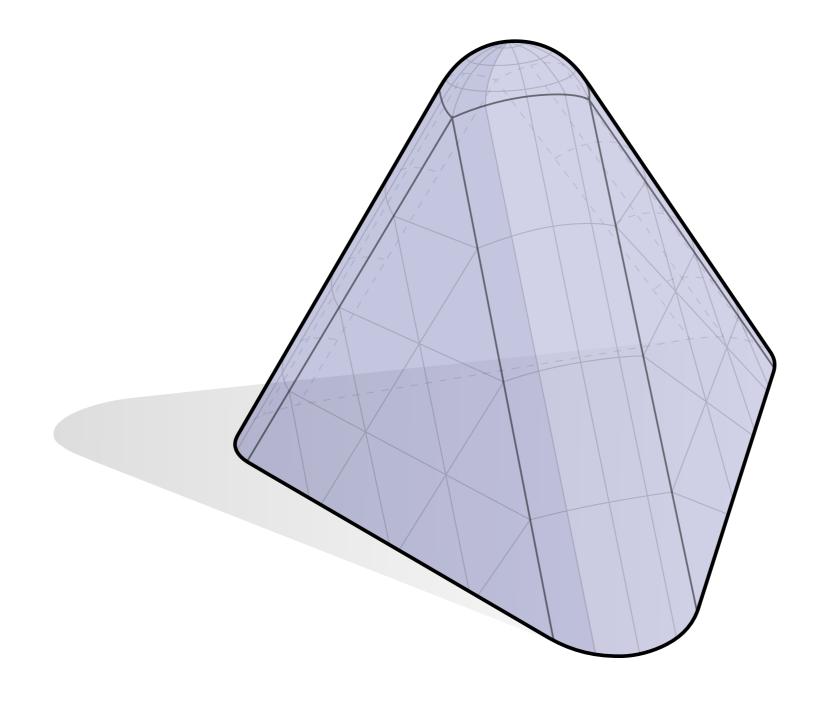


Figure 1: Cayley diagram of the Klein 4-group

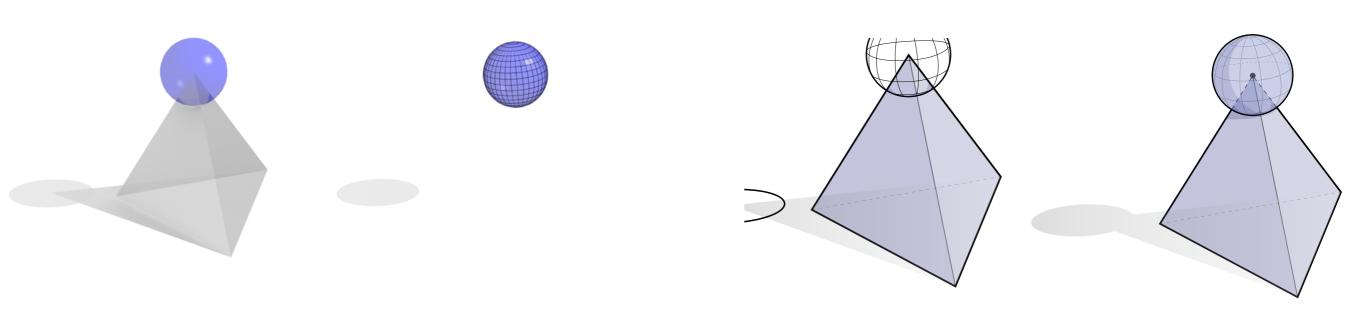
## The design space

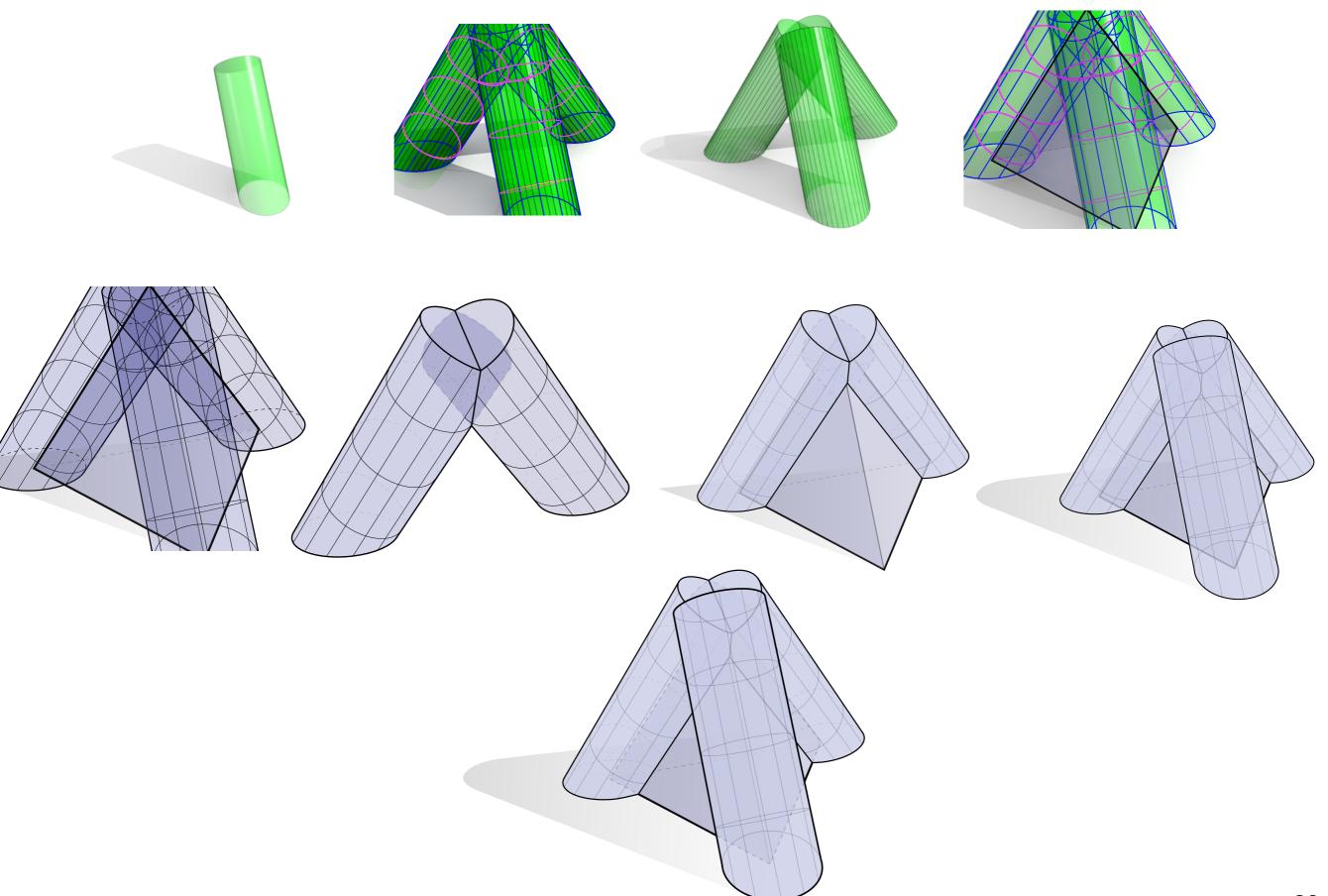


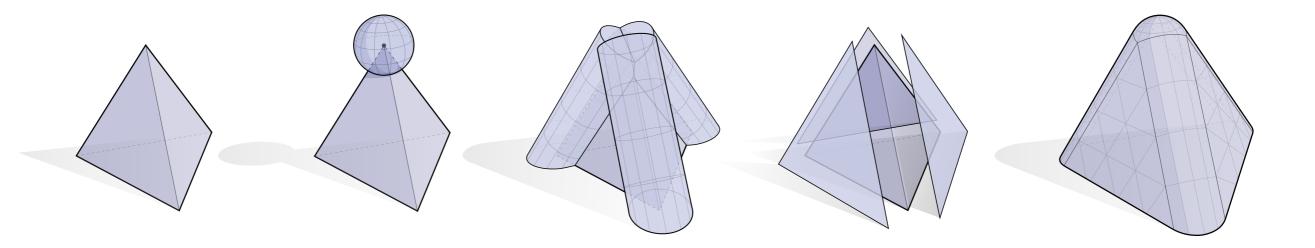
# Creating a diagram via a GUI

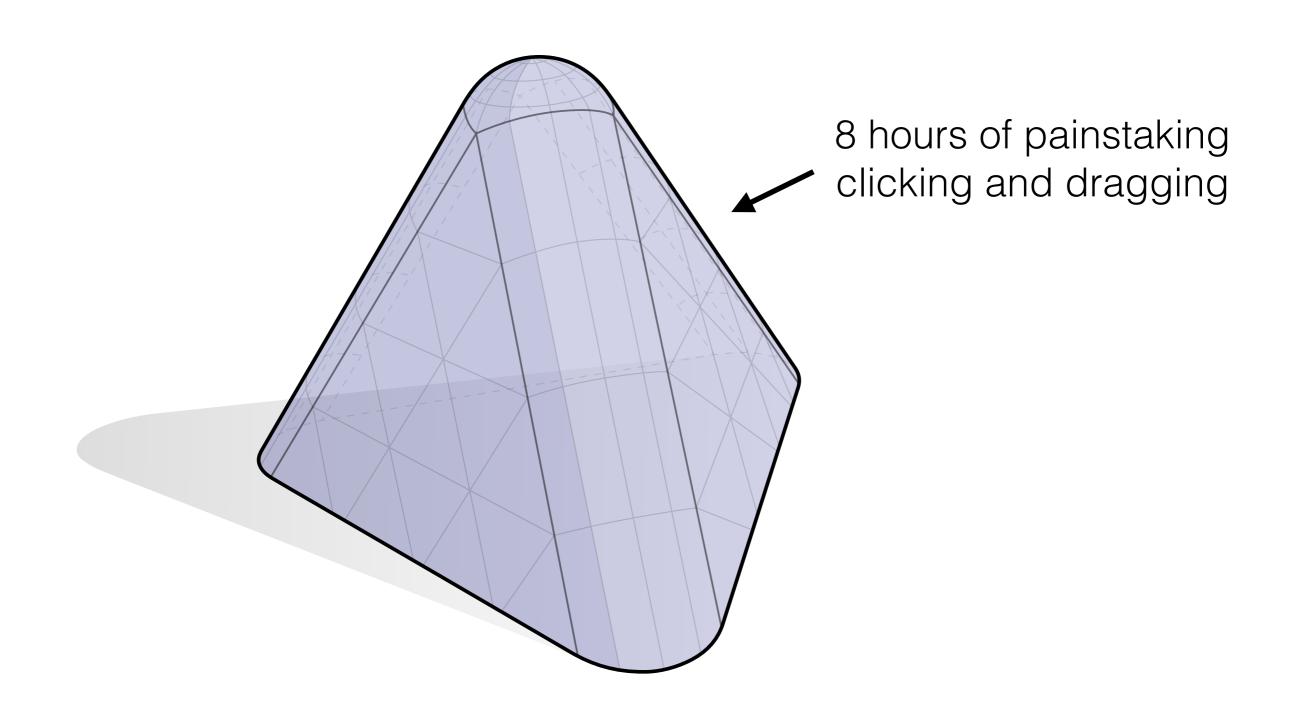


Source: Keenan Crane 37

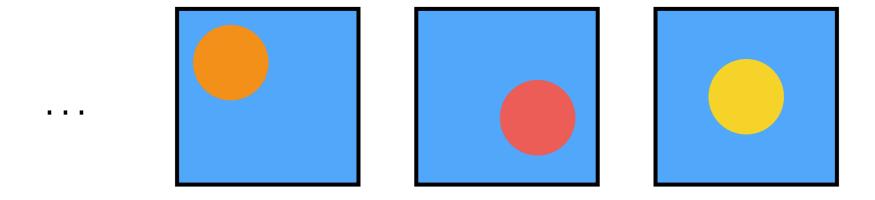


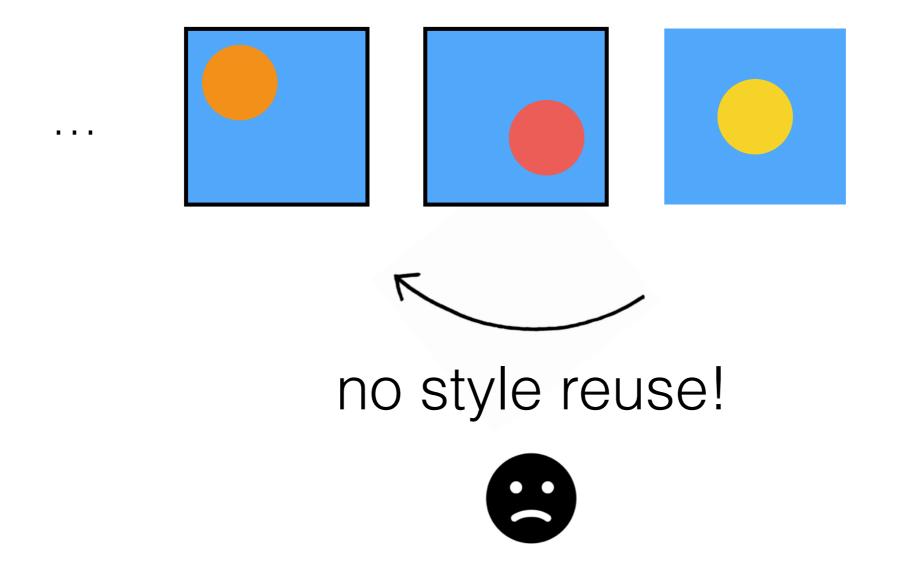




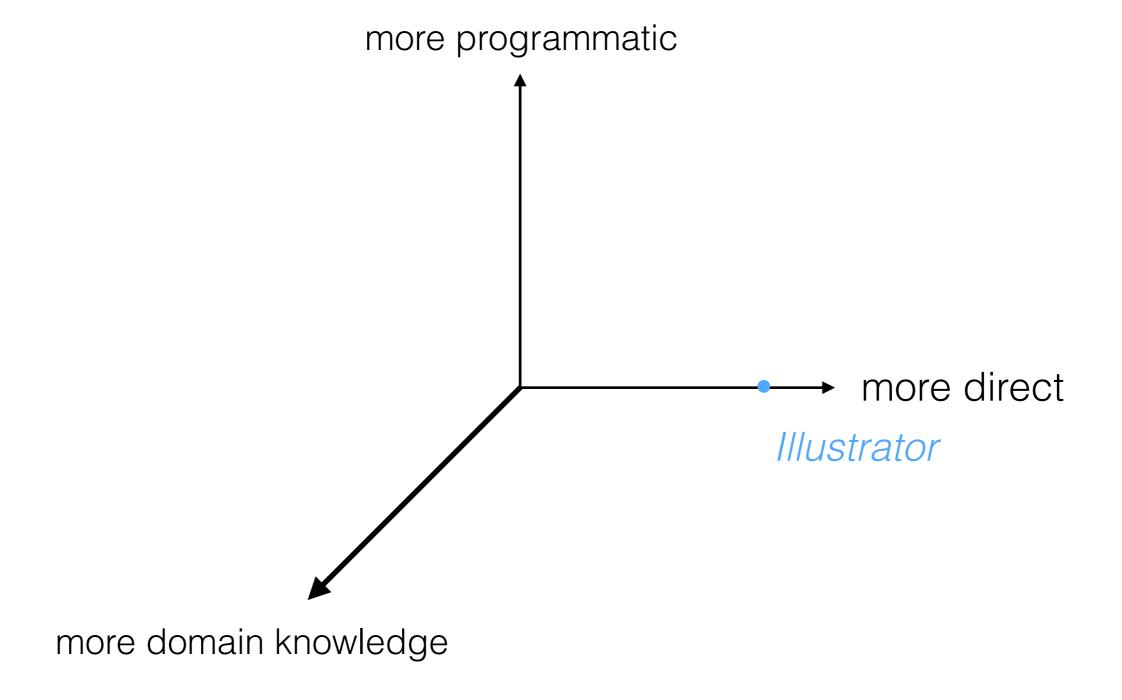


Source: Keenan Crane

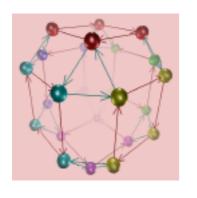




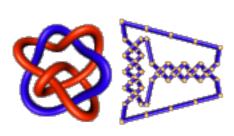
# The design space



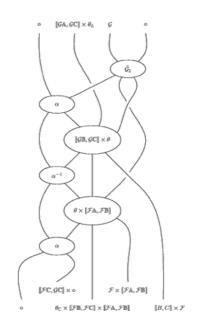
Burn it all down and write your own software or DSL?



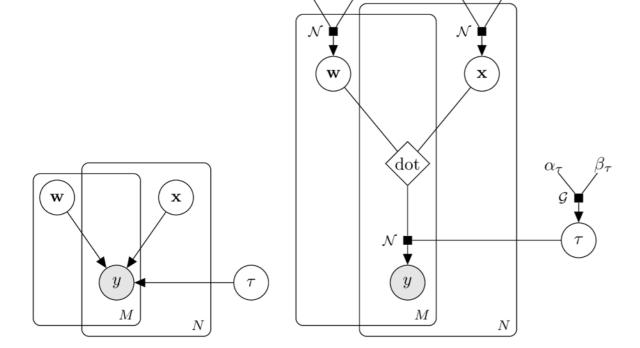




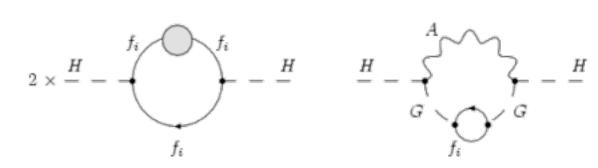
KnotPlot



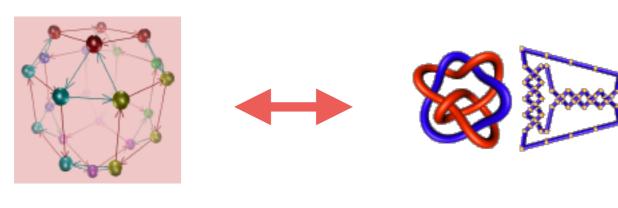
strid



BayesNet

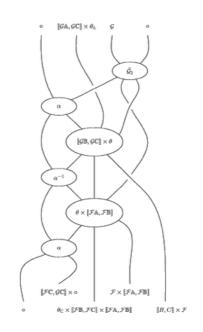


JaxoDraw

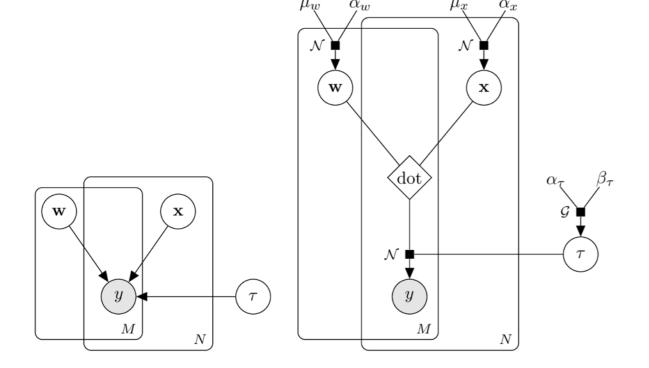


**Group Explorer** 

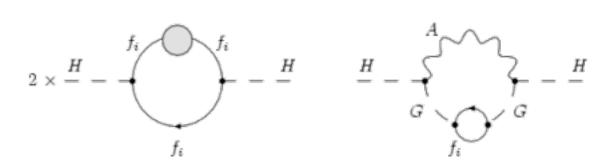
KnotPlot



strid

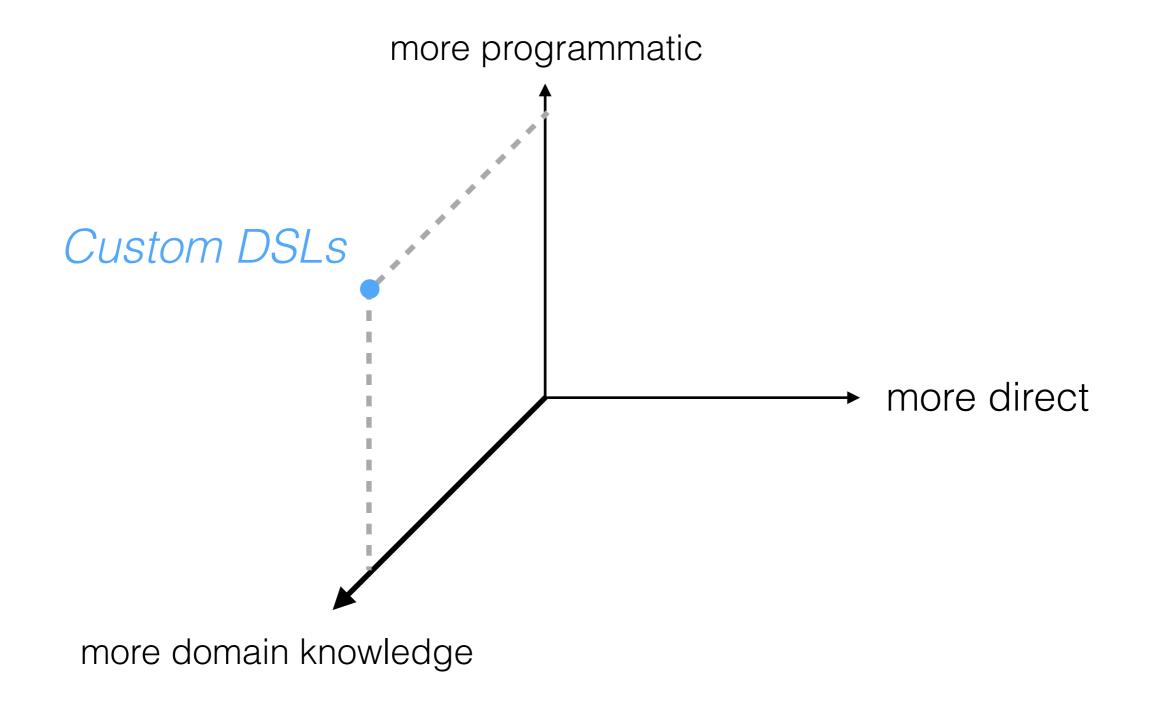


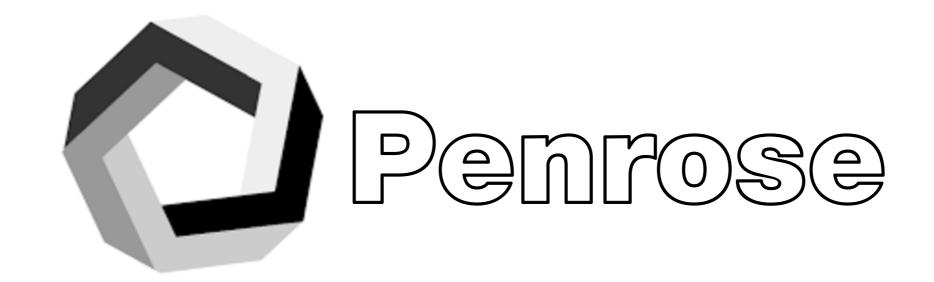
BayesNet



JaxoDraw

# The design space





"Let A and B be sets, and p be a point in A."

"Let A and B be sets, and p be a point in A."

#### Notation

Set A Set B Point p  $p \in A$ 

"Let A and B be sets, and p be a point in A."

"Let A and B be sets, and p be a point in A."

Substance View Style

"Let A and B be sets, and p be a point in A."

Substance View Style

Set A
Set B
Point p
In p A

"Let A and B be sets, and p be a point in A."

Style

Substance View

Set A
Set B
Point p
In p A

Dimension 2
Coordinates None

"Let A and B be sets, and p be a point in A."

Substance

View

Style

Set A
Set B
Point p
In p A

Dimension 2
Coordinates None

Shape Set Circle
Color A Blue
Color B Orange
Shape Point Circle
Label All Auto

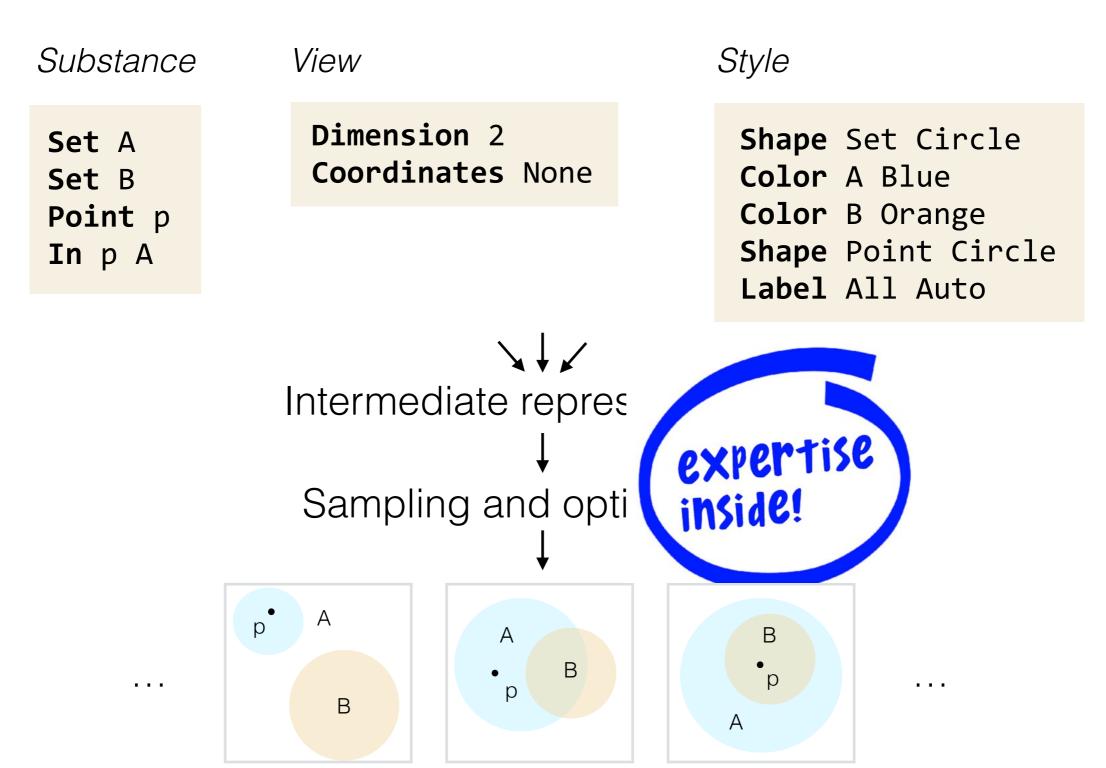
"Let A and B be sets, and p be a point in A."

Substance View Style Dimension 2 **Shape** Set Circle Set A Coordinates None Color A Blue Set B Color B Orange **Point** p **Shape** Point Circle In p A Label All Auto 1 Intermediate representations Sampling and optimization

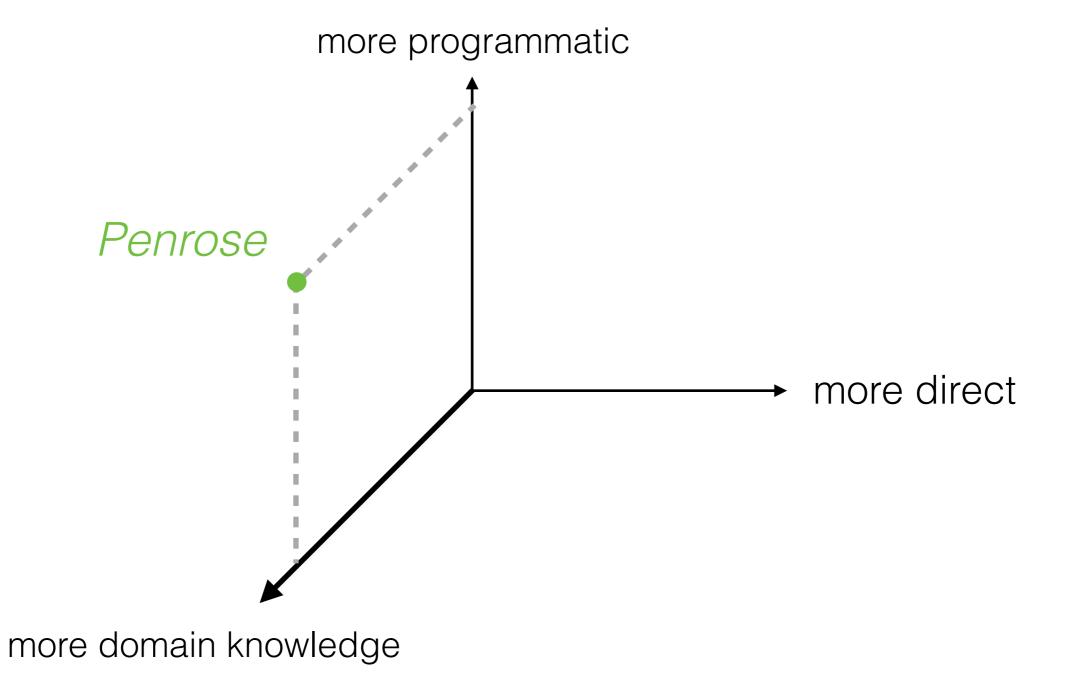
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Substance View Style Dimension 2 **Shape** Set Circle Set A Coordinates None Color A Blue Set B Color B Orange **Point** p **Shape** Point Circle In p A Label All Auto Intermediate representations Sampling and optimization Α

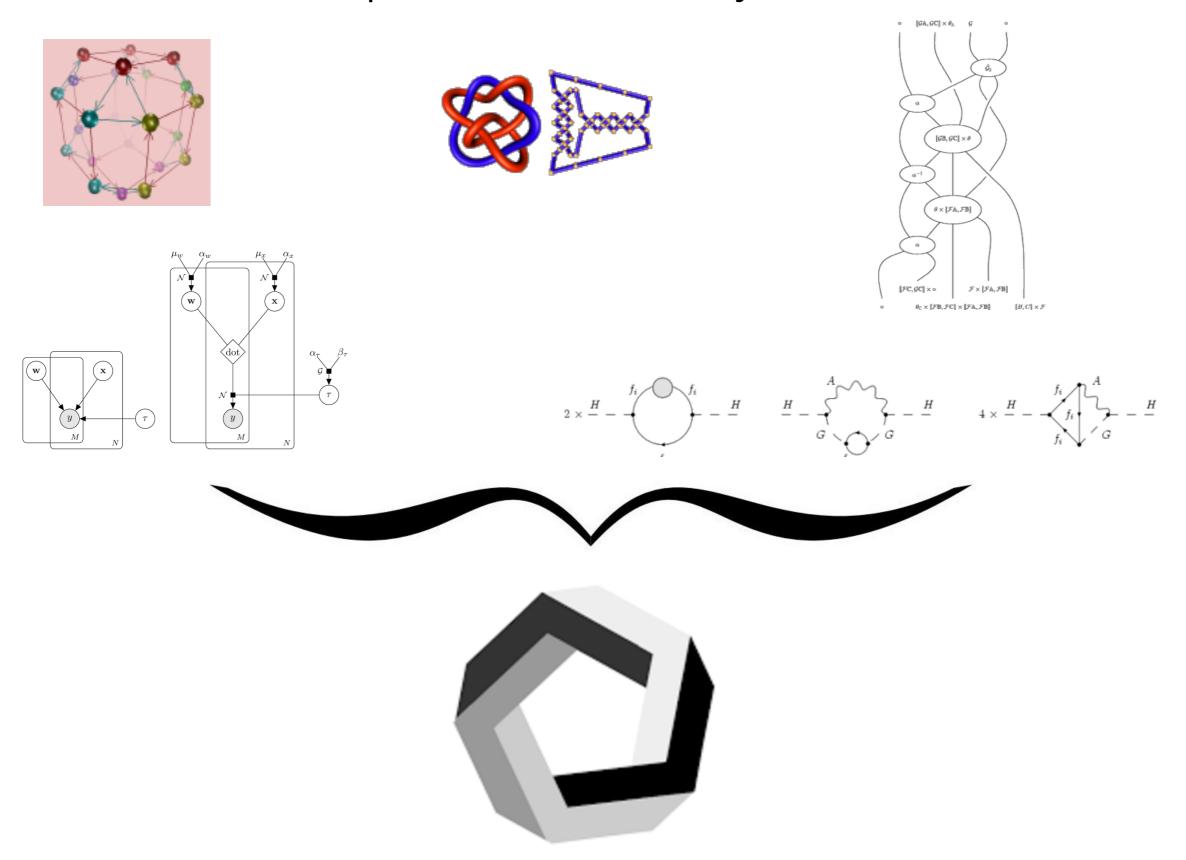
"Let A and B be sets, and p be a point in A."



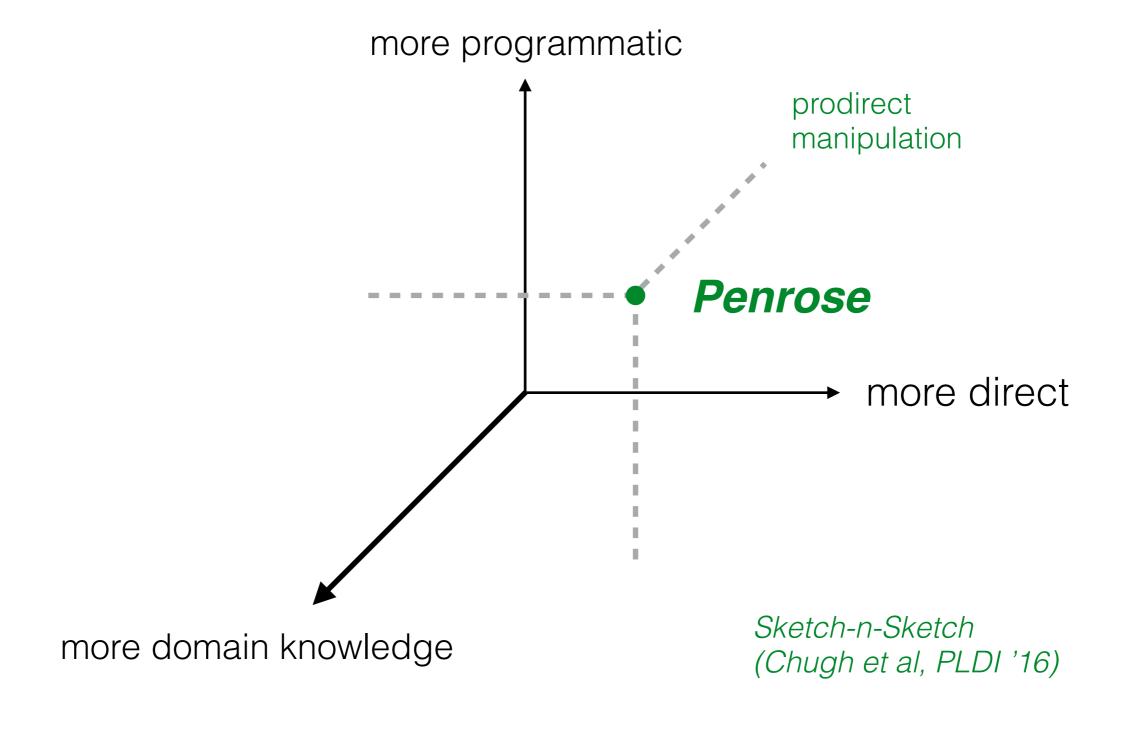
# The design space



## ...plus extensibility!



# The design space



# An example of our vision

I'm demonstrating a **mockup** of a workflow that I hope to demo live in a few months.



We have a very simple prototype for set theory, which I won't demo.



Domain:
Set theory →
Point-set topology

# Phases: Implementing the DSL → using the DSL

# Implementing the DSL

**Declarations** 

**Declarations** 

data SubDecl = Decl Object

#### Declarations

```
data SubDecl = Decl Object
data Object = OS Set | OP Point
data Set = Set' String SetType
```

data Pt = Pt' String

#### Declarations

```
data SubDecl = Decl Object
data Object = OS Set | OP Point
data Set = Set' String SetType
data SetType = Open | Closed | Clopen | Unspecified
data Pt = Pt' String
```

#### Implementing a set theory DSL

#### Declarations

```
data SubDecl = Decl Object
data Object = OS Set | OP Point
data Set = Set' String SetType
data SetType = Open | Closed | Clopen | Unspecified
data Pt = Pt' String
```

#### Constraints

#### Implementing a set theory DSL

#### Declarations

```
data SubDecl = Decl Object
data Object = OS Set | OP Point
data Set = Set' String SetType
data SetType = Open | Closed | Clopen | Unspecified
data Pt = Pt' String
```

#### Constraints

data Constraint = Subset String String

#### Implementing a set theory DSL

#### Declarations

```
data SubDecl = Decl Object
data Object = OS Set | OP Point
data Set = Set' String SetType
data SetType = Open | Closed | Clopen | Unspecified
data Pt = Pt' String
```

#### Constraints

### Set theory DSL

Set A
Set B
Point p
In p A

View

Dimension 2
Coordinates None

Style

**Shape** Set Circle

Color A Blue

Color B Orange

**Shape** Point Circle

Label All Auto

#### Extending the set theory DSL

#### Declarations

```
data SubDecl = Decl Object
data Object = OS Set | OP Point | OM Map
data Set = Set' String SetType
data SetType = Open | Closed | Clopen | Unspecified
data Pt = Pt' String
data Map = Map' String String String
```

#### Constraints

# Using the DSL

# Open sets, closed sets, continuous maps



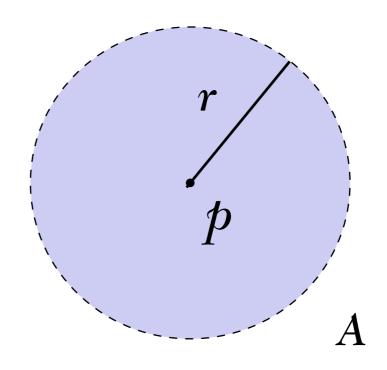
A First Course in Geometric Topology and Differential Geometry

Ethan D. Bloch

Birkhäuser

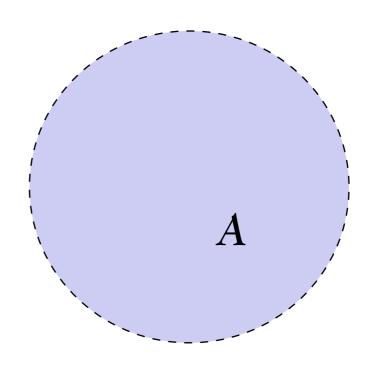
# An open ball in R<sup>2</sup>

$$A = O_r(p, R^n)$$



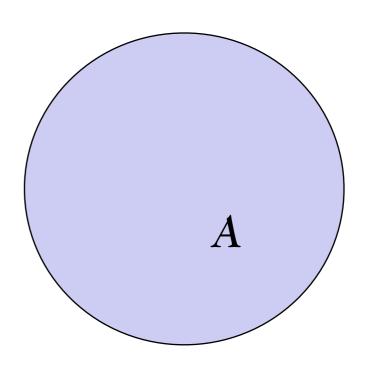
# An open ball in R<sup>2</sup>

OpenBall A



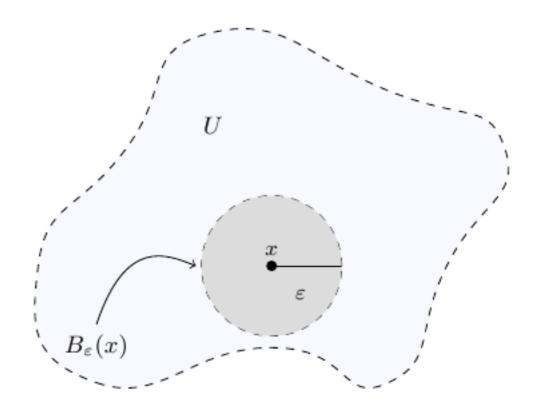
## An closed ball in R<sup>2</sup>

ClosedBall A



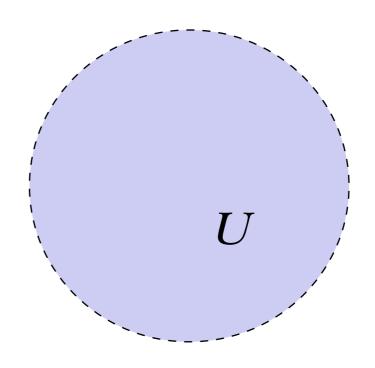
# An open set in R<sup>2</sup>

OpenSet U



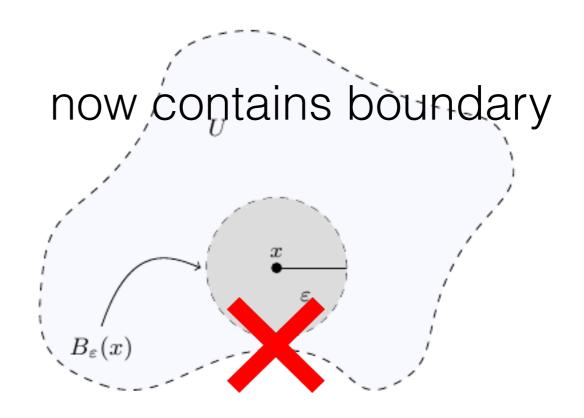
# An open set in R<sup>2</sup>, simplified

OpenSet U



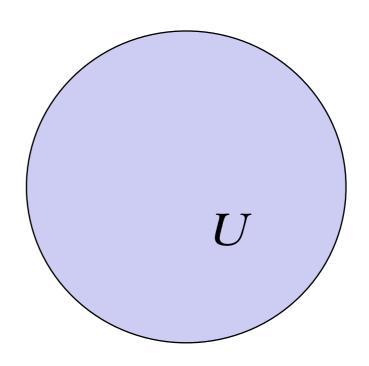
#### An closed set in R<sup>2</sup>

ClosedSet U

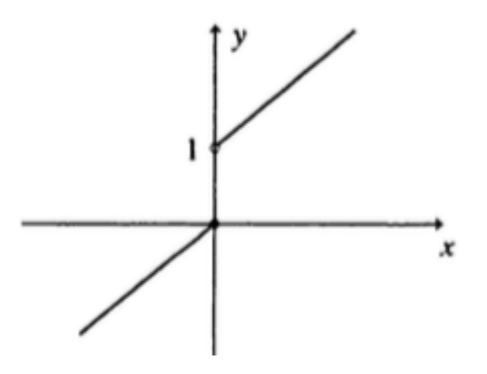


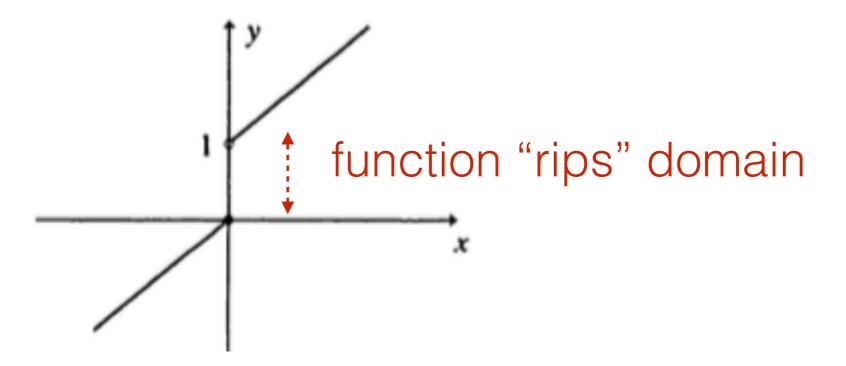
# An closed set in R<sup>2</sup>, simplified

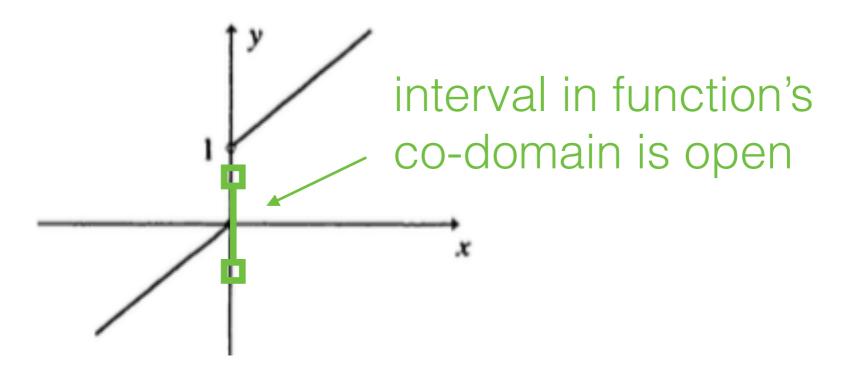
ClosedSet U

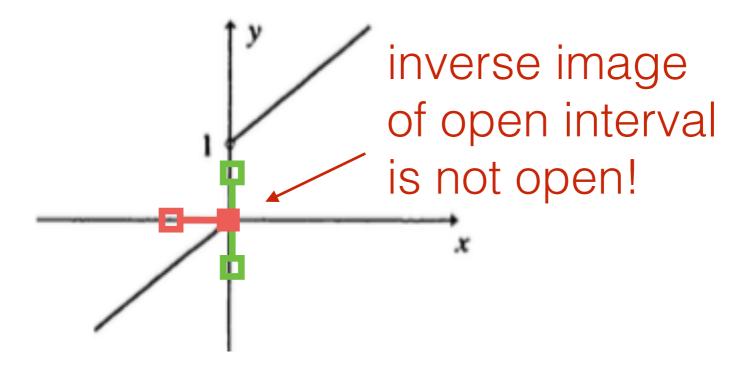


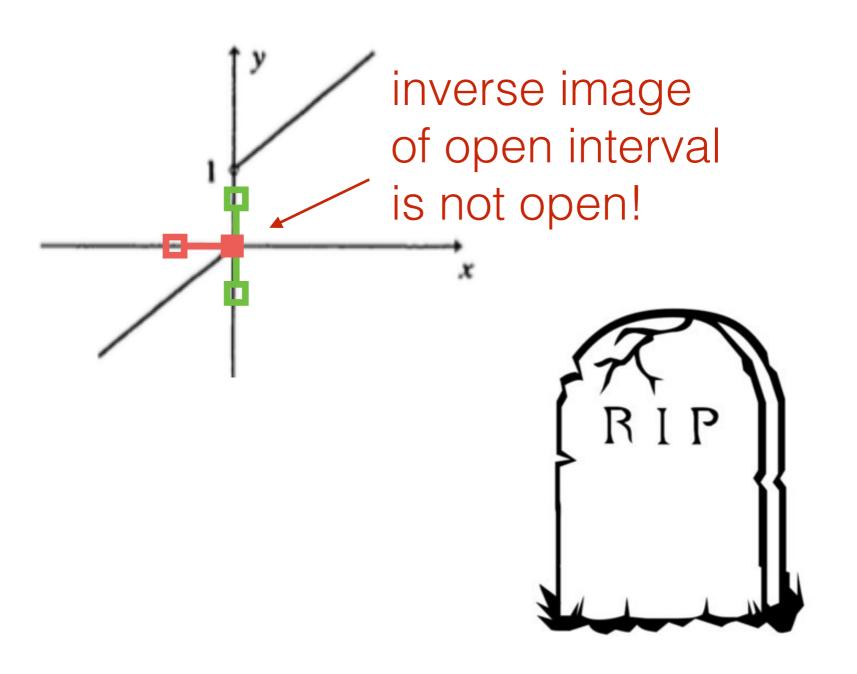
# Continuous maps











# Generalizing the intuition

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**Definition.** Let  $A \subset \mathbb{R}^n$  and  $B \subset \mathbb{R}^m$  be sets, and let  $f: A \to B$  be a map. The map f is **continuous** if for every open subset  $U \subset B$ , the set  $f^{-1}(U)$  is open in A.  $\diamondsuit$ 

#### Picture this!



#### Substance

```
Set A
Set B
Set R^n
Set R^m
Subset A R^n
Subset B R^m
Map f A B
OpenSet U
Subset U B
Set f^{-1}(U)
Subset f^{-1}(U) A
```

#### Substance

```
Set A
Set B
Set R^n
Set R^m
Subset A R^n
Subset B R^m
Map f A B
OpenSet U
Subset U B
Set f^{-1}(U)
Subset f^{-1}(U) A
```

declarations

#### Substance

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Set B
Set R^n
Set R^m
Subset A R^n
Subset B R^m
Map f A B
OpenSet U
Subset U B
Set f^{-1}(U)
Subset f^{-1}(U) A
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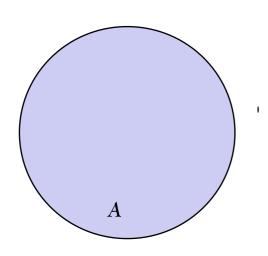
constraints

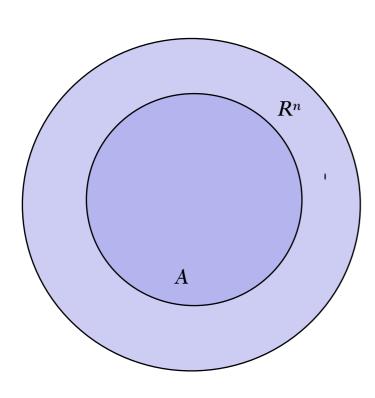
#### Substance

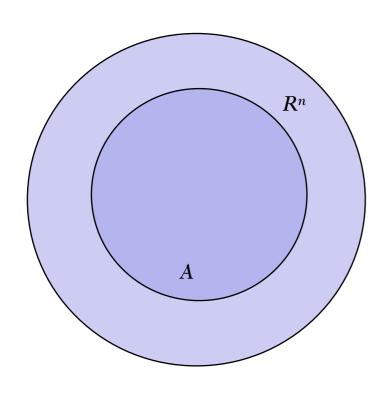
Set A
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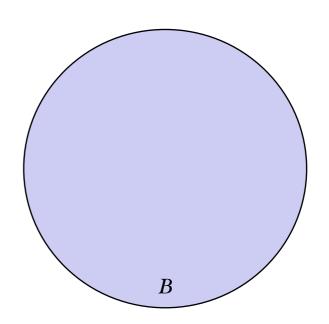
#### Style

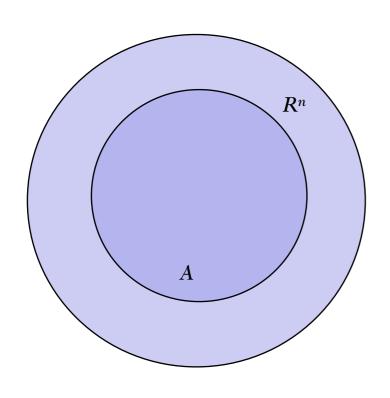
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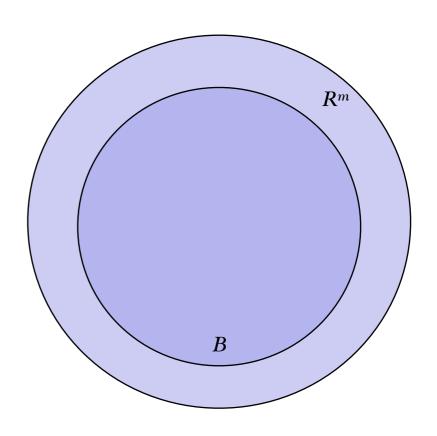


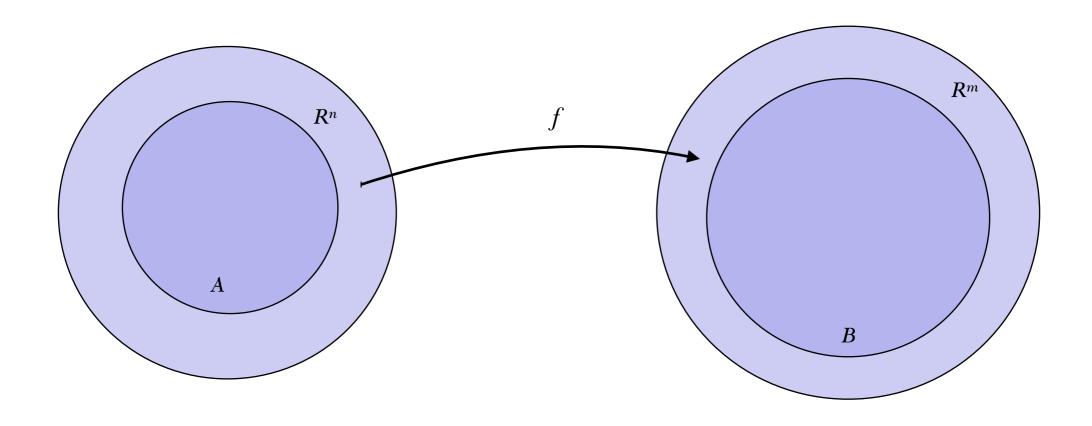


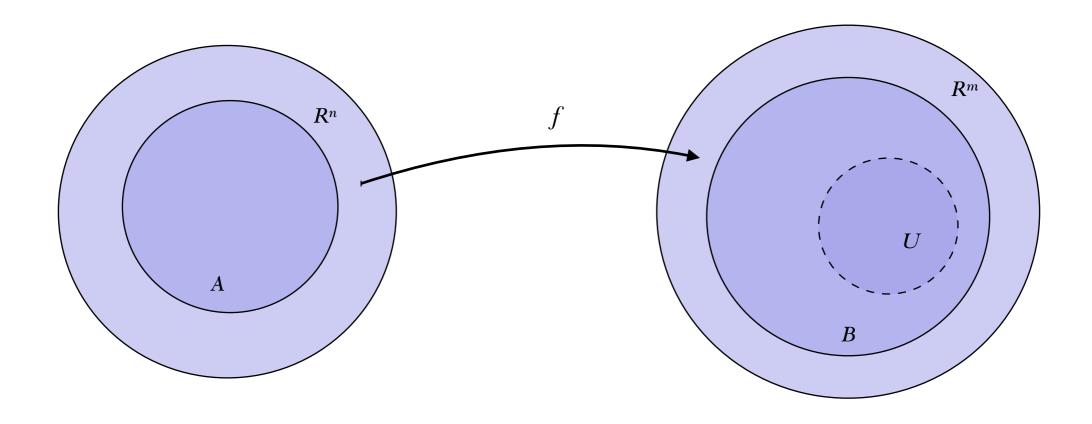


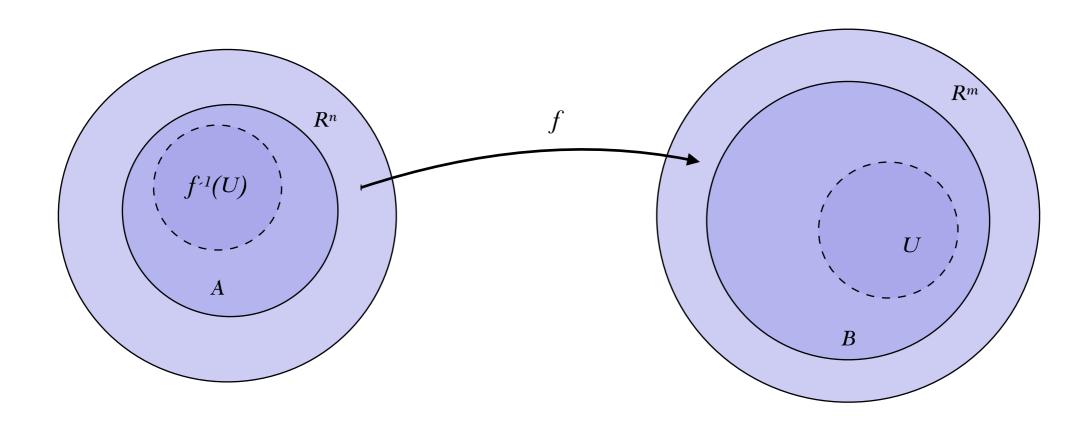


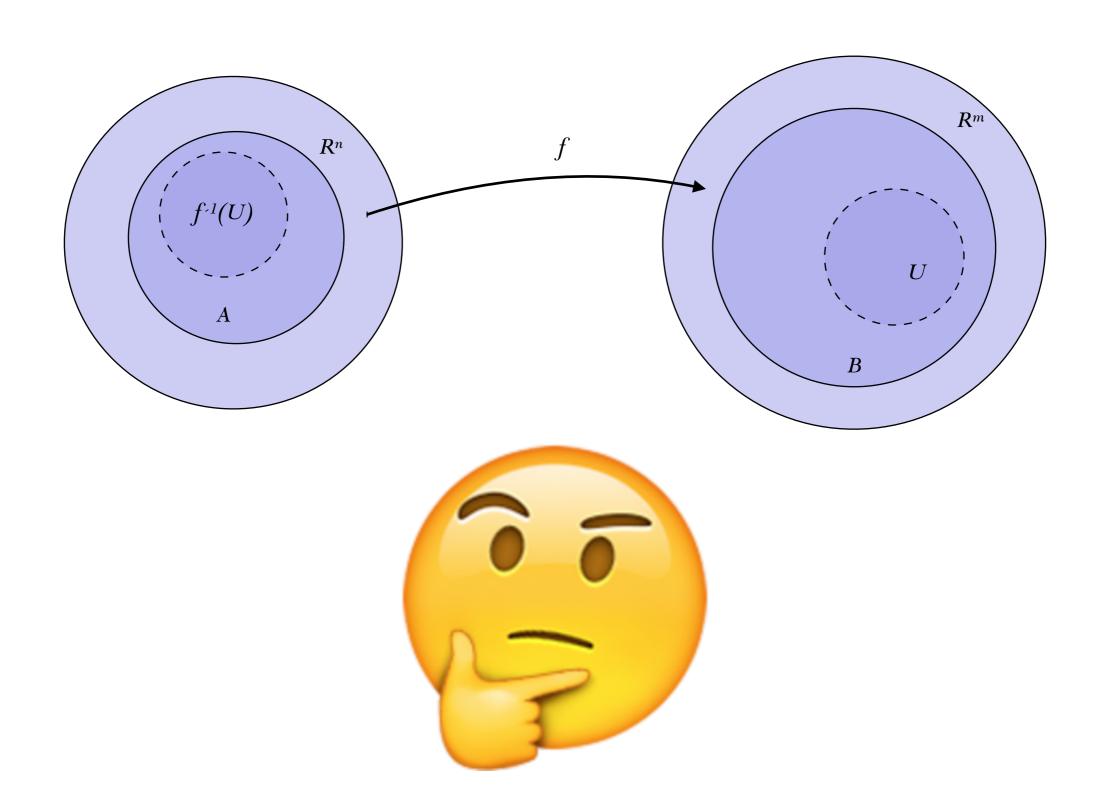










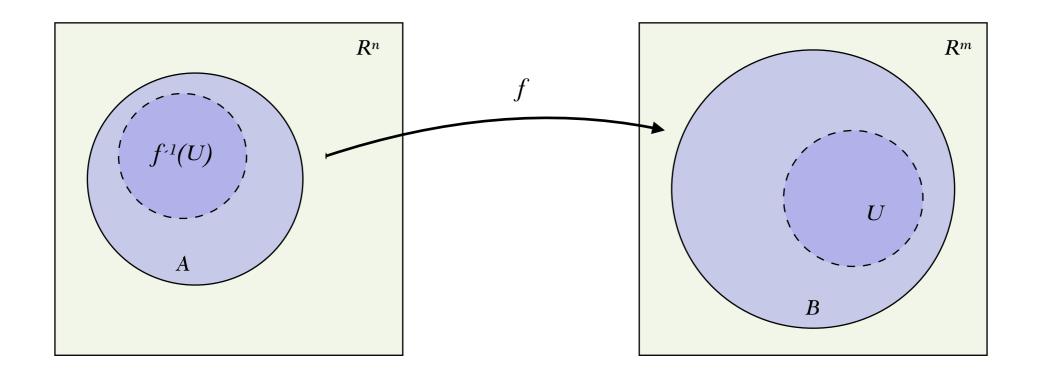


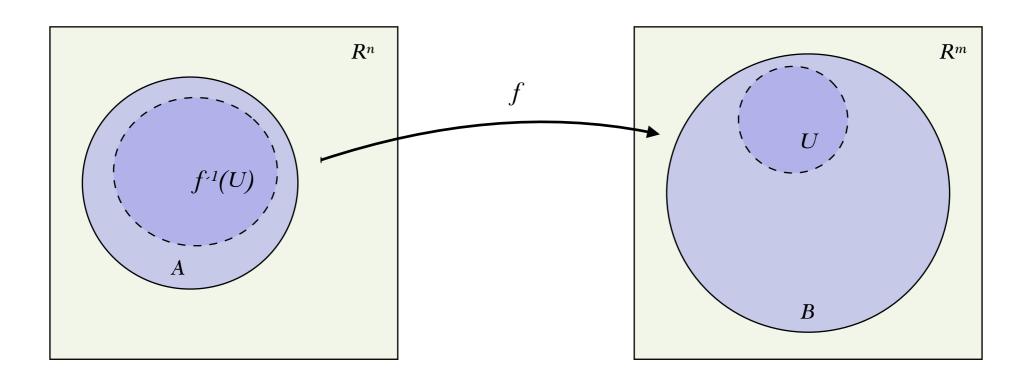
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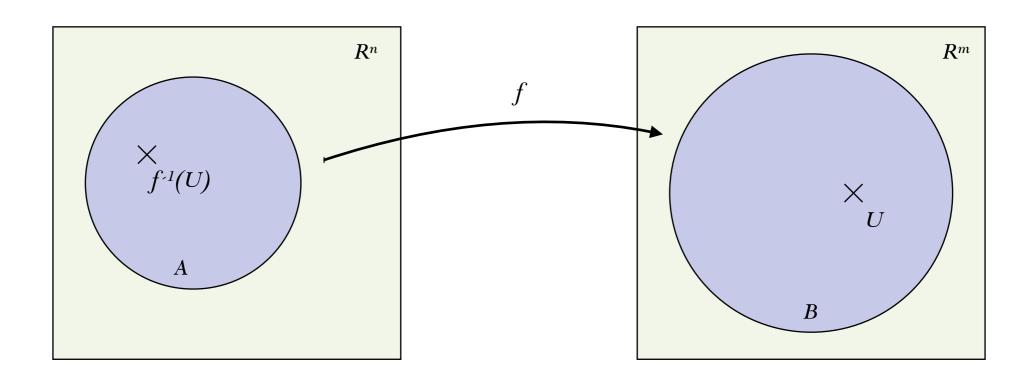
### Style

```
Style All Auto
Shape R^n Square
Shape R^m Square
Color R^n Yellow
Color R^m Yellow
```





animation, interactivity, fuzzing



definition holds for the "edge case" of empty sets

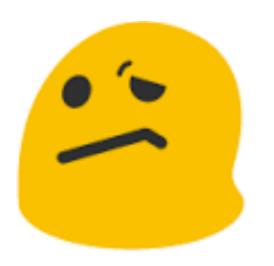
animation, interactivity, fuzzing

# A harder definition...

- (1) The map f is continuous.
- (2) For every point  $p \in A$ , and every open subset  $U \subset B$  containing f(p), there is an open subset  $V \subset A$  containing p such that  $f(V) \subset U$ .
- (3) For every point  $p \in A$  and every number  $\epsilon > 0$ , there is a number  $\delta > 0$  such that if  $x \in A$  and  $||x p|| < \delta$  then  $||f(x) f(p)|| < \epsilon$ .

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### Picture this!



- (1) for every open subset  $U \subset B$ , the set  $f^{-1}(U)$  is open in A.
- (2) For every point  $p \in A$ , and every open subset  $U \subset B$  containing f(p), there is an open subset  $V \subset A$  containing p such that  $f(V) \subset U$ .



(3) For every point  $p \in A$  and every number  $\epsilon > 0$ , there is a number  $\delta > 0$  such that if  $x \in A$  and  $||x - p|| < \delta$  then  $||f(x) - f(p)|| < \epsilon$ .

### Picture this!



- (1) for every open subset  $U \subset B$ , the set  $f^{-1}(U)$  is open in A.
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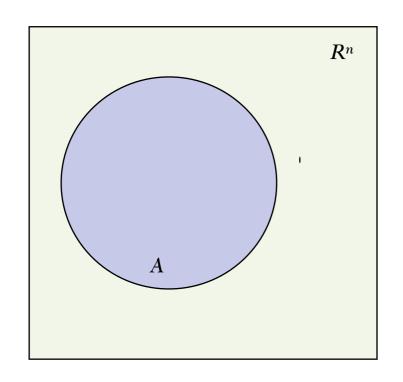
# Set A Set B Set R^n Set R^m Subset A R^n Subset B R^m Map f A B OpenSet U Subset U B

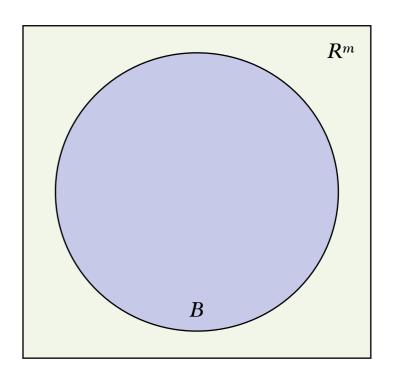
```
Point p
In p A
Point f(p)
In f(p) U
OpenSet V
Subset V A
In p V
Set f(V)
Subset f(V) U
In f(p) f(V)
```

### Style

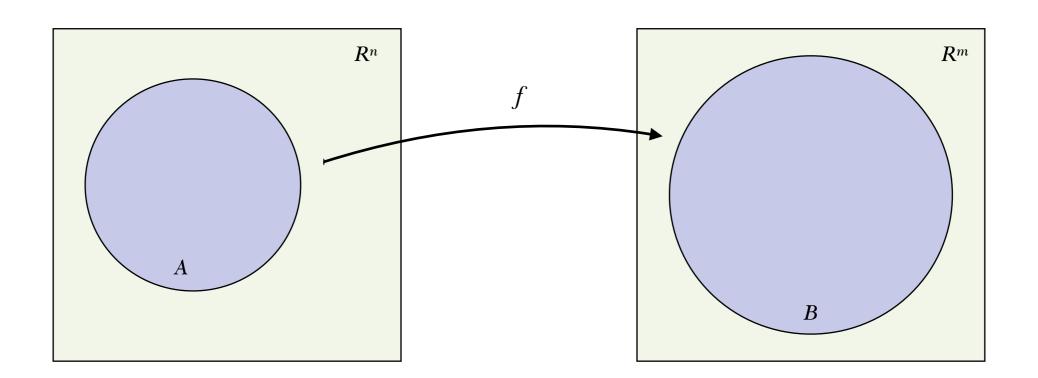
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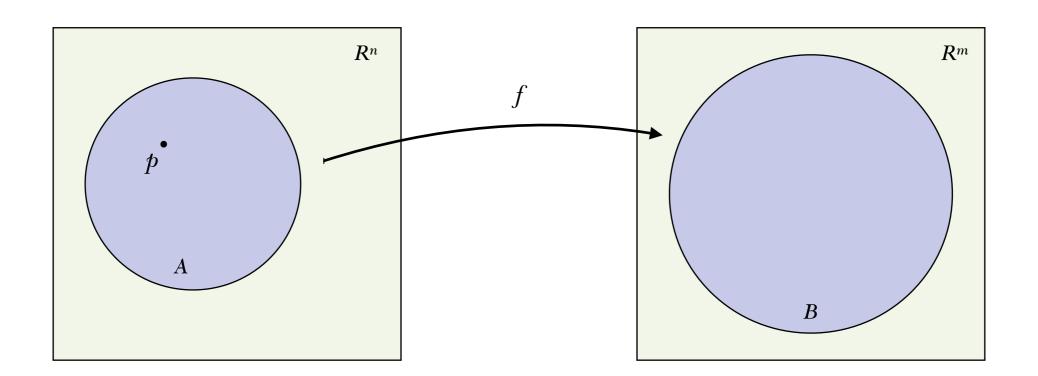




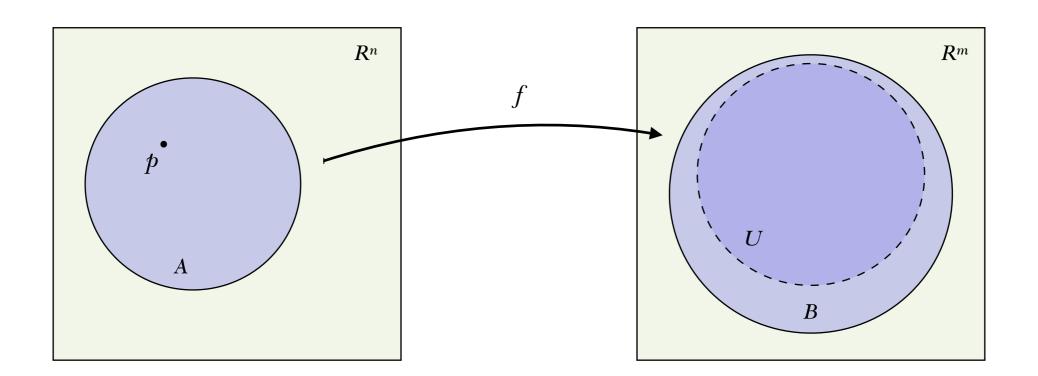
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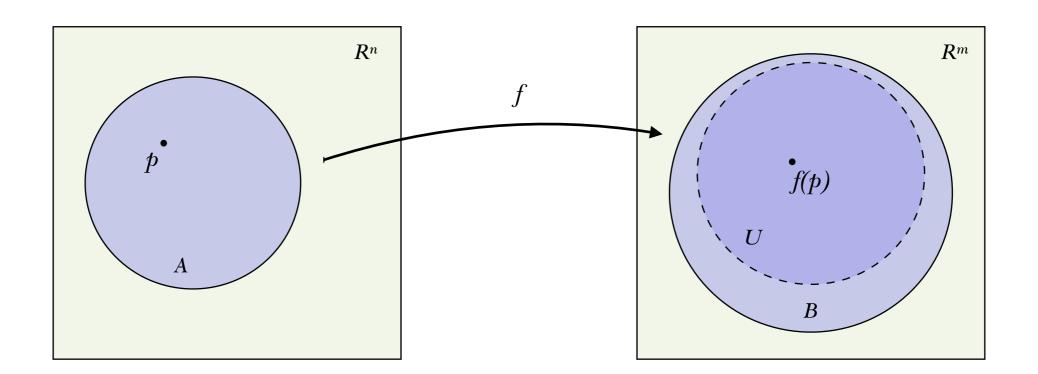
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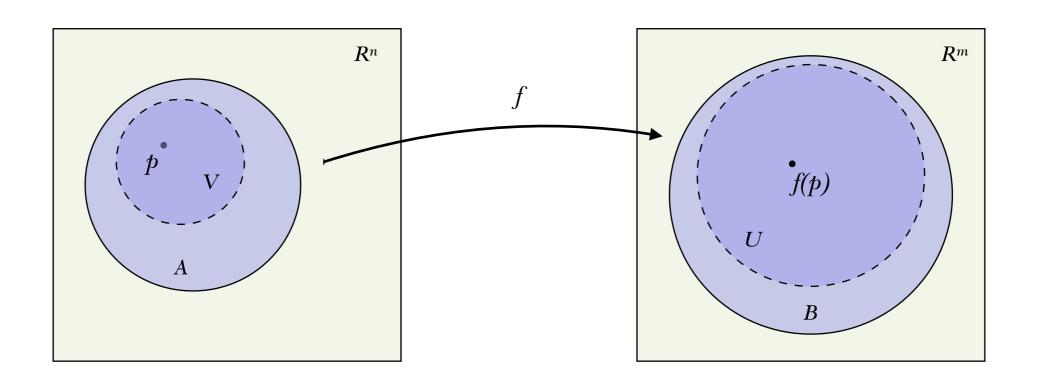
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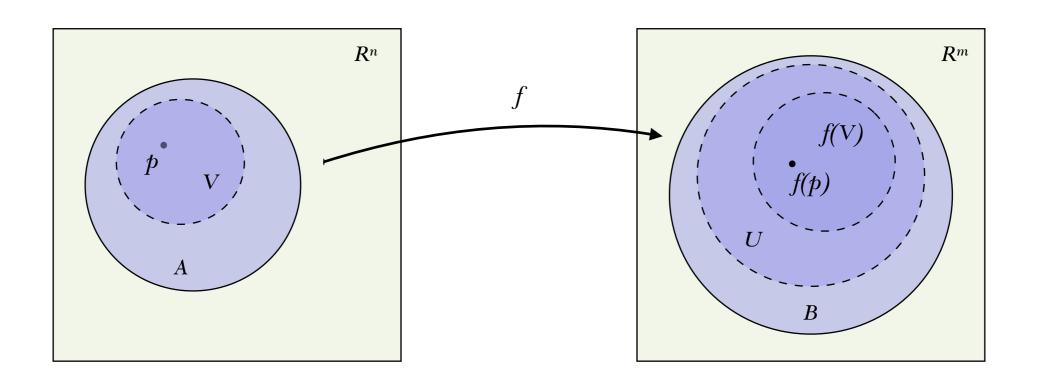
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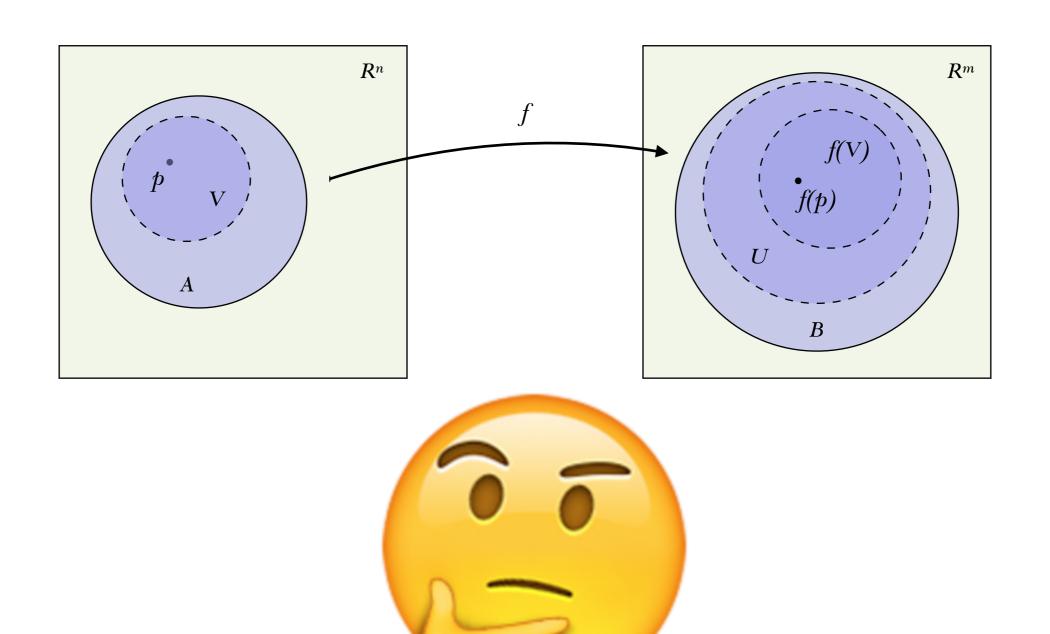
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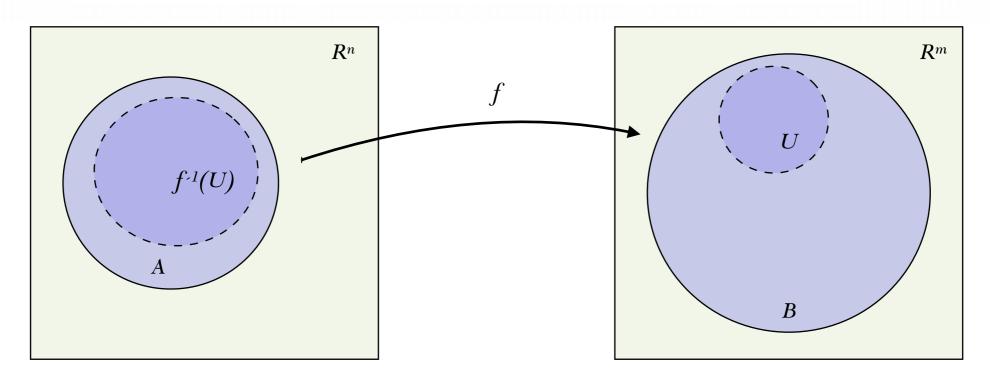


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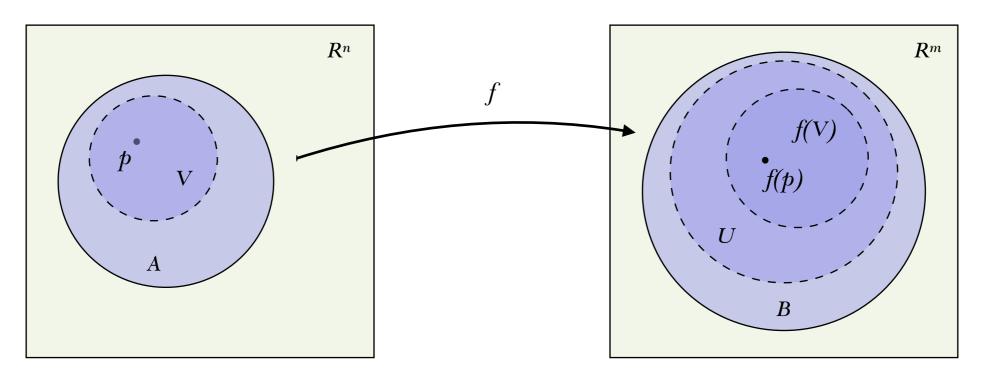


# How do the two definitions relate?

(1) for every open subset  $U \subset B$ , the set  $f^{-1}(U)$  is open in A.



(2) For every point  $p \in A$ , and every open subset  $U \subset B$  containing f(p), there is an open subset  $V \subset A$  containing p such that  $f(V) \subset U$ .



# Can prove (1) $\rightarrow$ (2) via diagram!

We now prove  $(1) \Rightarrow (2) \Rightarrow (3') \Rightarrow (1)$ .

(1)  $\Rightarrow$  (2). Let  $p \in A$  and  $U \subset B$  containing f(p) be given. By assumption, the map f is continuous, so that  $f^{-1}(U)$  is an open subset of A. Observe that  $p \in f^{-1}(U)$ . By the definition of openness there is thus some open ball of the form  $O_{\delta}(p, A)$  contained in  $f^{-1}(U)$ . It follows that  $f(O_{\delta}(p, A)) \subset U$ . By Exercise 1.2.5 the open ball  $O_{\delta}(p, A)$  is an open subset of A, so let  $V = O_{\delta}(p, A)$ .

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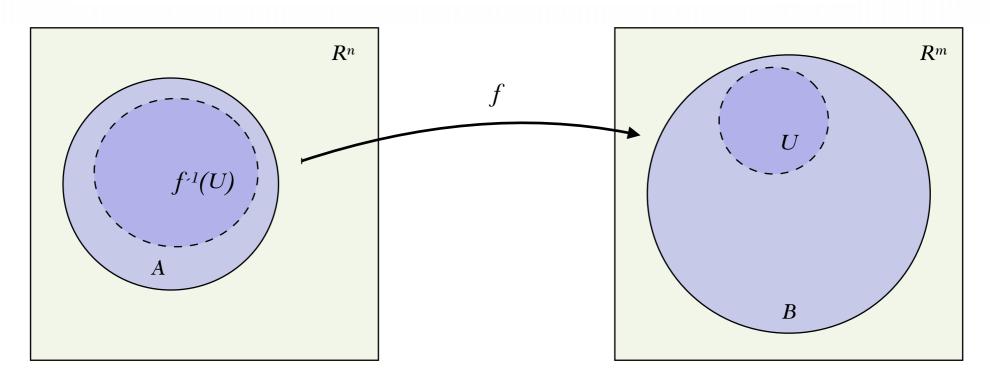
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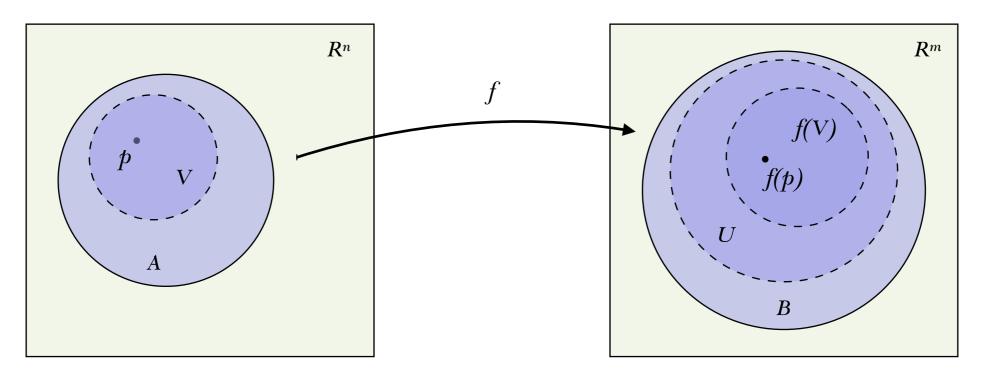
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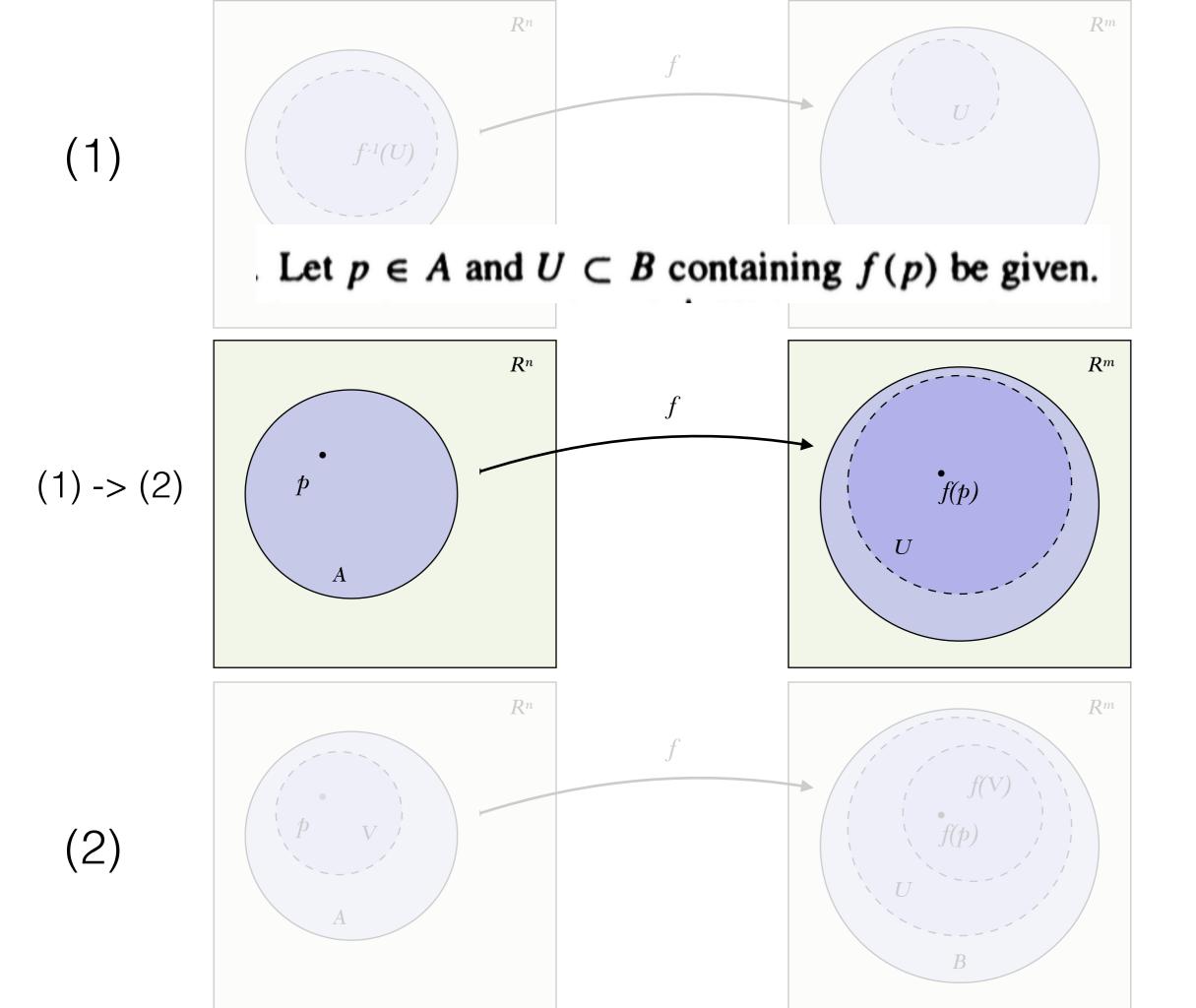
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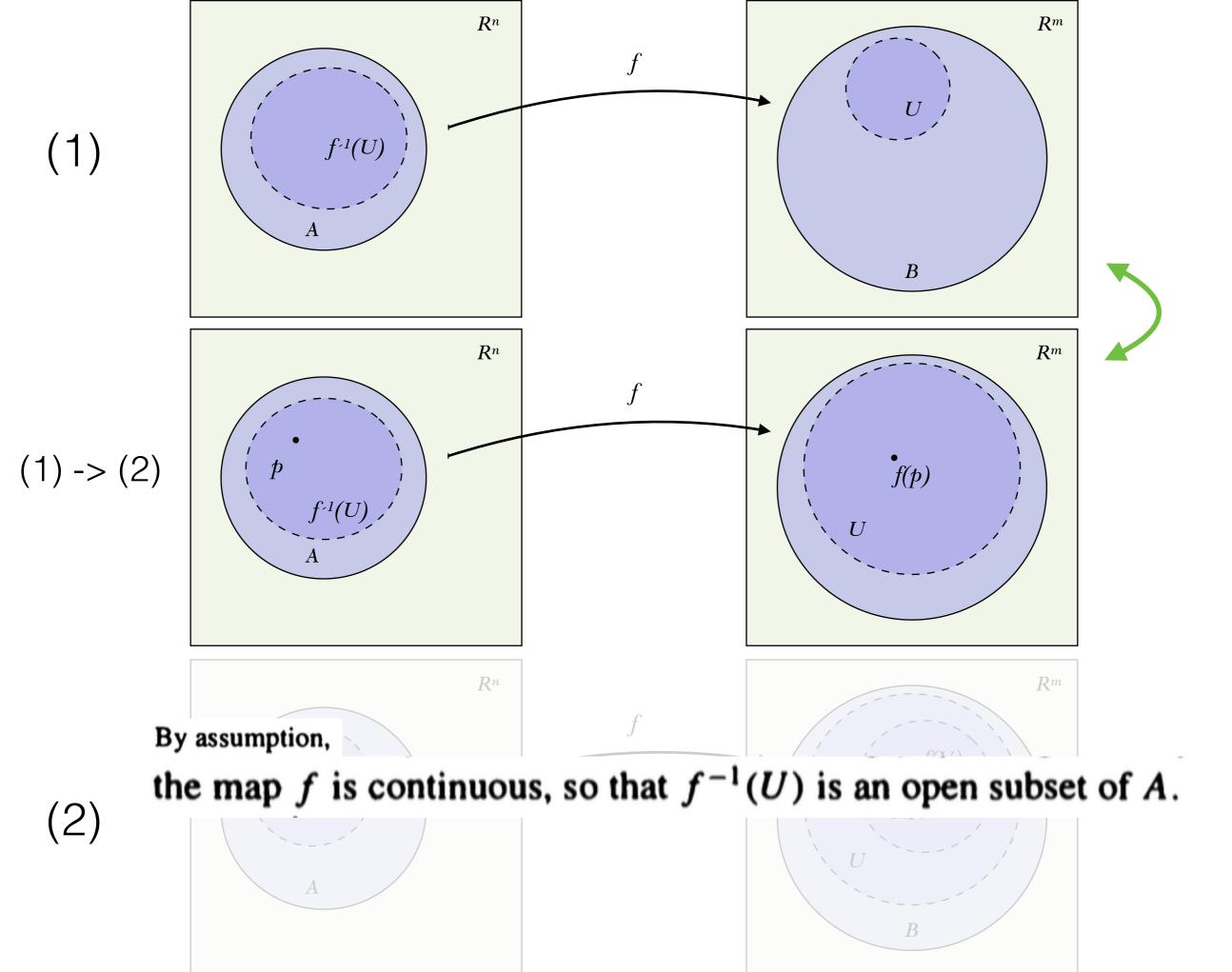
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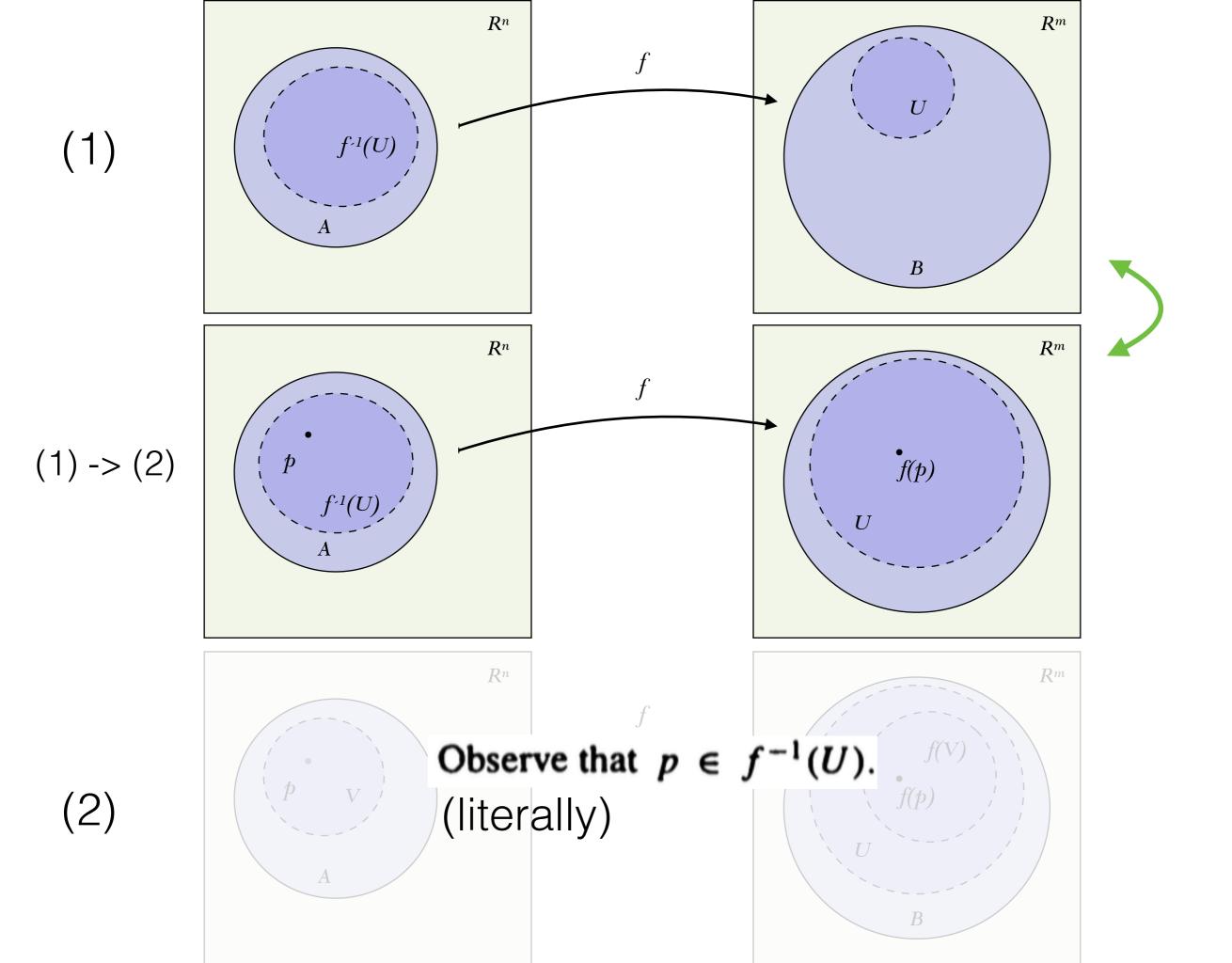


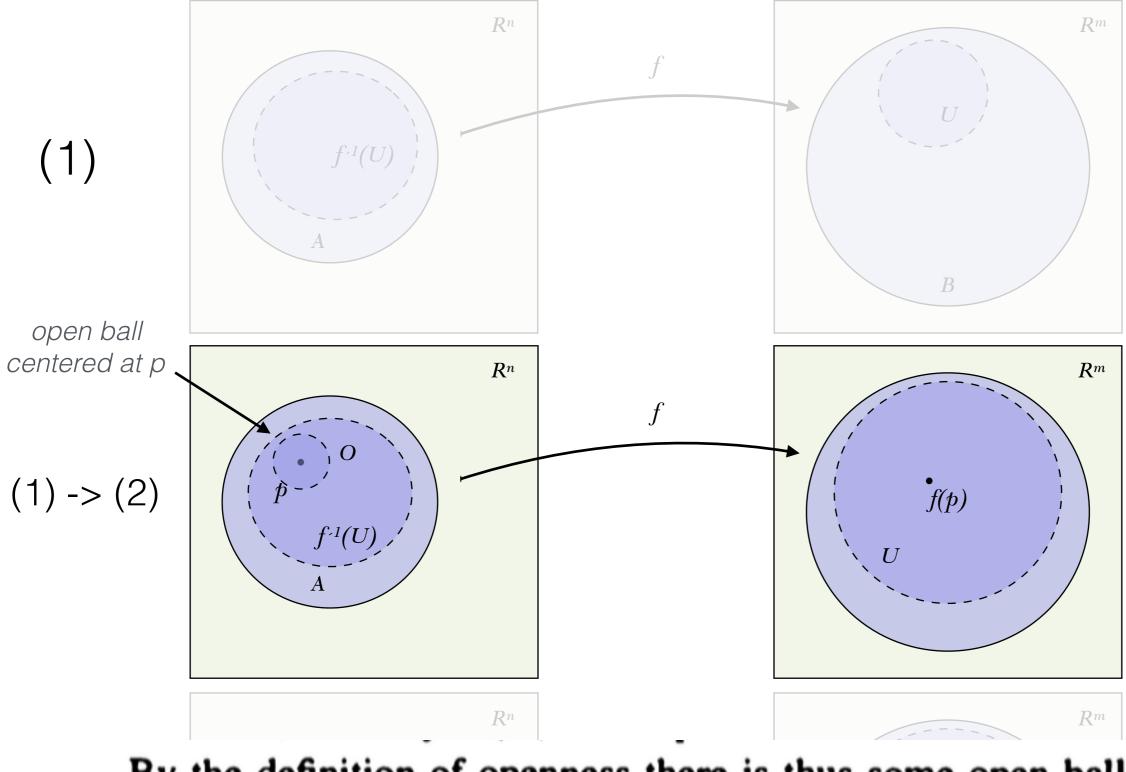
(2) For every point  $p \in A$ , and every open subset  $U \subset B$  containing f(p), there is an open subset  $V \subset A$  containing p such that  $f(V) \subset U$ .









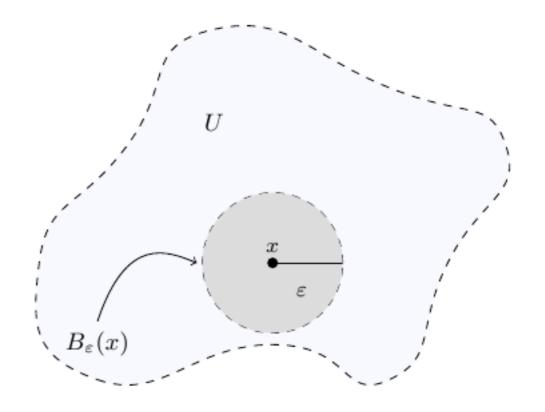


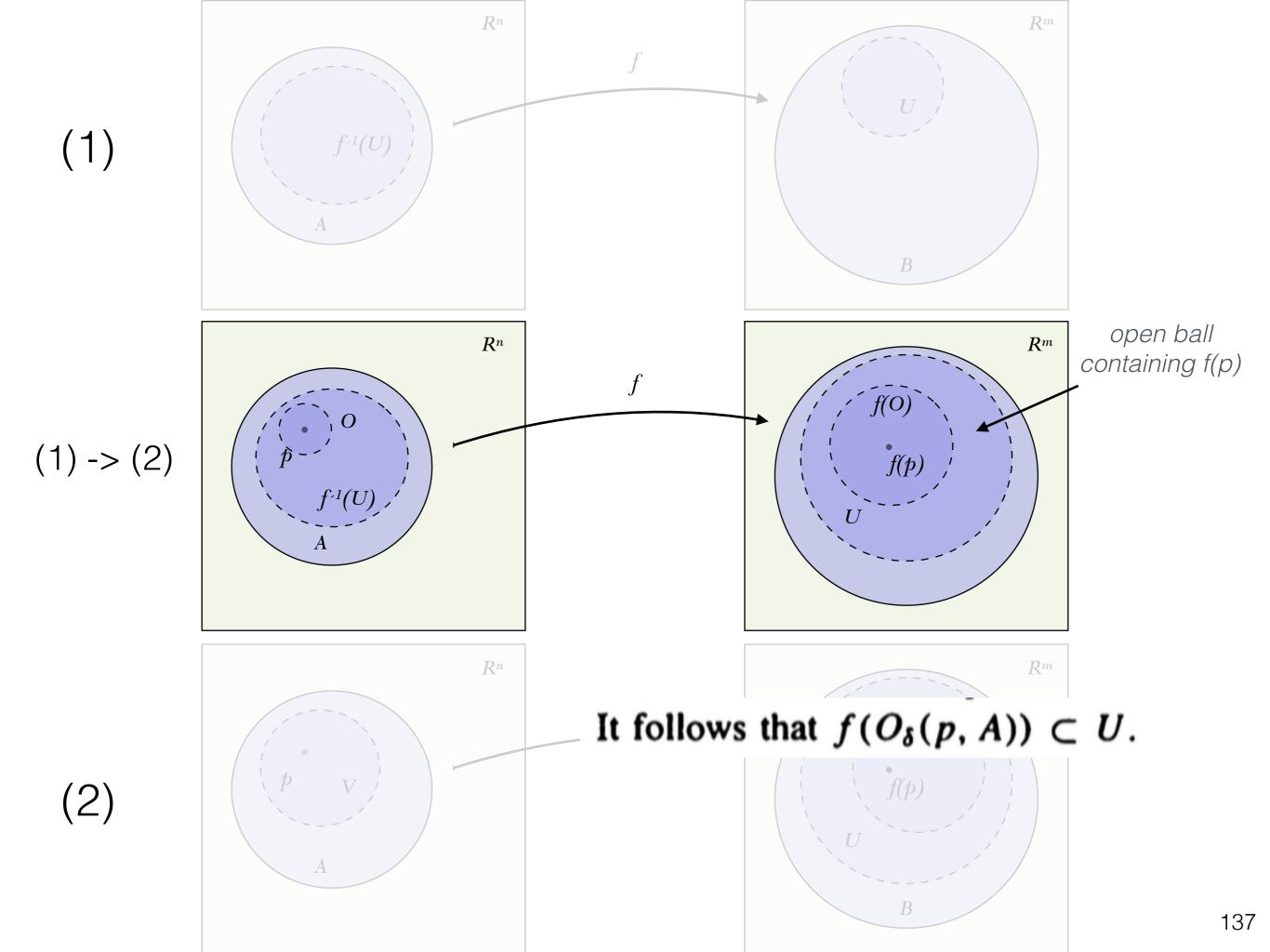
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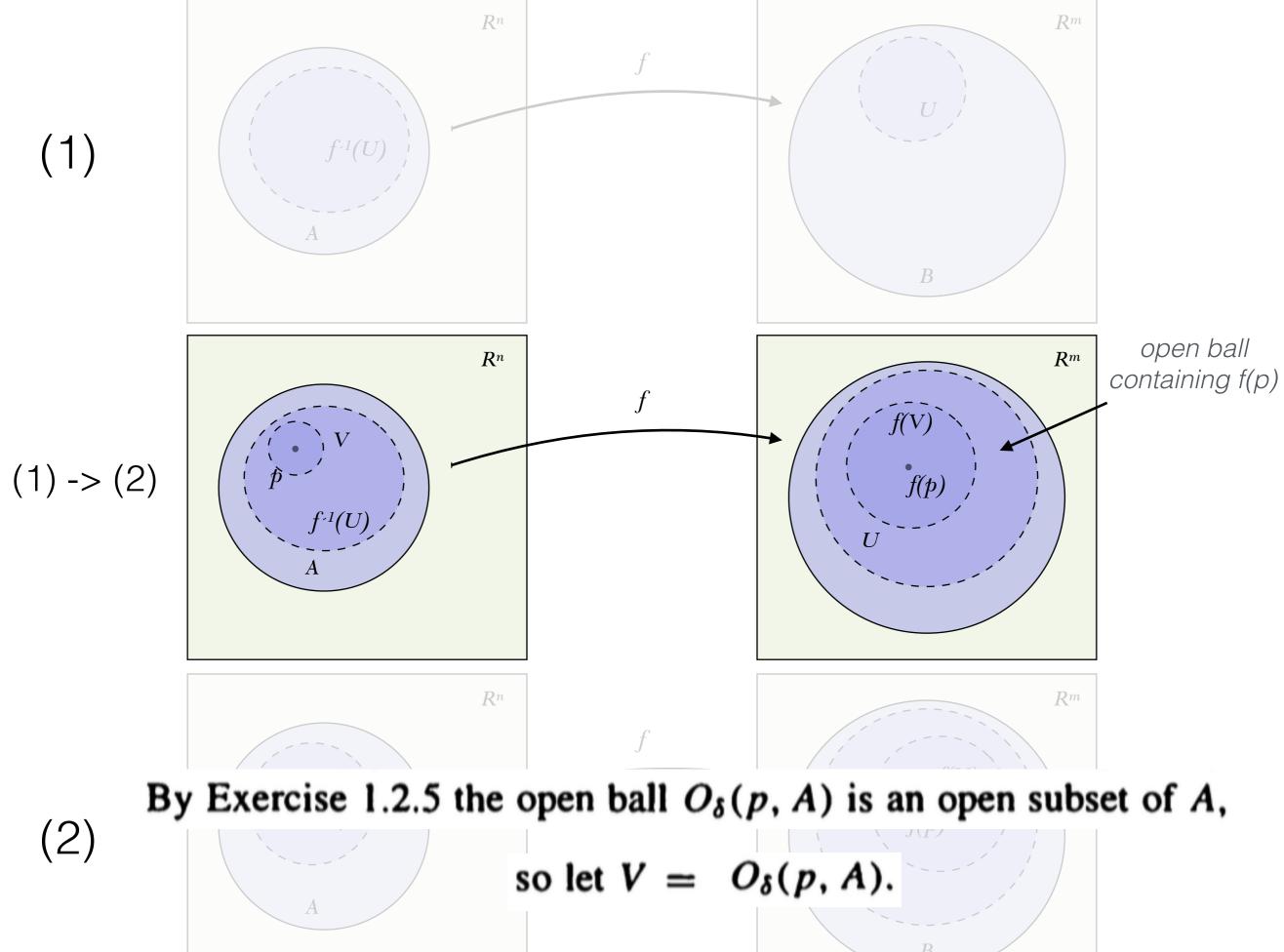
(2)

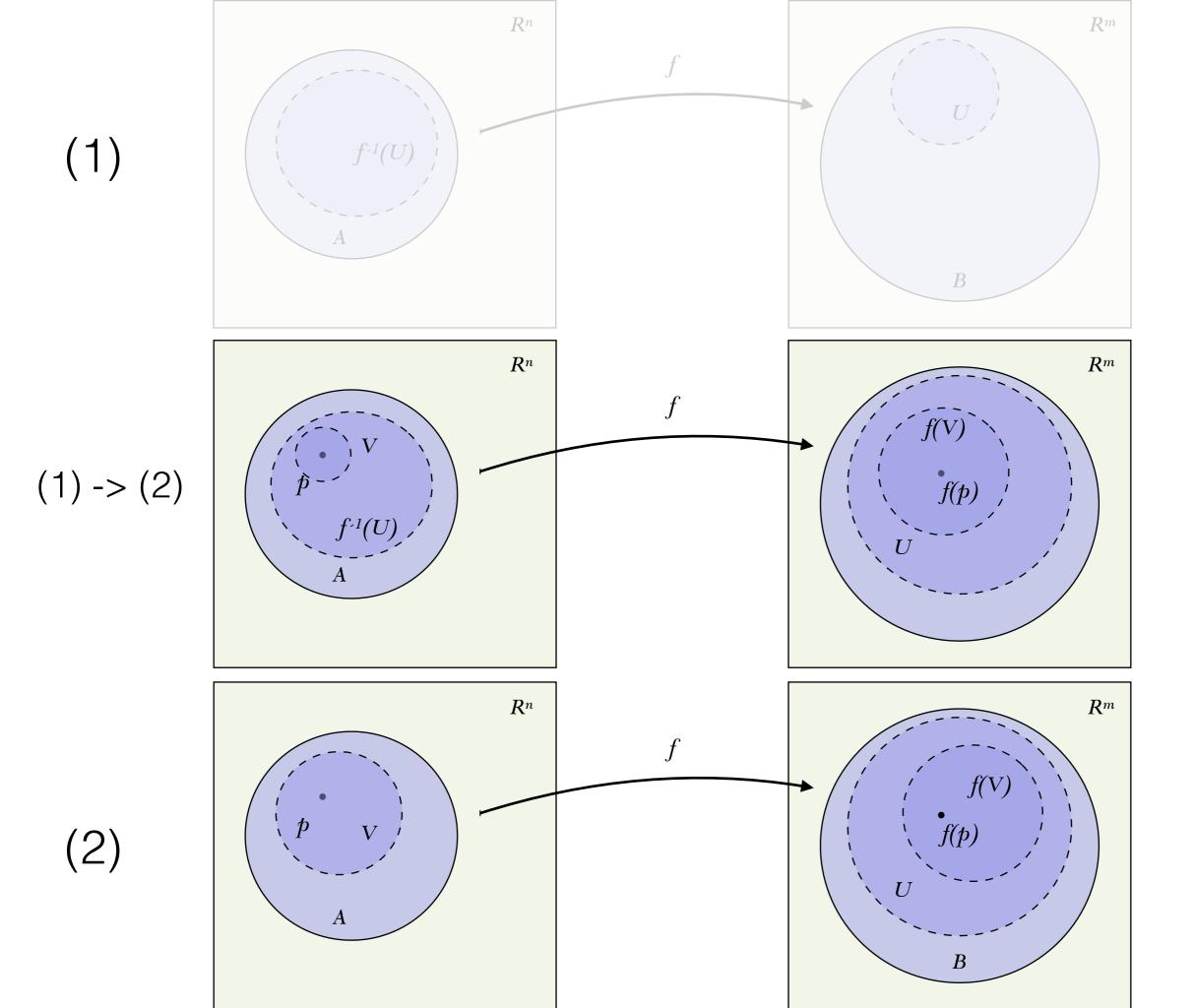
### recall: an open set in R<sup>2</sup>

OpenSet U





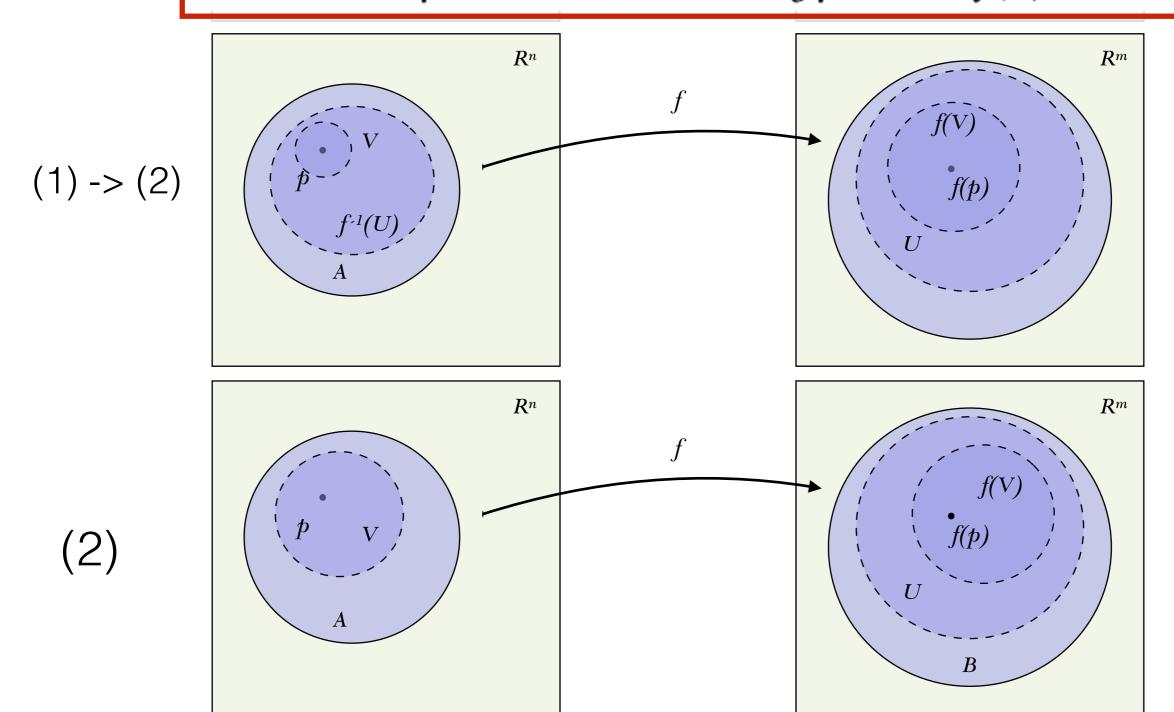


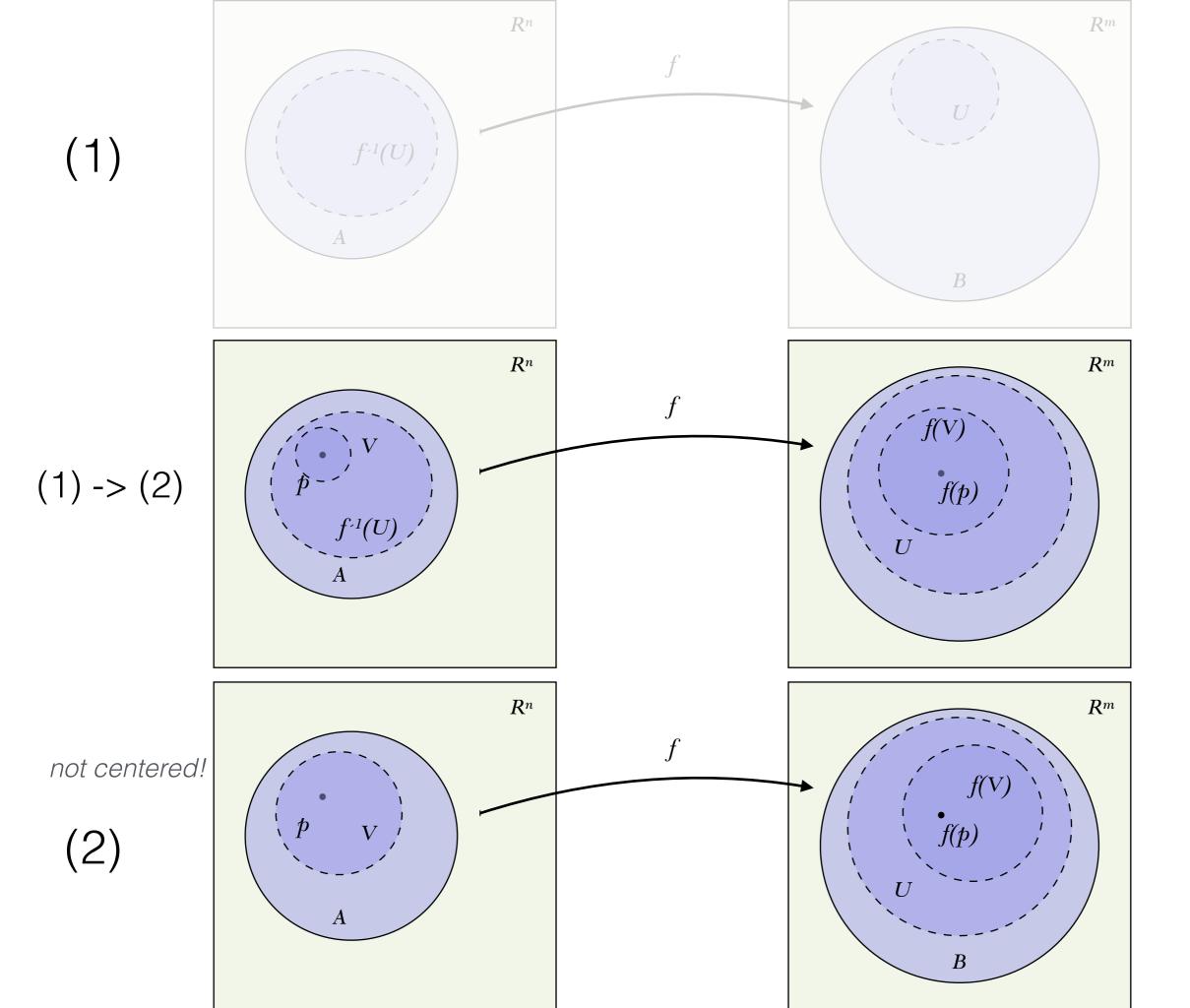


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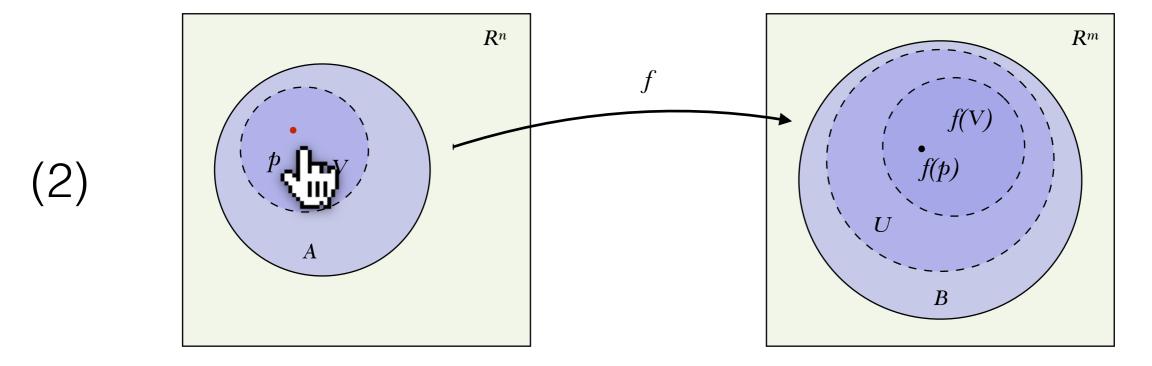
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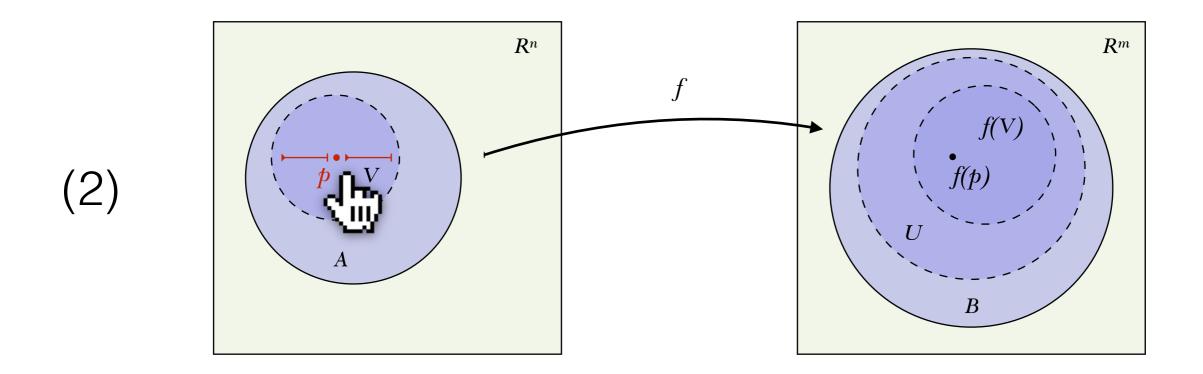


### Prodirect manipulation: center the points

want to center this point



## Prodirect manipulation: center the points



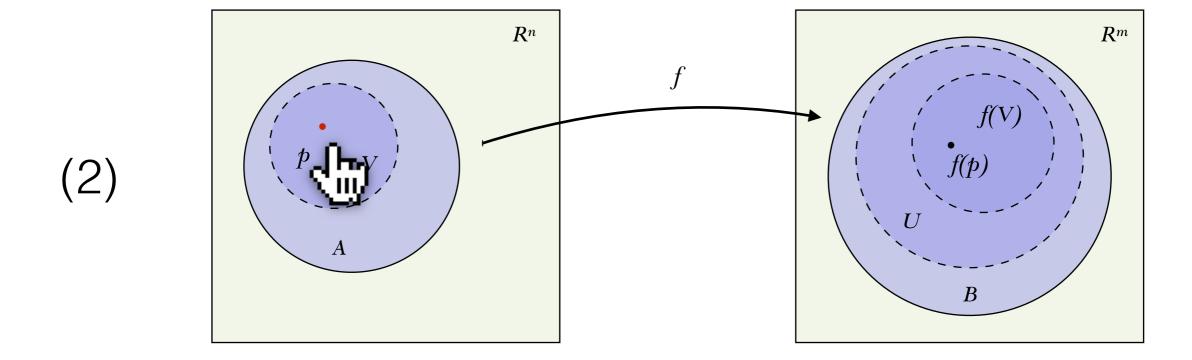
### Substance

Set A
Set B
Set R^n
Set R^m
Subset A R^n
Subset B R^m
Map f A B
OpenSet U
Subset U B

Point p
In p A
Point f(p)
In f(p) U
OpenSet V
Subset V A
In p V
Set f(V)
Subset f(V) U
In f(p) f(V)

### Style

Style All Auto
Shape R^n Square
Shape R^m Square
Color R^n Yellow
Color R^m Yellow



#### Substance

Set A
Set B
Set R^n
Set R^m
Subset A R^n
Subset B R^m
Map f A B
OpenSet U
Subset U B

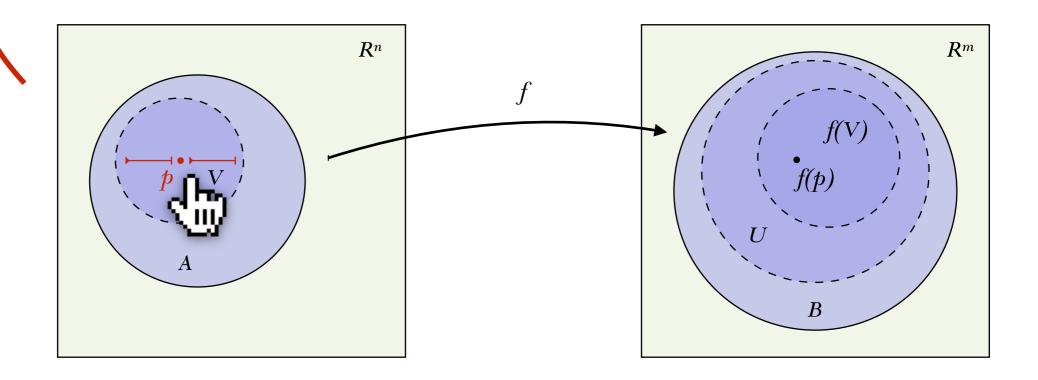
Point p
In p A
Point f(p)
In f(p) U
OpenSet V p
Subset V A
In p V
Set f(V)
Subset f(V) U
In f(p) f(V)

### Style

Style All Auto
Shape R^n Square
Shape R^m Square
Color R^n Yellow
Color R^m Yellow

runtime infers code changes!

(2)



#### Substance

Set A
Set B
Set R^n
Set R^m
Subset A R^n
Subset B R^m
Map f A B
OpenSet U
Subset U B

Point p
In p A
Point f(p) NOW (
In f(p) U
OpenSet V p
Subset V A
In p V
Set f(V)
Subset f(V) U
In f(p) f(V)

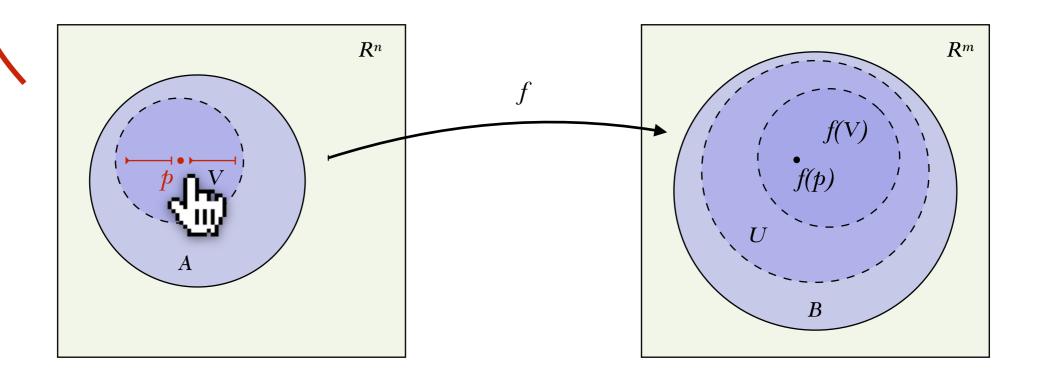
#### Style

Point p
In p A
Point f(p) now centered at p m
Square
In f(p) U
Color R^n Yellow
OpenSet V p

Style All Auto
Shape R^n Square
Color R^n Yellow

runtime infers code changes!

(2)



Easily implement and extend DSLs

- Easily implement and extend DSLs
- Easily reuse styles

- Easily implement and extend DSLs
- Easily reuse styles

Author's Guide to the ACM SIGPLAN Class

(sigplanconf.cls)

v3.6

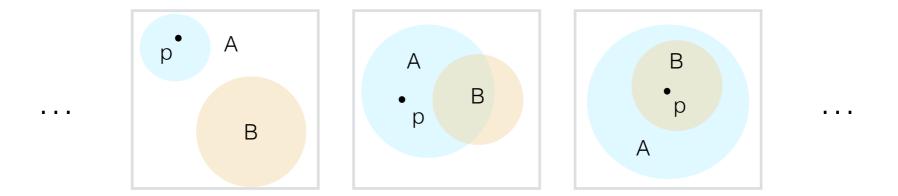
This guide describes Version 3.6 of the class file, which ends with a complete revision history that starts with the file's creation in September 2004.

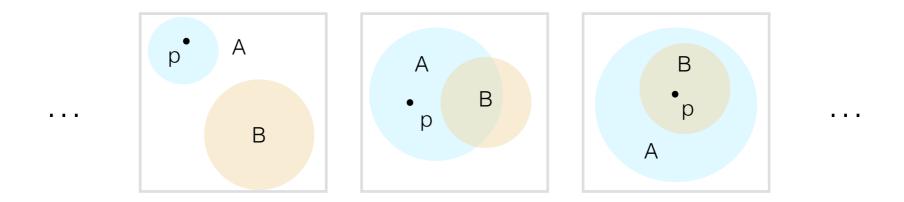
- Easily implement and extend DSLs
- Easily reuse styles
- Expertise encoded! Substantial benefits accrue with each new diagram

- Easily implement and extend DSLs
- Easily reuse styles
- Expertise encoded! Substantial benefits accrue with each new diagram
- DSL users can just throw in notation; get a useful illustration

# Language design challenges

Penrose provides an extensible visual semantics for mathematical notation.





# Where are the semantics?

### Notation

Set A Set B Point p  $p \in A$ 

Set A
Set B
Point p
In p A

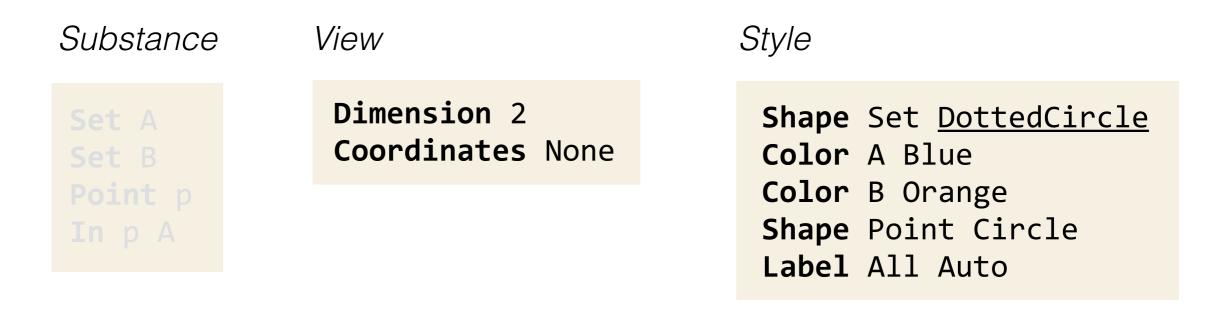
Style

Style

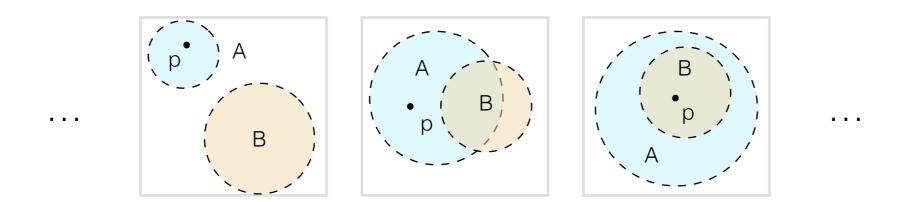
Shape Set Circle
Color A Blue
Color B Orange
Shape Point Circle
Label All Auto

Substance View Style Dimension 2 **Shape** Set Circle Coordinates None Color A Blue Color B Orange **Shape** Point Circle Label All Auto

В



### incorrect semantics?

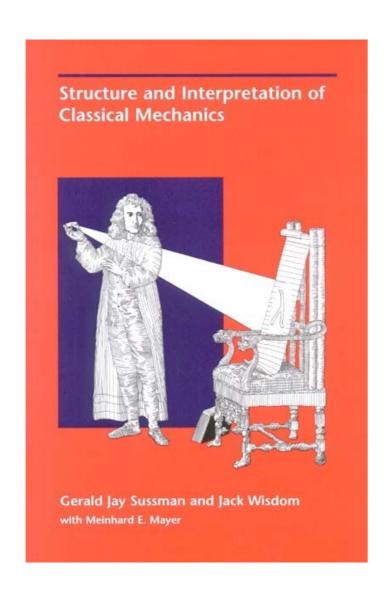


The derivative of the coordinate path Dq is the function that maps time to velocity components:

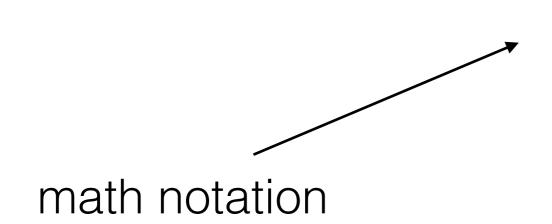
$$Dq(t) = (Dx(t), Dy(t), Dz(t)).$$

We can make and use the derivative of a function. $\frac{31}{2}$  For example, we can write:

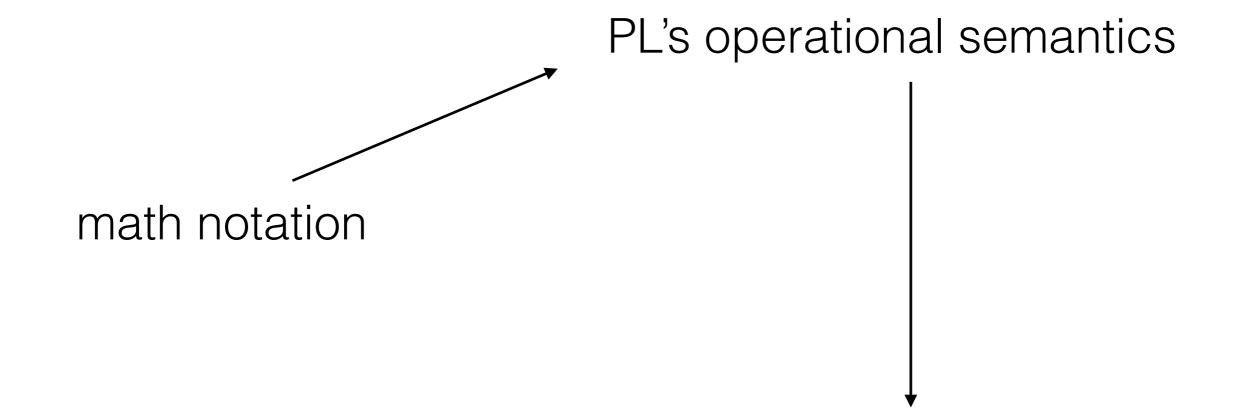
```
(print-expression ((D q) 't))
(up ((D x) t) ((D y) t) ((D z) t))
```



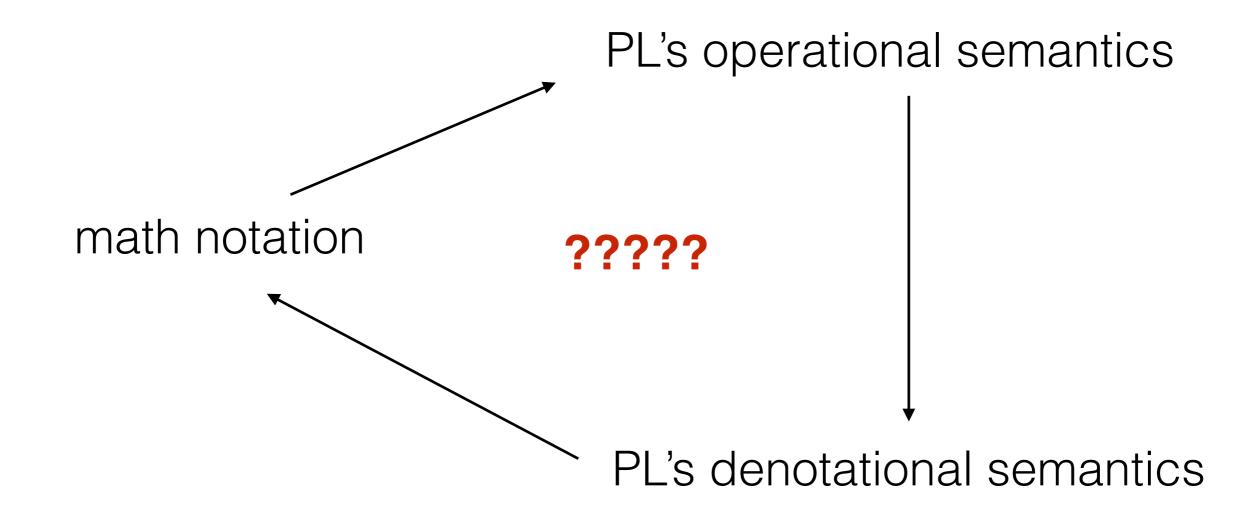
Once formalized as a procedure, a mathematical idea becomes a tool that can be used directly to compute results.



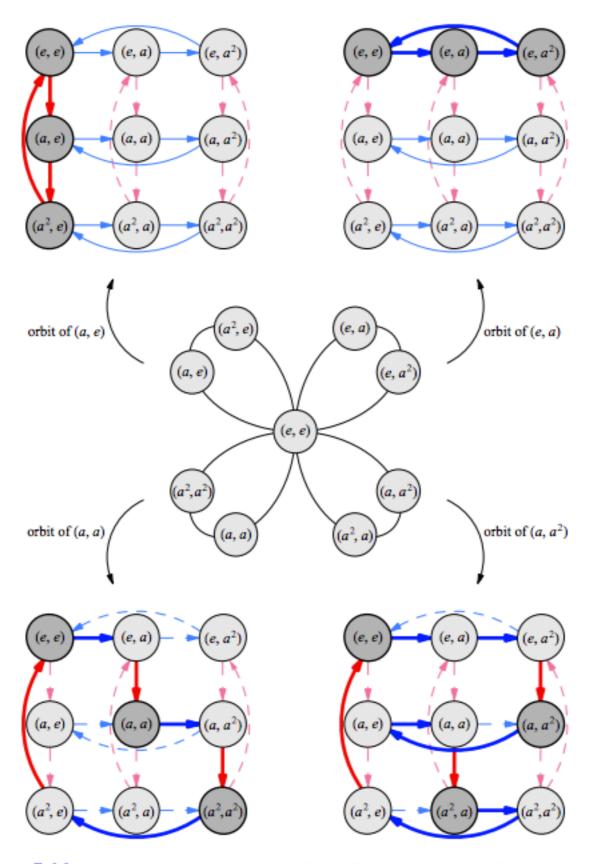
PL's operational semantics



PL's denotational semantics

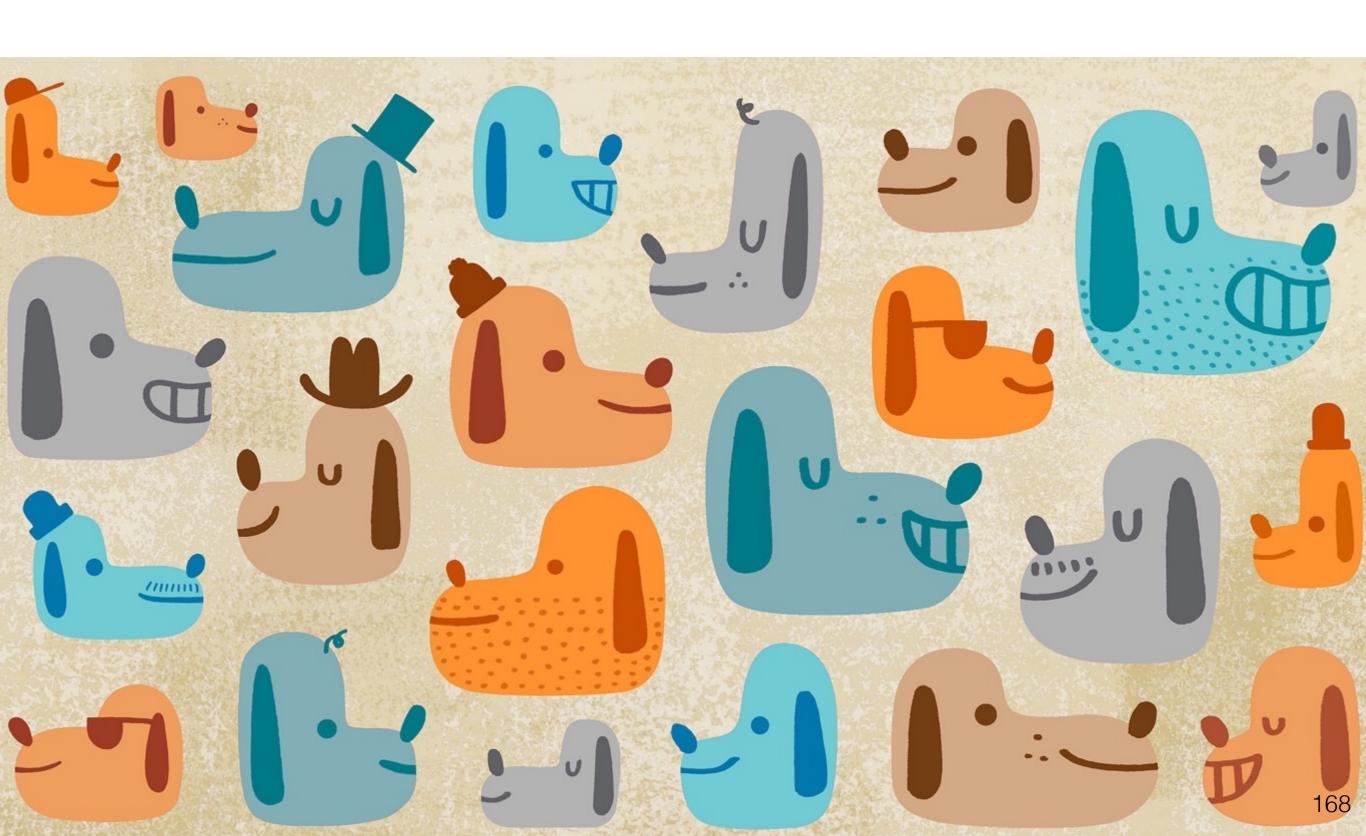


# Modularity, compositionality

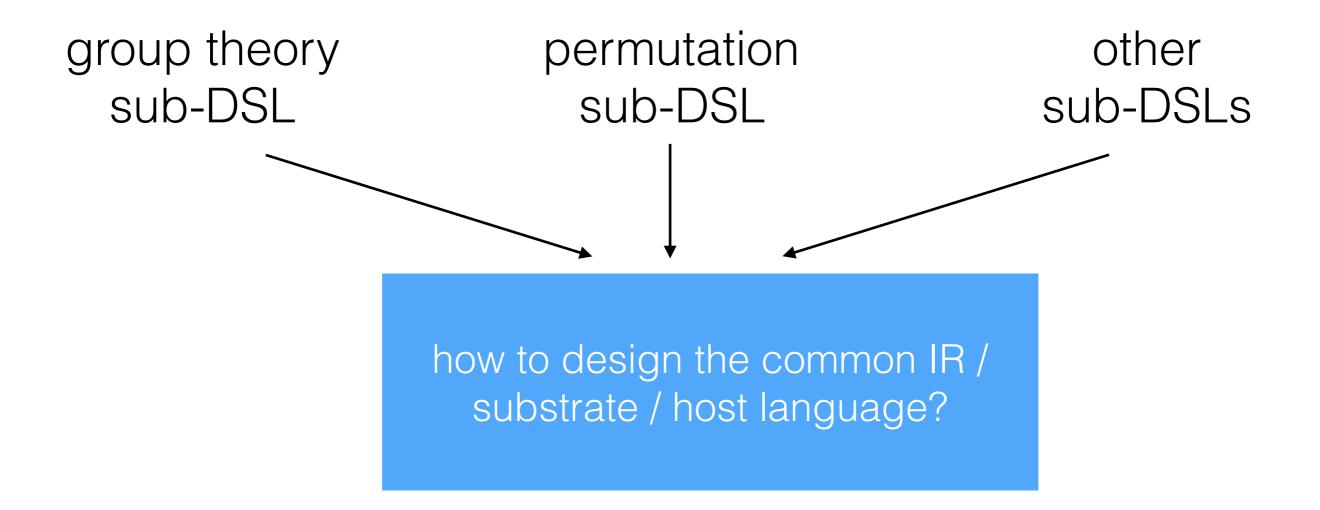


**Figure 5.14.** In the center is the cycle graph of  $C_3 \times C_3$ , and next to each of the four orbits is a copy of the Cayley diagram for  $C_3 \times C_3$  with the corresponding orbit highlighted. These Cayley diagrams are structured like the original in Figure 2.10 on page 22.

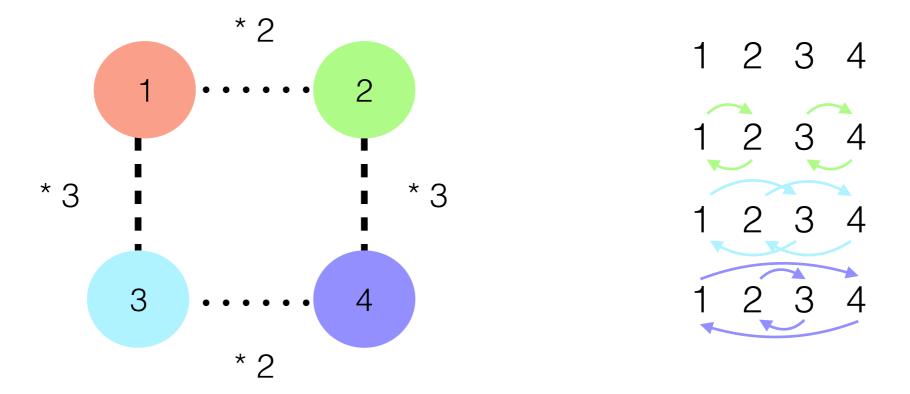
# Dynamic Symbols in Illustrator



# Interoperability



# Cayley's theorem

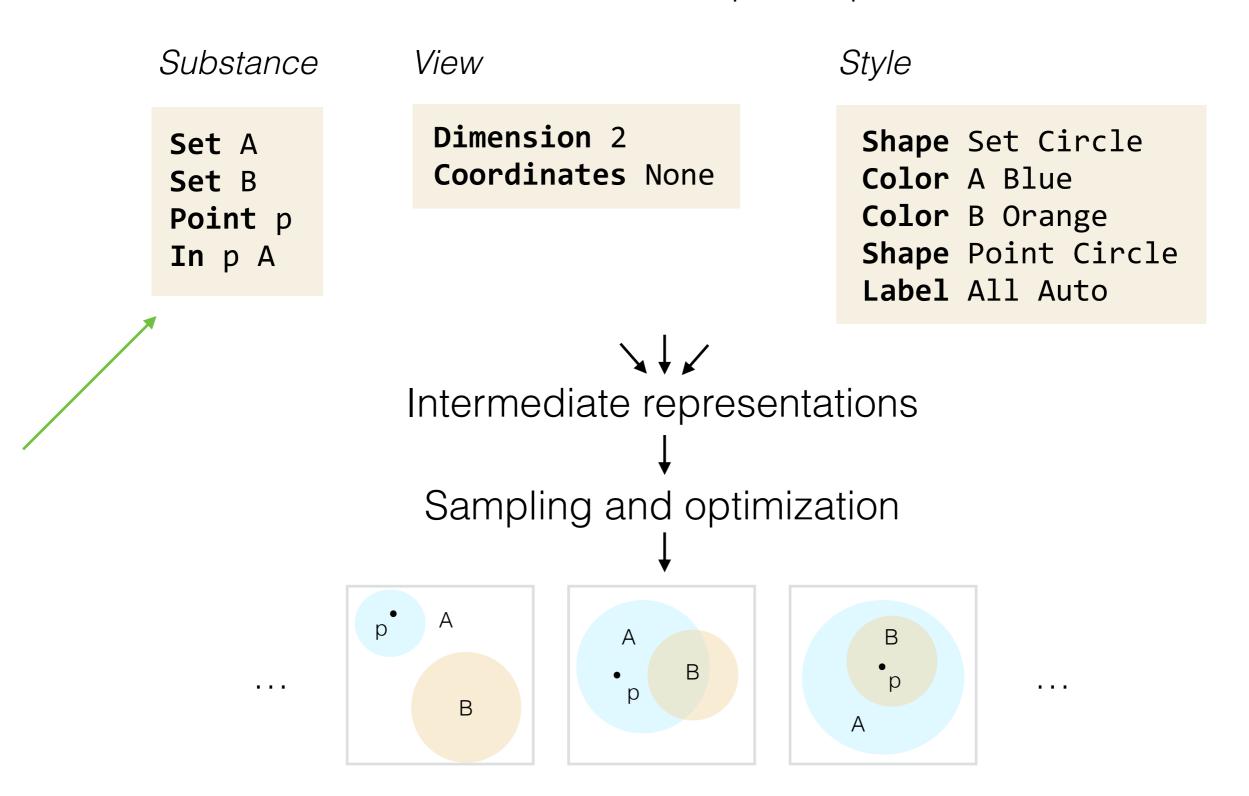


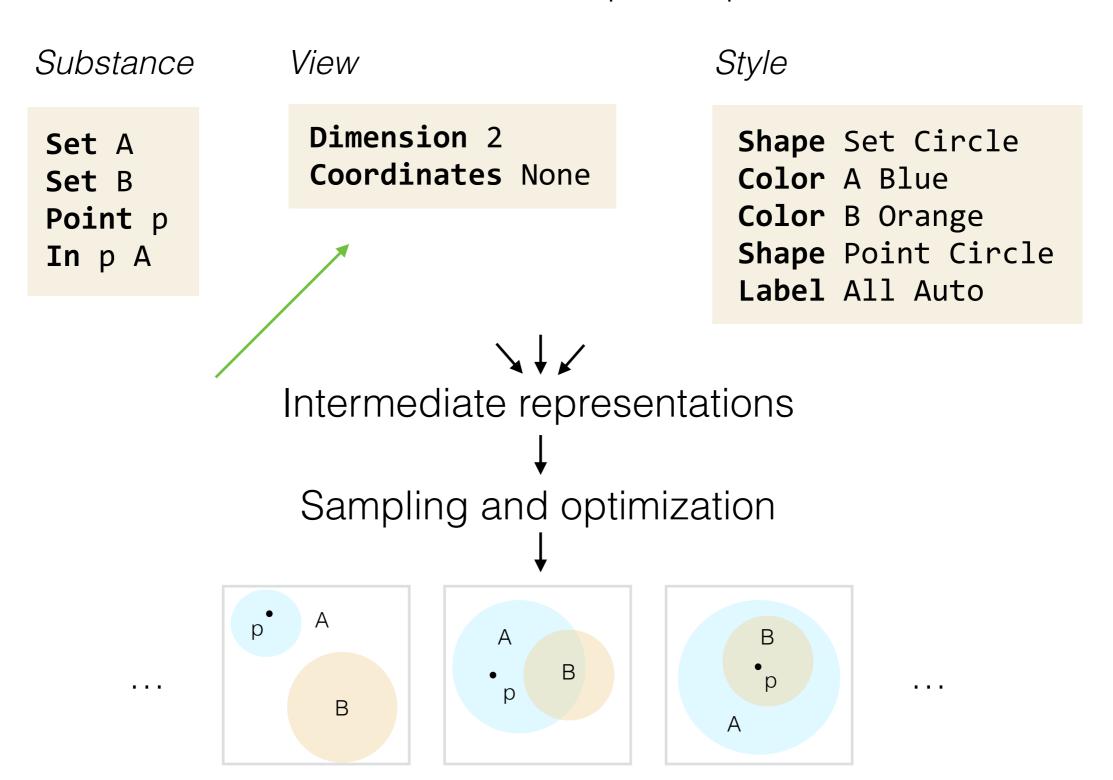
Cayley diagram for K<sub>4</sub>

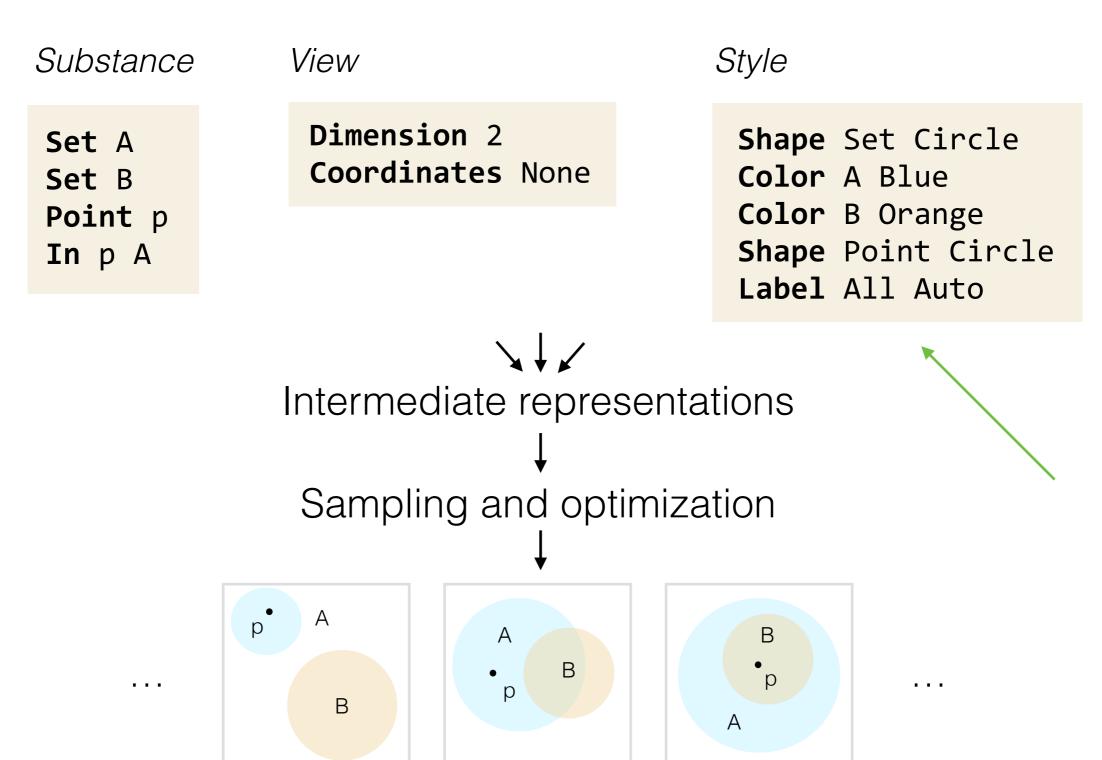
Permutation group

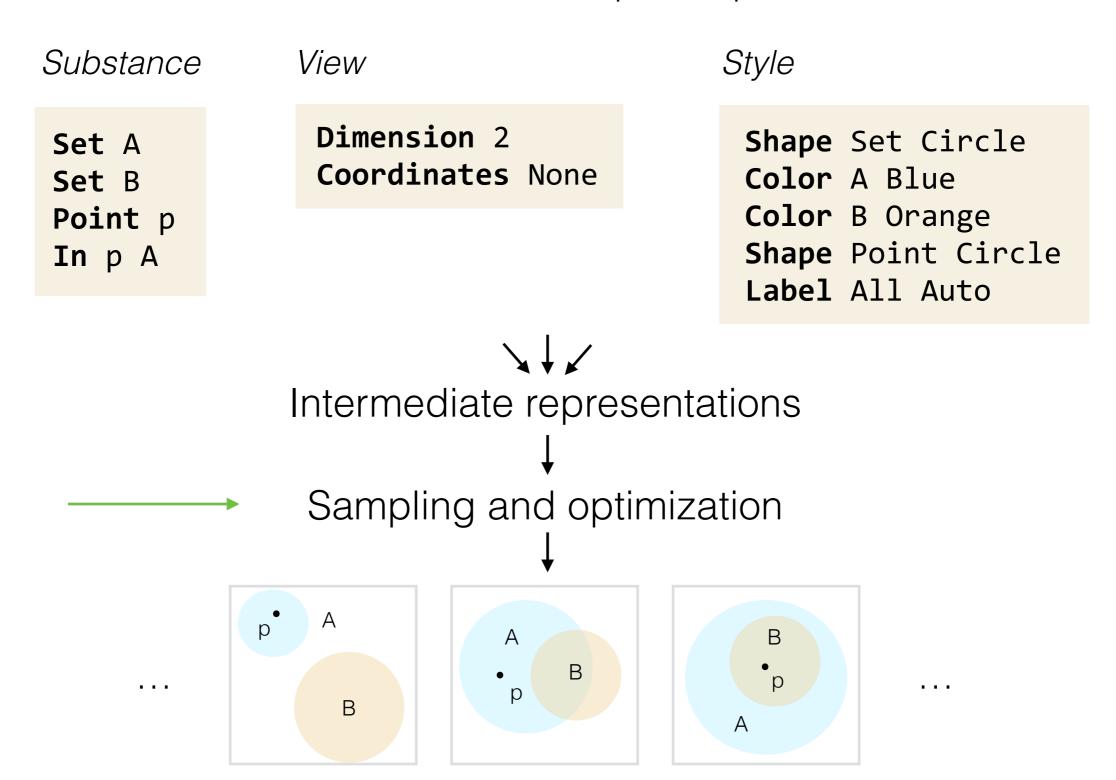
# Extensibility

### ...for each component of Penrose

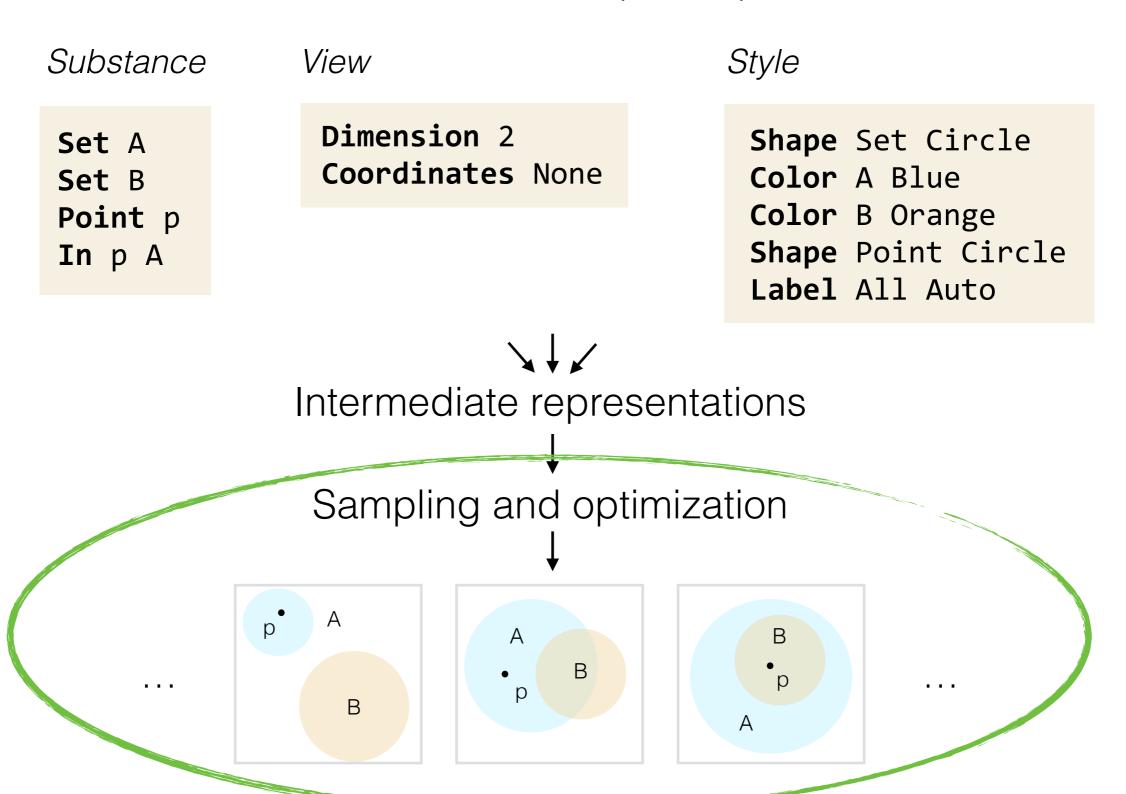


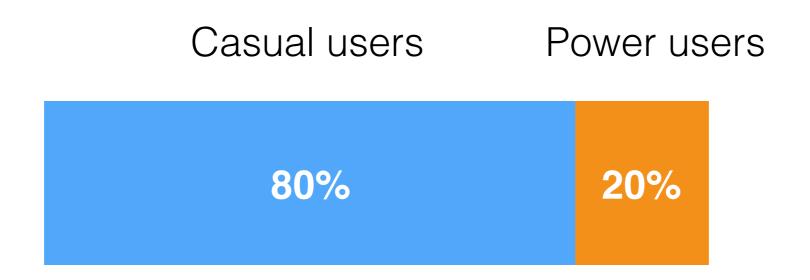






# Layout and runtime challenges



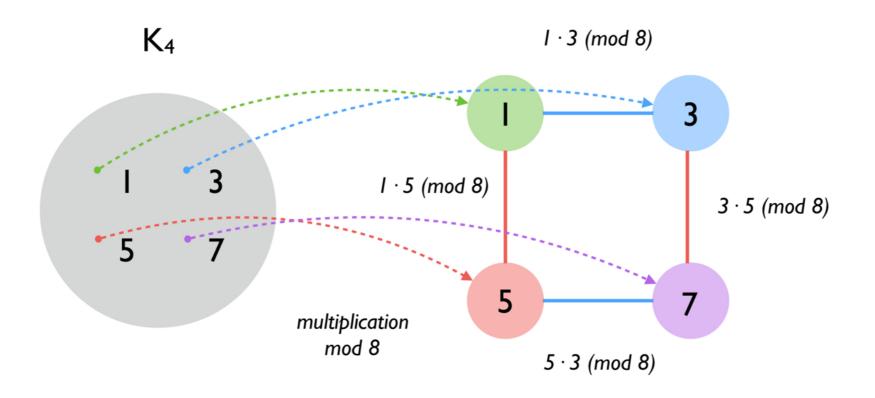


### Constraints

```
p in A
A and B do not intersect
curve 2 is not homotopic to 0
B is contained in the preimage of f
```

### Style

linked.layout = left\_to\_right



## Style

Label All Auto
Position All Auto



### Many common decisions for geometry

What's the geometry?

Where does the camera go, and what's the focal length?

Where's the light/shadow?

Do we want to include shading?

Do we want a separate light for the shadow vs. the shading?

Which "geometric" lines should be included?

Always the silhouette.

What about other contours?

Which "combinatorial" lines should be included?

E.g., grid lines.

Should we automatically trace contours, or do it by hand?

Are there any components that can be algorithmically generated?

What's the line thickness, for various classes of lines?

What's the line style for hidden lines?

What's the line style for partial lines?

What's the fill color for various regions?

Should we use different colors/brightnesses to indicate lighting?

Should we use gradients for shading?

How should we color occluded regions?

How dark is the shadow?

What direction does the shadow gradient go?

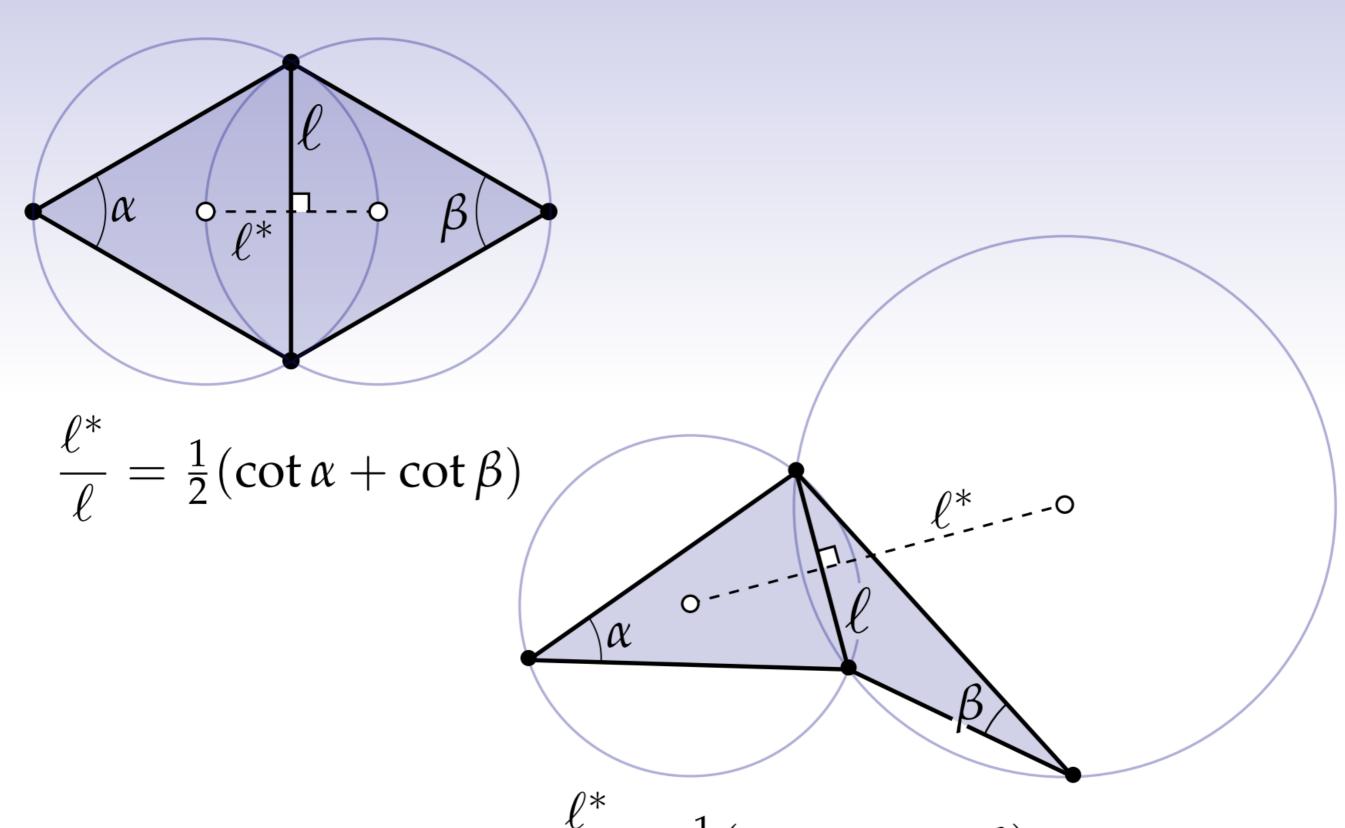
How do we label the figure?

What gets labeled?

Where do labels go? (How) should they occlude/be occluded by geometry?

Slide borrowed from Keenan Crane

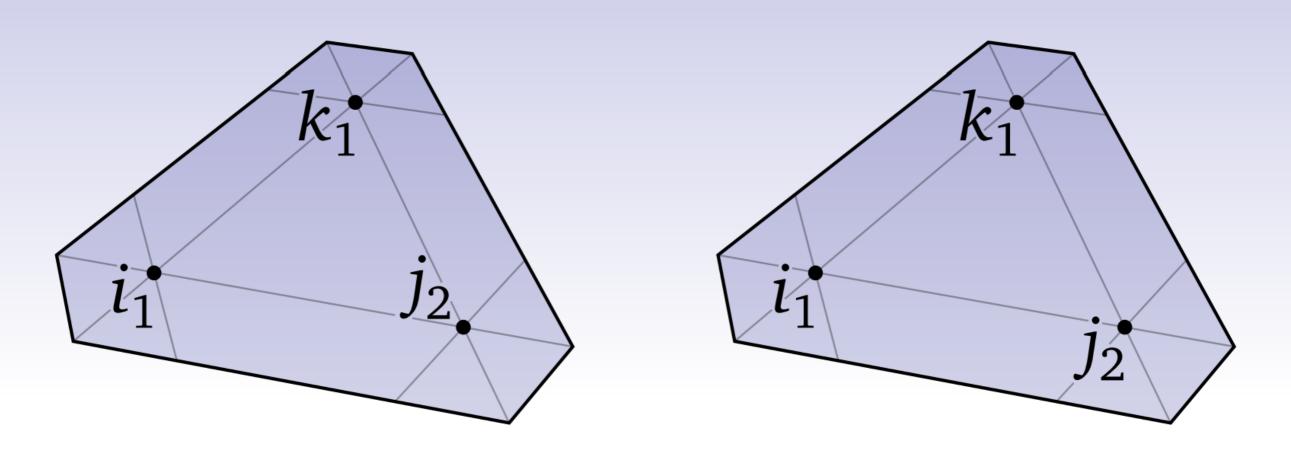
# Is This Figure Misleading?



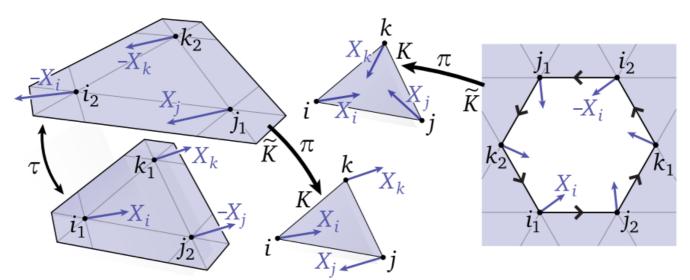
Slide borrowed from Keenan Crane

$$\frac{\ell^*}{\ell} = \frac{1}{2}(\cot\alpha + \cot\beta)$$

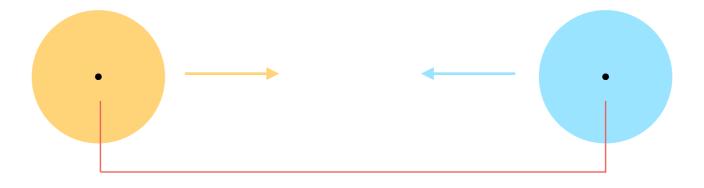
# Labels - Global Placement



### ...In context, lots of constraints!

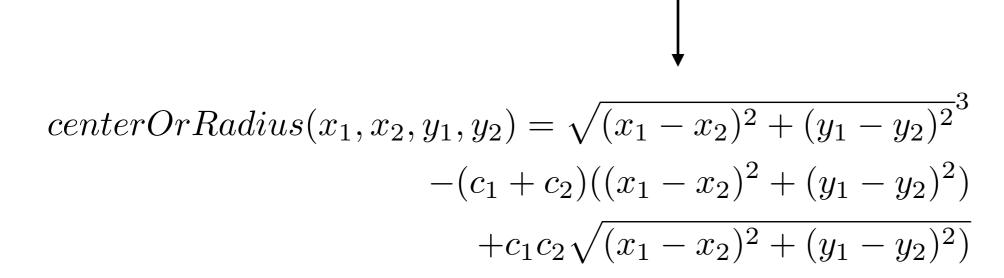


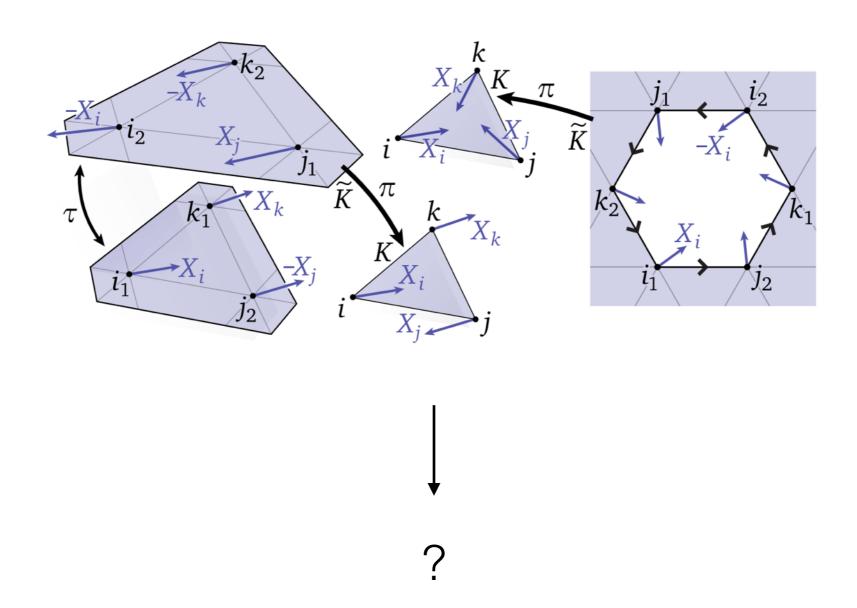
Slide borrowed from Keenan Crane



lay it out left to right center sets, but keep them away from each other position labels close to their object but far from others' smart line breaks

lay it out left to right center sets, but keep them away from each other position labels close to their object but far from others' smart line breaks

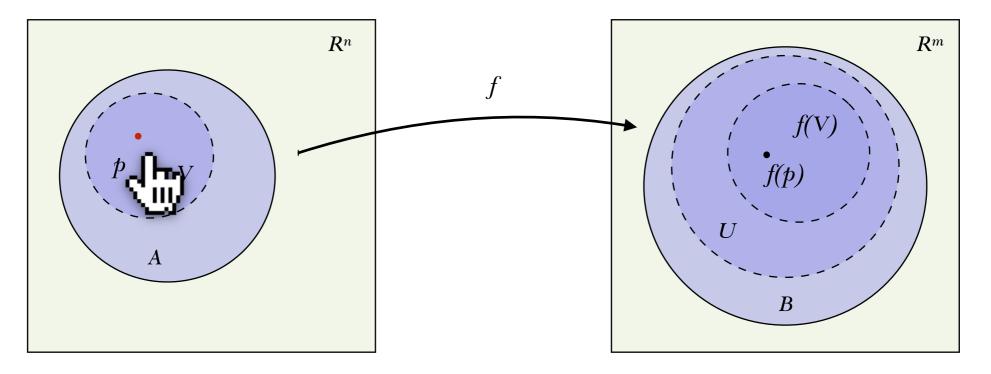


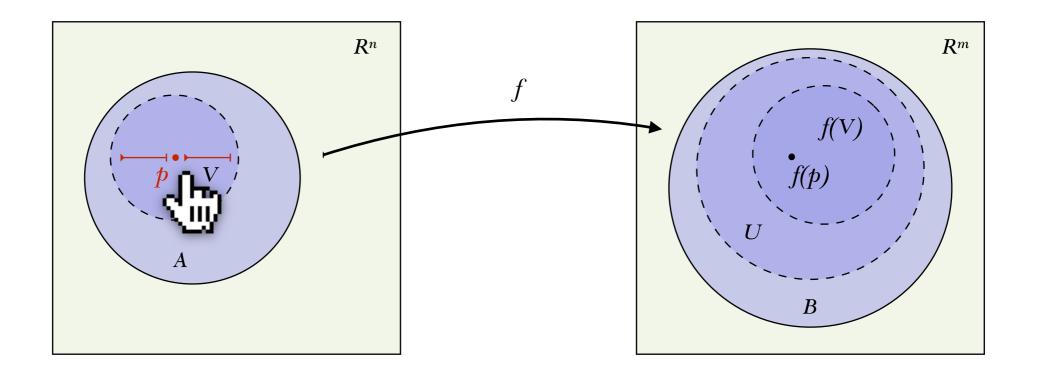


Source: Keenan Crane

### Prodirect manipulation: center the points

want to center this point





### Substance

Set A
Set B
Set R^n
Set R^m
Subset A R^n
Subset B R^m
Map f A B
OpenSet U
Subset U B

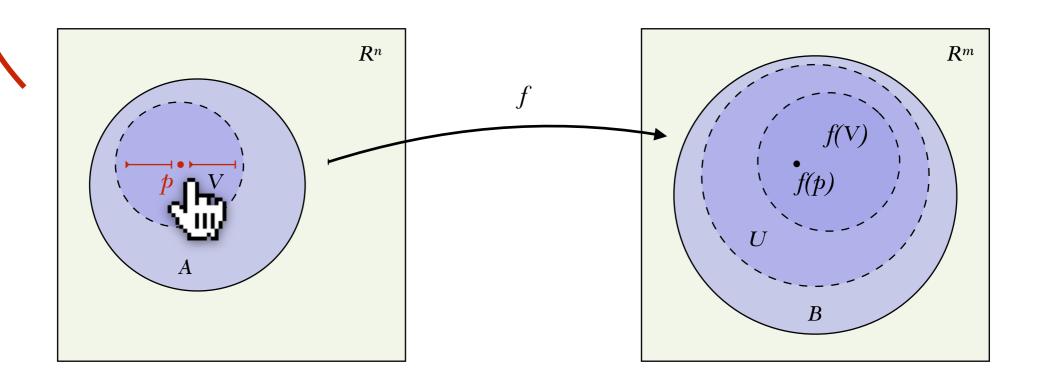
Point p
In p A
Point f(p) now o
In f(p) U
OpenSet V p
Subset V A
In p V
Set f(V)
Subset f(V) U
In f(p) f(V)

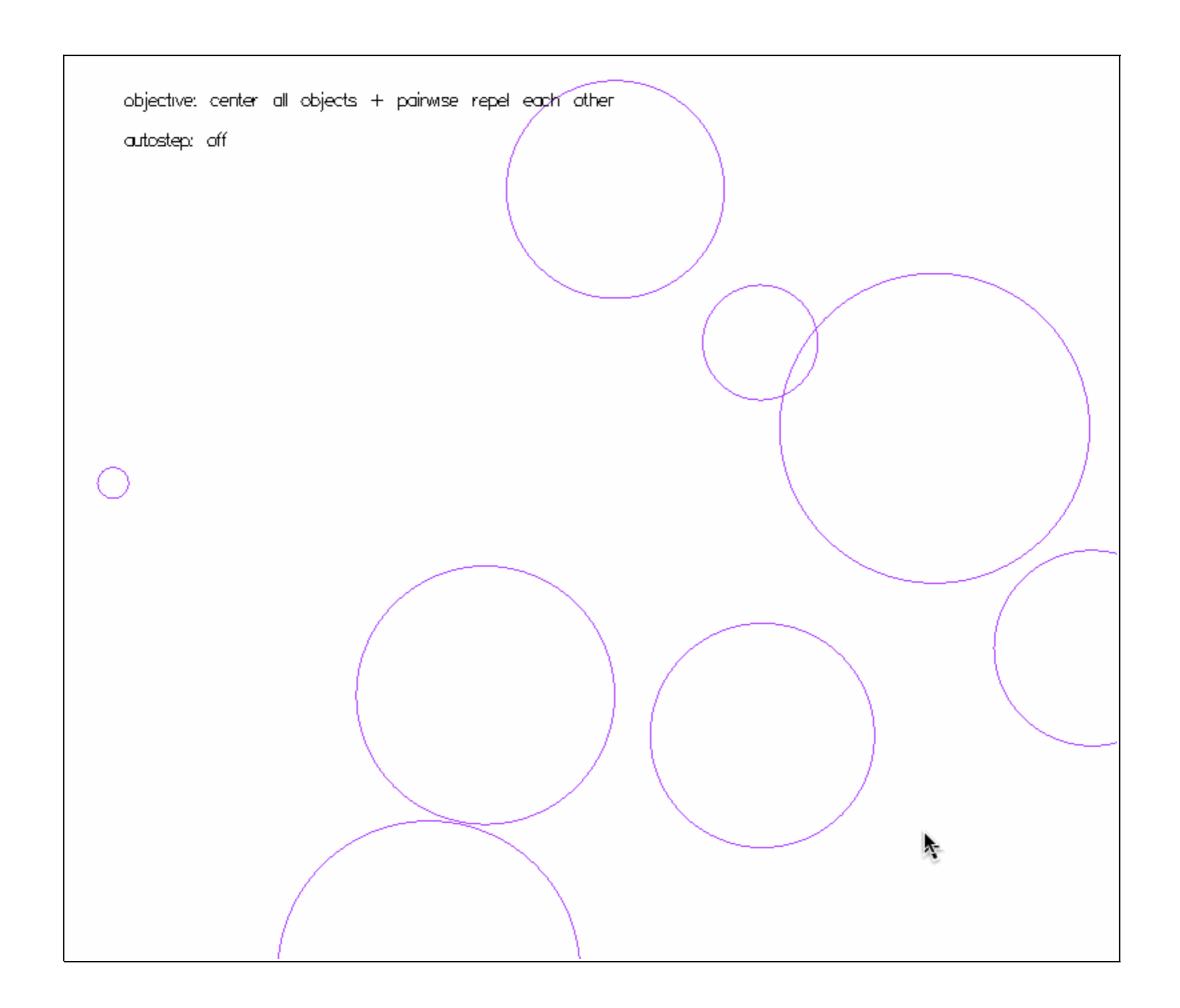
### Style

Point p
In p A
Point f(p) now centered at p m
Square
In f(p) U
Color R^n Yellow
OpenSet V p

Style All Auto
Shape R^n Square
Color R^n Yellow

runtime infers code changes!





# import Expertise

# Our broader vision

# Automatically parse and visualize mathematics

#### MEAN VALUE THEOREMS

5.7 Definition Let f be a real function defined on a metric space X. We say that f has a local maximum at a point  $p \in X$  if there exists  $\delta > 0$  such that  $f(q) \le f(p)$  for all  $q \in X$  with  $d(p, q) < \delta$ .

Local minima are defined likewise.

Our next theorem is the basis of many applications of differentiation.

**5.8 Theorem** Let f be defined on [a, b]; if f has a local maximum at a point  $x \in (a, b)$ , and if f'(x) exists, then f'(x) = 0.

The analogous statement for local minima is of course also true.

**Proof** Choose  $\delta$  in accordance with Definition 5.7, so that

$$a < x - \delta < x < x + \delta < b$$
.

If  $x - \delta < t < x$ , then

$$\frac{f(t) - f(x)}{t - x} \ge 0.$$

Letting  $t \to x$ , we see that  $f'(x) \ge 0$ .

If  $x < t < x + \delta$ , then

$$\frac{f(t) - f(x)}{t - x} \le 0,$$

which shows that  $f'(x) \le 0$ . Hence f'(x) = 0.

5.9 Theorem If f and g are continuous real functions on [a,b] which are differentiable in (a,b), then there is a point  $x \in (a,b)$  at which

$$[f(b) - f(a)]g'(x) = [g(b) - g(a)]f'(x).$$

Note that differentiability is not required at the endpoints.

Proof Put

$$h(t) = [f(b) - f(a)]g(t) - [g(b) - g(a)]f(t) \qquad (a \le t \le b).$$

Then h is continuous on [a, b], h is differentiable in (a, b), and

(12) 
$$h(a) = f(b)g(a) - f(a)g(b) = h(b).$$

To prove the theorem, we have to show that h'(x) = 0 for some  $x \in (a, b)$ . If h is constant, this holds for every  $x \in (a, b)$ . If h(t) > h(a) for some  $t \in (a, b)$ , let x be a point on [a, b] at which h attains its maximum (Theorem 4.16). By (12),  $x \in (a, b)$ , and Theorem 5.8 shows that h'(x) = 0. If h(t) < h(a) for some  $t \in (a, b)$ , the same argument applies if we choose for x a point on [a, b] where h attains its minimum.

This theorem is often called a generalized mean value theorem; the following special case is usually referred to as "the" mean value theorem:

**5.10 Theorem** If f is a real continuous function on [a, b] which is differentiable in (a, b), then there is a point  $x \in (a, b)$  at which

$$f(b) - f(a) = (b - a)f'(x).$$

**Proof** Take g(x) = x in Theorem 5.9.

- **5.11** Theorem Suppose f is differentiable in (a, b).
  - (a) If  $f'(x) \ge 0$  for all  $x \in (a, b)$ , then f is monotonically increasing.
  - (b) If f'(x) = 0 for all  $x \in (a, b)$ , then f is constant.
  - (c) If  $f'(x) \le 0$  for all  $x \in (a, b)$ , then f is monotonically decreasing.

Proof All conclusions can be read off from the equation

$$f(x_2) - f(x_1) = (x_2 - x_1)f'(x),$$

which is valid, for each pair of numbers  $x_1$ ,  $x_2$  in (a, b), for some x between  $x_1$  and  $x_2$ .

### THE CONTINUITY OF DERIVATIVES

We have already seen [Example 5.6(b)] that a function f may have a derivative f' which exists at every point, but is discontinuous at some point. However, not every function is a derivative. In particular, derivatives which exist at every point of an interval have one important property in common with functions which are continuous on an interval: Intermediate values are assumed (compare Theorem 4.23). The precise statement follows.

**5.12 Theorem** Suppose f is a real differentiable function on [a, b] and suppose  $f'(a) < \lambda < f'(b)$ . Then there is a point  $x \in (a, b)$  such that  $f'(x) = \lambda$ .

A similar result holds of course if f'(a) > f'(b).

**Proof** Put  $g(t) = f(t) - \lambda t$ . Then g'(a) < 0, so that  $g(t_1) < g(a)$  for some  $t_1 \in (a, b)$ , and g'(b) > 0, so that  $g(t_2) < g(b)$  for some  $t_2 \in (a, b)$ . Hence g attains its minimum on [a, b] (Theorem 4.16) at some point x such that a < x < b. By Theorem 5.8, g'(x) = 0. Hence  $f'(x) = \lambda$ .

### Computation of $\pi_4$ of simple Lie groups



Below we assume any simple Lie group G to be simply connected.



 $\pi_3(G) = \mathbb{Z}$  for any simple Lie group G and there is a uniform proof for that.



Now the textbooks say  $\pi_4(G)$  is trivial except for G = Sp(n), for which it is  $\mathbb{Z}/2\mathbb{Z}$ .



My question is the following: is there a uniform way to derive this fact, in such a way to show which property of the symplectic group is so special?

₩ 4

at.algebraic-topology lie-groups homotopy-theory

share cite improve this question

edited 3 hours ago



2 you probably mean "for G locally isomorphic to Sp(n)" – YCor 8 hours ago

Thanks, corrected. - Yuji Tachikawa 3 hours ago

add a comment

### 2 Answers

active

oldest

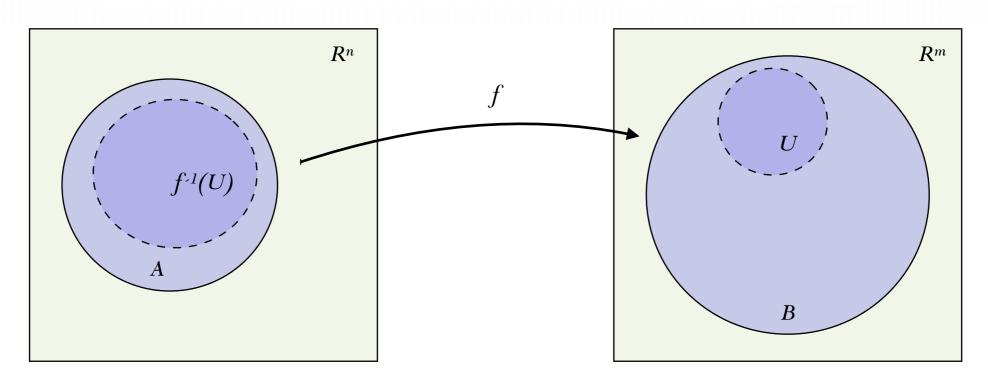
votes



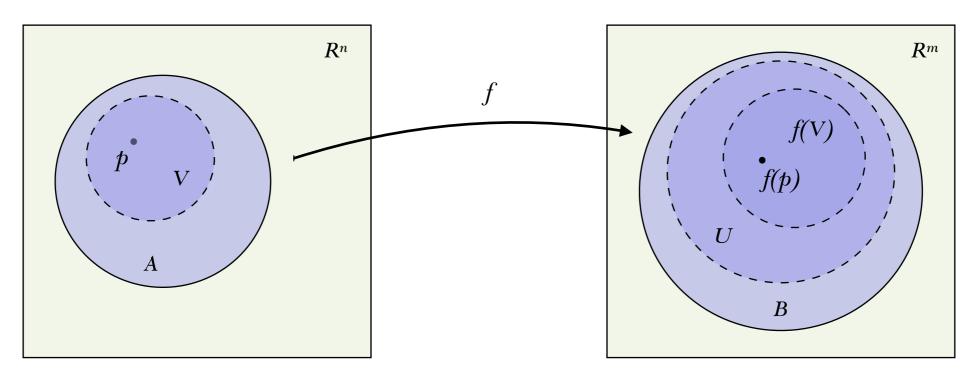
As a partial answer, here's at least a uniform statement, which can be found as Theorem 3.10 in Mimura's survey "Homotopy theory of Lie groups" in the Handbook of Algebraic Topology:

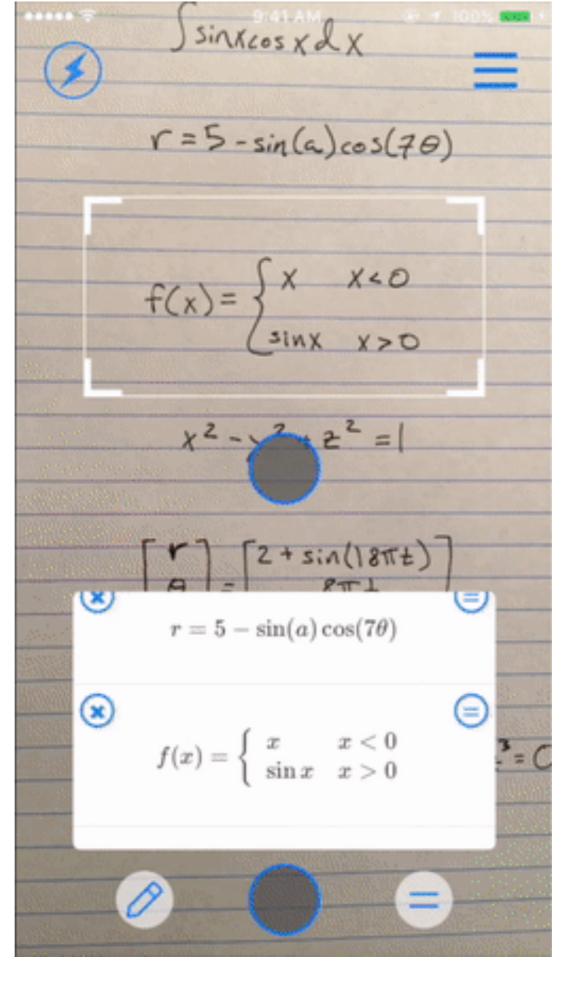
**Proposition 1.3.3.** Let  $A \subset \mathbb{R}^n$  and  $B \subset \mathbb{R}^m$  be sets, and let  $f: A \to B$  be a map. The following statements are equivalent.

(1) for every open subset  $U \subset B$ , the set  $f^{-1}(U)$  is open in A.



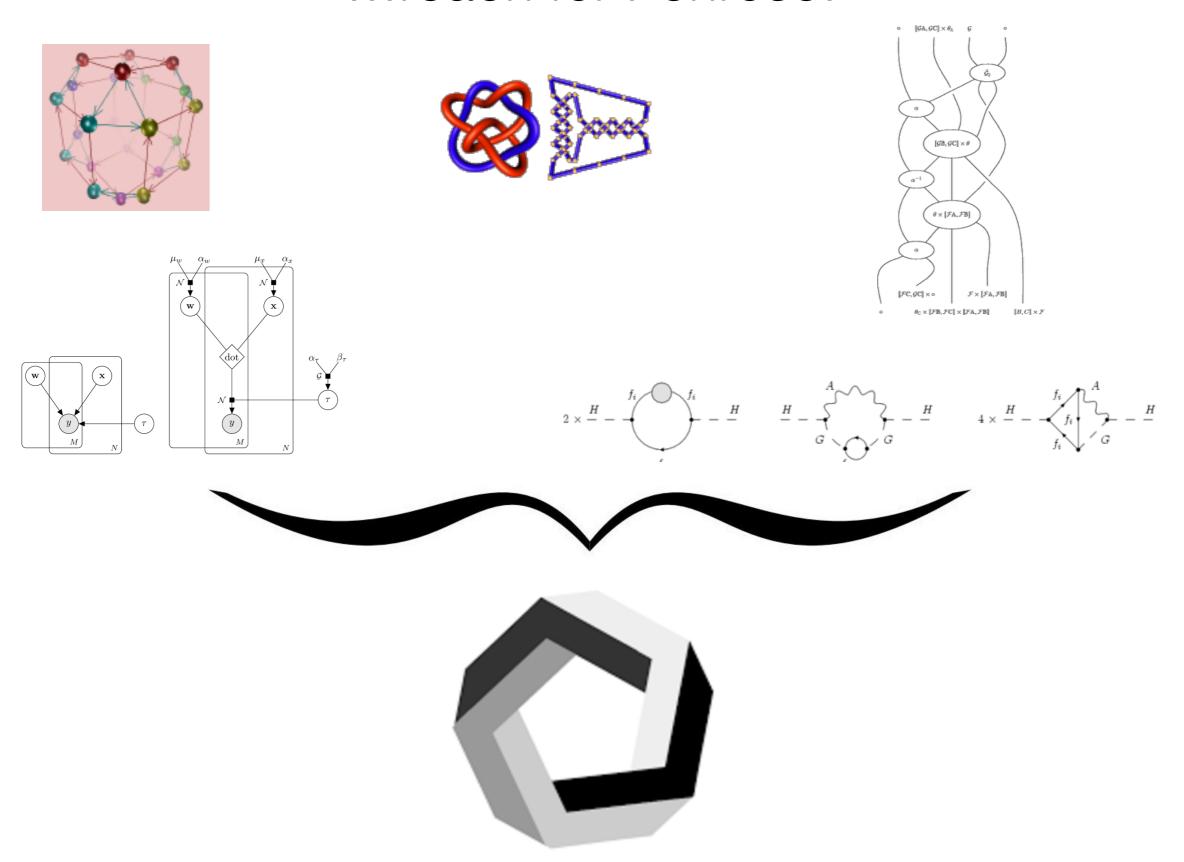
(2) For every point  $p \in A$ , and every open subset  $U \subset B$  containing f(p), there is an open subset  $V \subset A$  containing p such that  $f(V) \subset U$ .





Next time you want to illustrate your paper or talk...

### ...reach for Penrose!



# Gallery

### illustrating a theorem by example

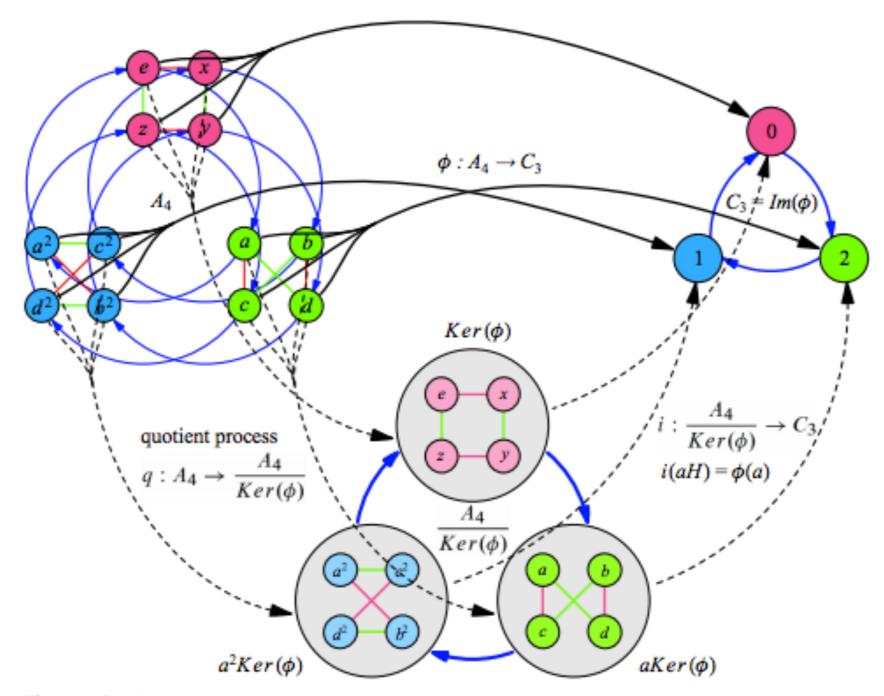


Figure 8.14. The Fundamental Homomorphism Theorem (Theorem 8.5) exemplified using the group  $A_4$  and the quotient map  $\phi$  whose kernel is the subgroup  $\{e, x, y, z\}$  of  $A_4$ 

### illustrating an algorithm

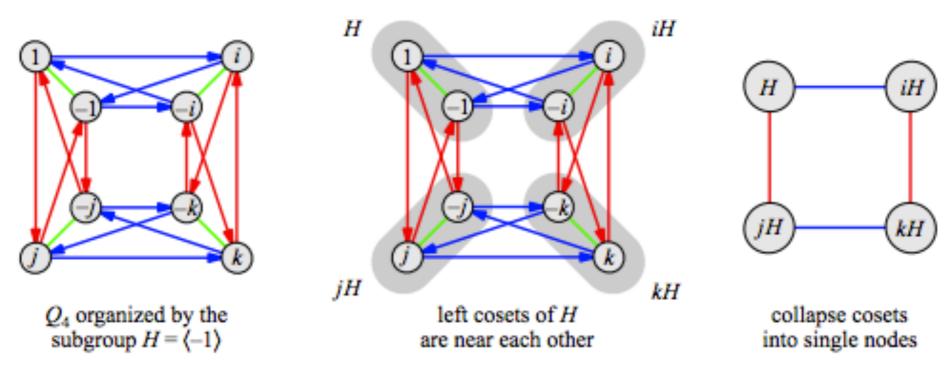


Figure 8.12. The quotient of  $Q_4$  by  $C_2$  corresponding to the homomorphism  $\tau_1$  in Figure 8.9. It follows the quotient procedure from Chapter 7, as depicted in Figure 7.20 (and others).

### illustrating an algorithm

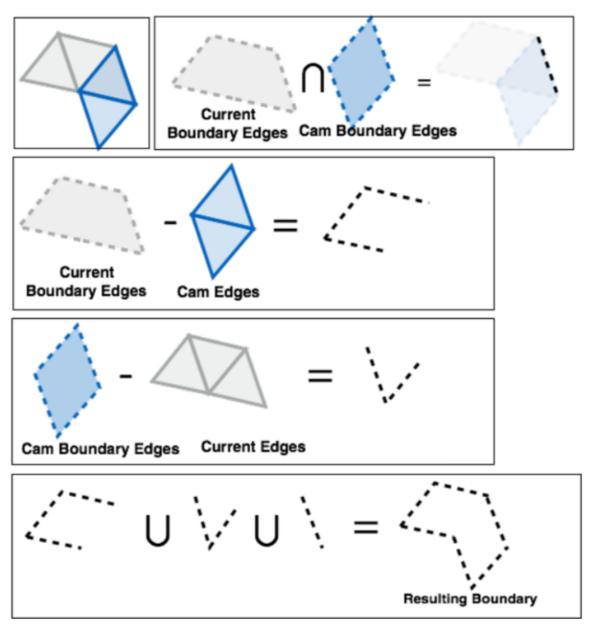


Figure 3. Illustration of calculating the new boundary of the surface area seen by cameras after adding a new camera.

### more geometric diagrams

2.4 Surfaces via Gluing

69

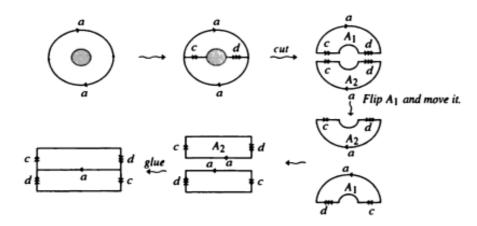


Figure 2.4.13



Figure 2.4.14

#### Exercises

**2.4.1.** Give a polygonal disk and a gluing scheme that will yield the surface pictured in Figure 2.4.15.



Figure 2.4.15

2.4.2\*. Prove Lemma 2.4.5.

### Source:

A first course in geometric topology

We need to find simplicial complexes whose underlying spaces are familiar objects such as  $T^2$  and  $P^2$ . For  $T^2$  this is easy, since it sits in  $\mathbb{R}^3$ ; see Figure 3.3.5. Surfaces that do no sit in  $\mathbb{R}^3$ , such as  $P^2$ , are harder to work with. We will eventually solve this problem using the following construction, which is a simplicial analog of quotient maps and identification spaces (discussed in Section 1.4).

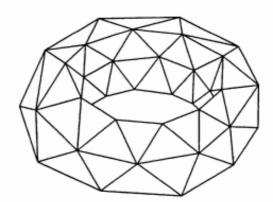
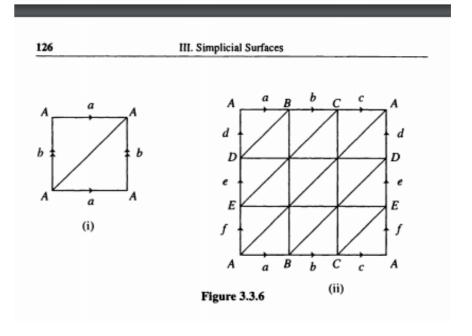
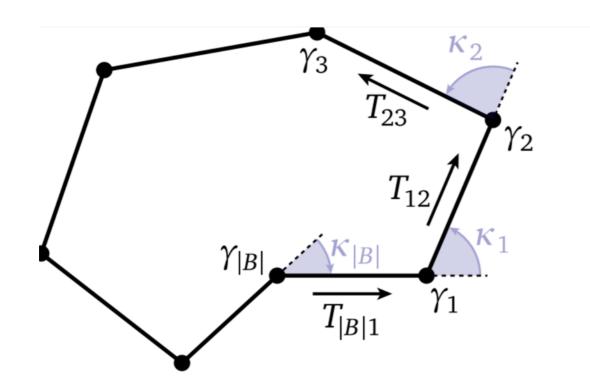


Figure 3.3.5

Consider the way in which  $T^2$  is formed out of gluing the edges of a square, as described in Section 2.4. If we want to obtain a simplicial complex with an underlying space that is a torus, it would be tempting to break up the square shown in Figures 2.4.2 and 2.4.3 into two 2-simplices, as shown in Figure 3.3.6 (i). Unfortunately, when the edges of this square are identified as prescribed by the gluing scheme, all three vertices of both triangles are identified to a single point; since a 2-simplex must have three distinct vertices, we have not produced a simplicial complex by this process of breaking up the original square and then gluing. However, if we break up the original square into 2-simplices a bit more



### geometric diagram described by notation



$$G = (V, E) G = C_{|B|}$$

$$\gamma : V \to \mathbb{R}^{2}$$

$$\kappa : V \to \mathbb{R}$$

$$T : E \to S^{1} \subset \mathbb{R}^{2}$$

$$T_{ij} := \frac{\gamma_{j} - \gamma_{i}}{|\gamma_{j} - \gamma_{i}|}$$

$$\kappa_{j} = \arg(T_{jk} T_{ij}^{-1})$$

### illustrating maps (functions) between objects

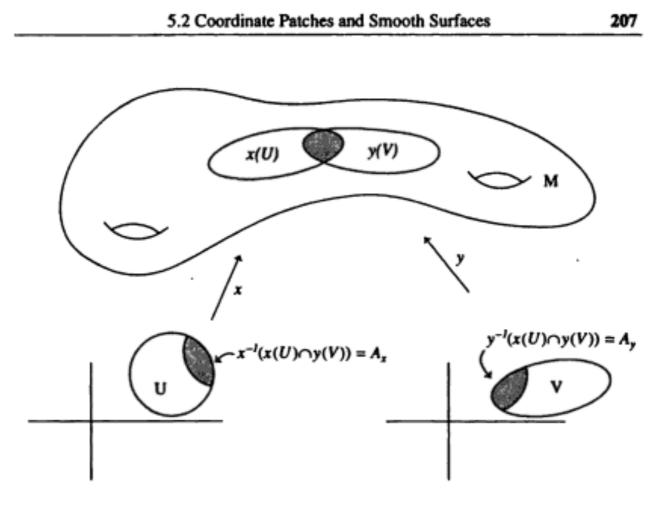
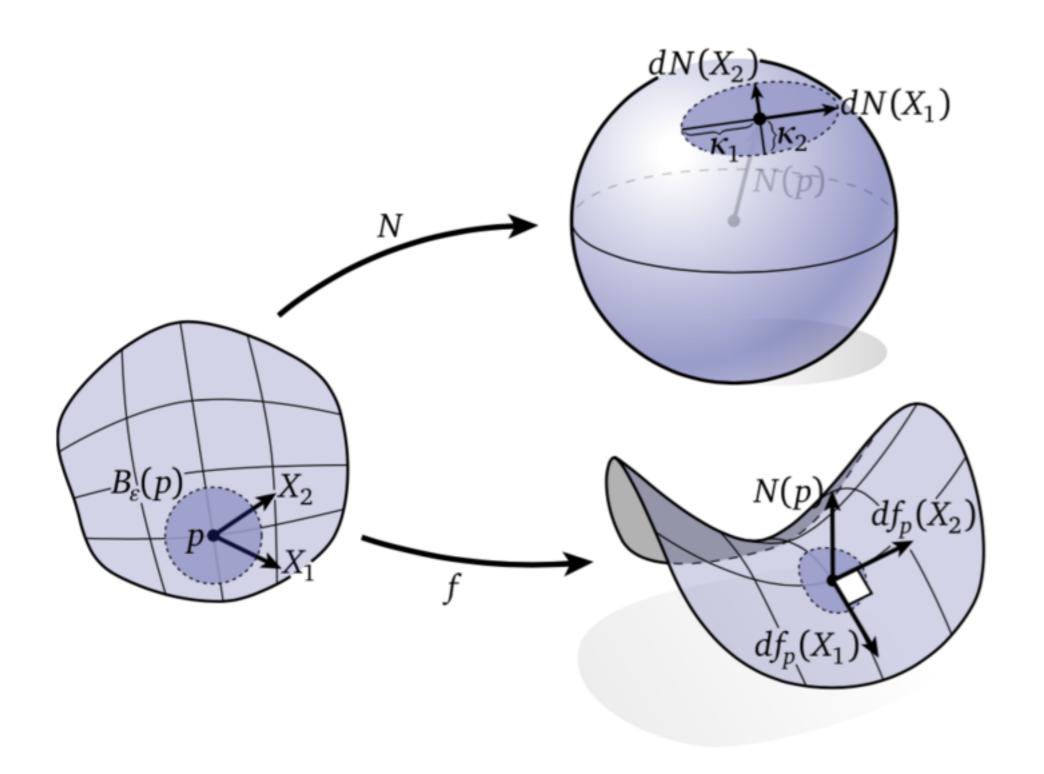


Figure 5.2.2

Since x and y in the above definition are injective, they are bijections onto their images, so we can validly refer to maps  $x^{-1}$  and  $y^{-1}$  in the above definition (though only as maps of sets, with no mention of differentiability). It is easy to see that  $\phi_{x,y}$  is bijective. Also, note that

### illustrating maps (functions) between objects



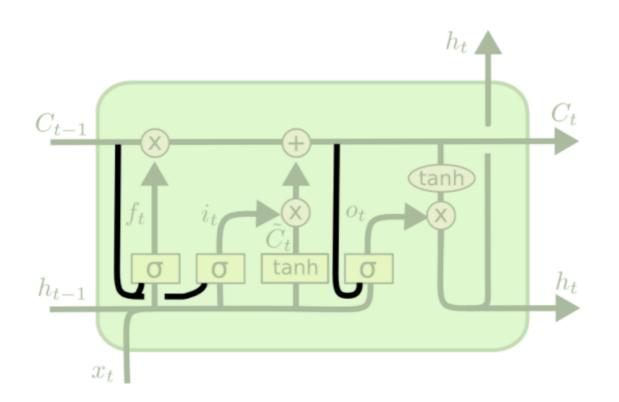
### function composition

**Proof**: The 'only if' statement is obvious since  $f_* = p_*\widetilde{f}_*$ . For the converse, let  $y \in Y$  and let y be a path in Y from  $y_0$  to y. The path fy in X starting at  $x_0$ has a unique lift  $\widetilde{fy}$  starting at  $\widetilde{x}_0$ . Define  $\widetilde{f}(y) = \widetilde{fy}(1)$ . To show this is welldefined, independent of the choice of  $\gamma$ , let  $\gamma'$  be another path from  $\gamma_0$  to  $\gamma$ . Then  $(f\gamma')\cdot(\overline{f\gamma})$  is a loop  $h_0$  at  $x_0$  with  $[h_0]\in f_*(\pi_1(Y,y_0))\subset p_*(\pi_1(\widetilde{X},\widetilde{x}_0))$ . This means there is a homotopy  $h_t$  of  $h_0$  to a loop  $h_1$  that lifts to a  $f \gamma'$  $\widetilde{f}(y)$ loop  $\widetilde{h}_1$  in  $\widetilde{X}$  based at  $\widetilde{x}_0$ . Apply the covering homotopy property to  $h_t$  to get a lifting  $h_t$ . Since  $h_1$  is a loop at  $\widetilde{x}_0$ , so is  $\widetilde{h}_0$ . By the uniqueness of lifted paths, the first half of  $\widetilde{h}_0$  is  $\widetilde{f}\widetilde{\gamma}'$  and the second half is  $\widetilde{fy}$  traversed backwards, with f(y)the common midpoint  $\widetilde{fy}(1) =$  $\widetilde{f\gamma}'(1)$ . This shows that  $\widetilde{f}$  is

Source: Algebraic Topology Hatcher

well-defined.

### circuit-type diagrams for equations



$$f_t = \sigma \left( W_f \cdot [\boldsymbol{C_{t-1}}, h_{t-1}, x_t] + b_f \right)$$

$$i_t = \sigma \left( W_i \cdot [\boldsymbol{C_{t-1}}, h_{t-1}, x_t] + b_i \right)$$

$$o_t = \sigma \left( W_o \cdot [\boldsymbol{C_t}, h_{t-1}, x_t] + b_o \right)$$

### circuit-type diagrams for equations

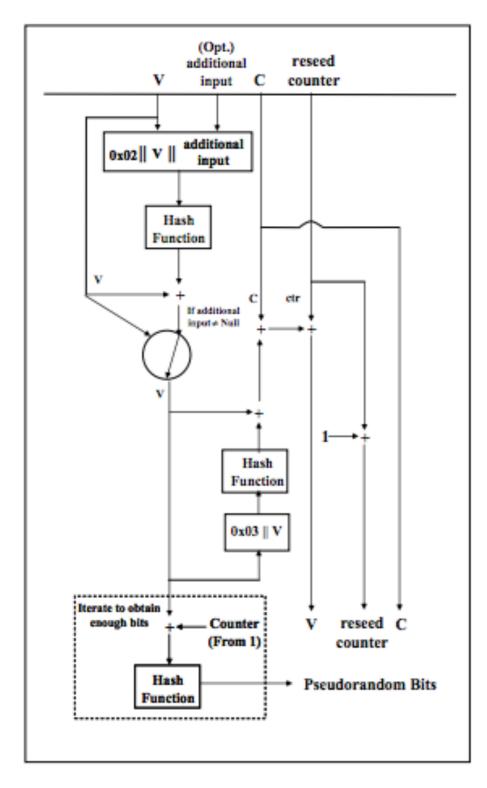


Figure 8: Hash\_DRBG

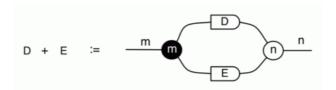
#### Hash\_DRBG Generate Process:

- If reseed\_counter > reseed\_interval, then return an indication that a reseed is required.
- 2. If  $(additional\_input \neq Null)$ , then do 2.1  $w = \mathbf{Hash} (0x02 \parallel V \parallel additional\_input)$ . 2.2  $V = (V + w) \mod 2^{seedlen}$ .
- 3. (returned\_bits) = **Hashgen** (requested\_number\_of\_bits, V).
- 4.  $H = \mathbf{Hash} \ (0x03 \parallel V)$ .
- 5.  $V = (V + H + C + reseed\_counter) \mod 2^{seedlen}$ .
- 6.  $reseed\_counter = reseed\_counter + 1$ .
- Return SUCCESS, returned\_bits, and the new values of V, C, and reseed\_counter for the new\_working\_state.

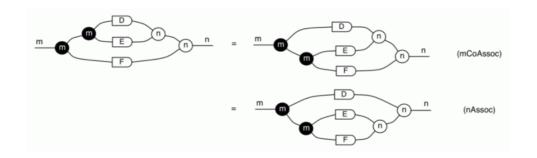
Source: NIST 800-90A 215

### math done via diagrams

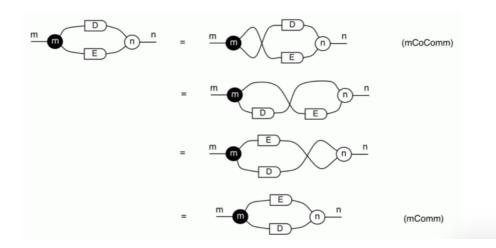
These properties intrigued Fibonchi, and he went ahead and defined an operation on *compatible* diagrams—those that agreed in their respective numbers of dangling wires on the left and right. He called the operation **sum**. This operation would be rediscovered in different notation a few centuries later and given the name **matrix addition**.



He proved that sum was associative



and commutative:



### sometimes we don't know what we want to draw...

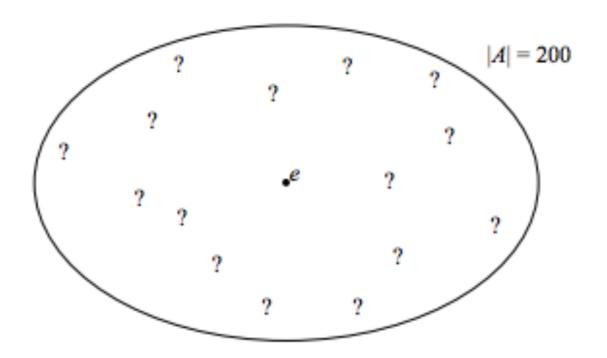


Figure 9.1. If we assume only that a group A has order 200, we do not yet know anything about its internal structure. This chapter teaches several theorems that will reveal much of that structure.

# "Testimonials"

(from real people!)

I would certainly be an eager user of such a system.

The current state of the art for my workflow is a lot of copy-pasting and tweaking of TikZ code.

Edward Morehouse category theorist

This looks fabulous! (...)

I'd love to use it to rewrite Group Explorer for the web!

Nathan Carter author of Visual Group Theory



# FAQ

How will you evaluate the effectiveness of Penrose?

How does Penrose compare to Asymptote, TikZ, Ipe, etc.?

How does Penrose compare to Mathematica, Wolfram Alpha, GeoGebra, graphing calculators, etc.?

What other domains of math might you target?

How will Penrose deal with continuous domains?

How much of the system have you actually implemented so far?

Do you plan to put Penrose on the web, or integrate it with TeX?

What rendering backend will you use?

Have you considered <related work>?

Can I use Penrose right now? No? Then when??

How much about mathematics does Penrose need to "know"?

Have you considered hooking it up to a system for formal reasoning, such as a theorem prover (e.g. Coq)?

What about programs that represent invalid, incorrect, or inconsistent diagrams?

How can Penrose enable counterfactual or nonconstructive reasoning?

(Comment: "I found the anti-example of continuity more helpful than the illustrations of the definition.")

How will you deal with quantifiers (e.g. for all, there exists) and their nesting and ordering?

How will you separate Substance from Style? I don't find HTML/CSS to be very convincing.

Should all branches of mathematics be illustrated? Some seem more amenable to illustration than others.

How easy would it be for you to extend Penrose to do interactive diagrams, animated diagrams, sequential diagrams, algorithmic diagrams, or parameterized diagrams? What about illustrating theorems and proofs?

# I don't know!

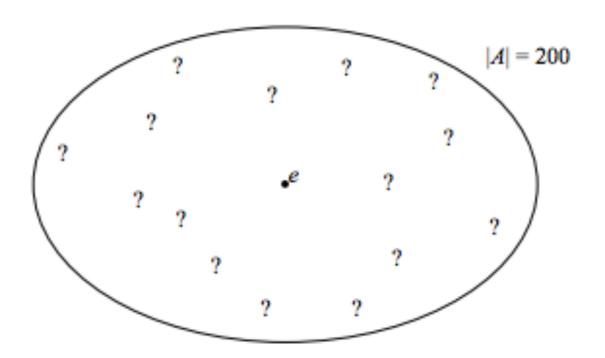


Figure 9.1. If we assume only that a group A has order 200, we do not yet know anything about its internal structure. This chapter teaches several theorems that will reveal much of that structure.