

Programmable MCMC with Soundly Composed Guide Programs

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OOPSLA 2024

October 25, 2024

Probabilistic Programming

1. Probabilistic model $p(\boldsymbol{\theta}, \mathbf{y})$ as a (model) program

- $\boldsymbol{\theta}$: latent variable
- \mathbf{y} : observed variable

May contain conditional statements and loops

2. Posterior distribution by Bayes' rule:

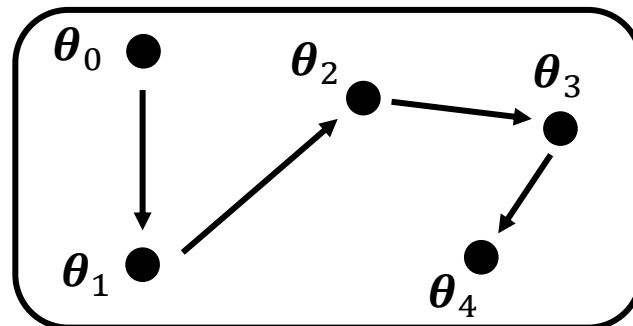
$$p(\boldsymbol{\theta} \mid \mathbf{y} = D) = \frac{p(\boldsymbol{\theta}, \mathbf{y} = D)}{\int p(\boldsymbol{\theta}, \mathbf{y} = D) d\boldsymbol{\theta}}$$

Integral over the space of $\boldsymbol{\theta}$ is difficult

Sampling-based inference algorithms (e.g., Markov-chain Monte Carlo)

Proposal distribution

$$\boldsymbol{\theta}_1^* \sim q(\boldsymbol{\theta} \mid \boldsymbol{\theta}_0)$$
$$\boldsymbol{\theta}_1 := \begin{cases} \boldsymbol{\theta}_1^* & \text{some probability} \\ \boldsymbol{\theta}_0 & \text{otherwise} \end{cases}$$



Example: Bayesian Polynomial Regression

Model program probabilistically generates a polynomial degree and coefficients

```
 $d \sim \text{Categorical}(0.3, 0.5, 0.2)$   
 $c_0 \sim \text{Normal}(\dots)$   
if  $d \geq 1$  then  
   $c_1 \sim \text{Normal}(\dots)$   
  if  $d \geq 2$  then  
     $c_2 \sim \text{Normal}(\dots)$   
 $y \sim \text{polynomial}_{\vec{c}}(x)$ 
```

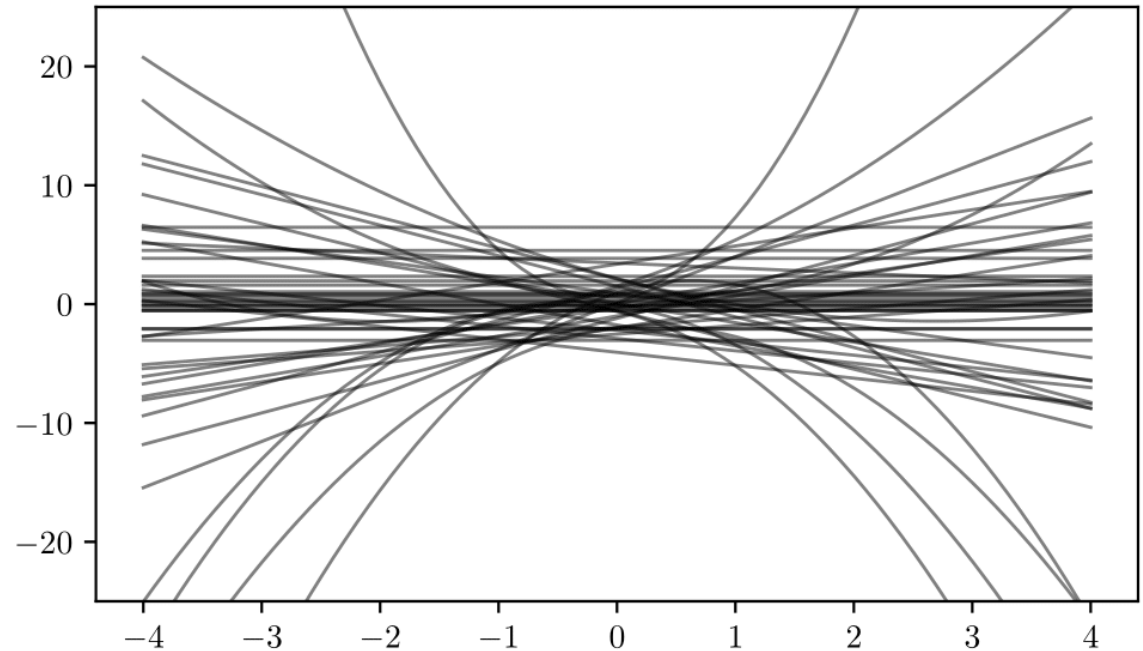
1. Polynomial degree d
2. Constant coefficient c_0
3. Generate coefficients c_1 and c_2 if necessary
4. Generate observed variable y

The set of random variables depends on the execution path

Example: Bayesian Polynomial Regression

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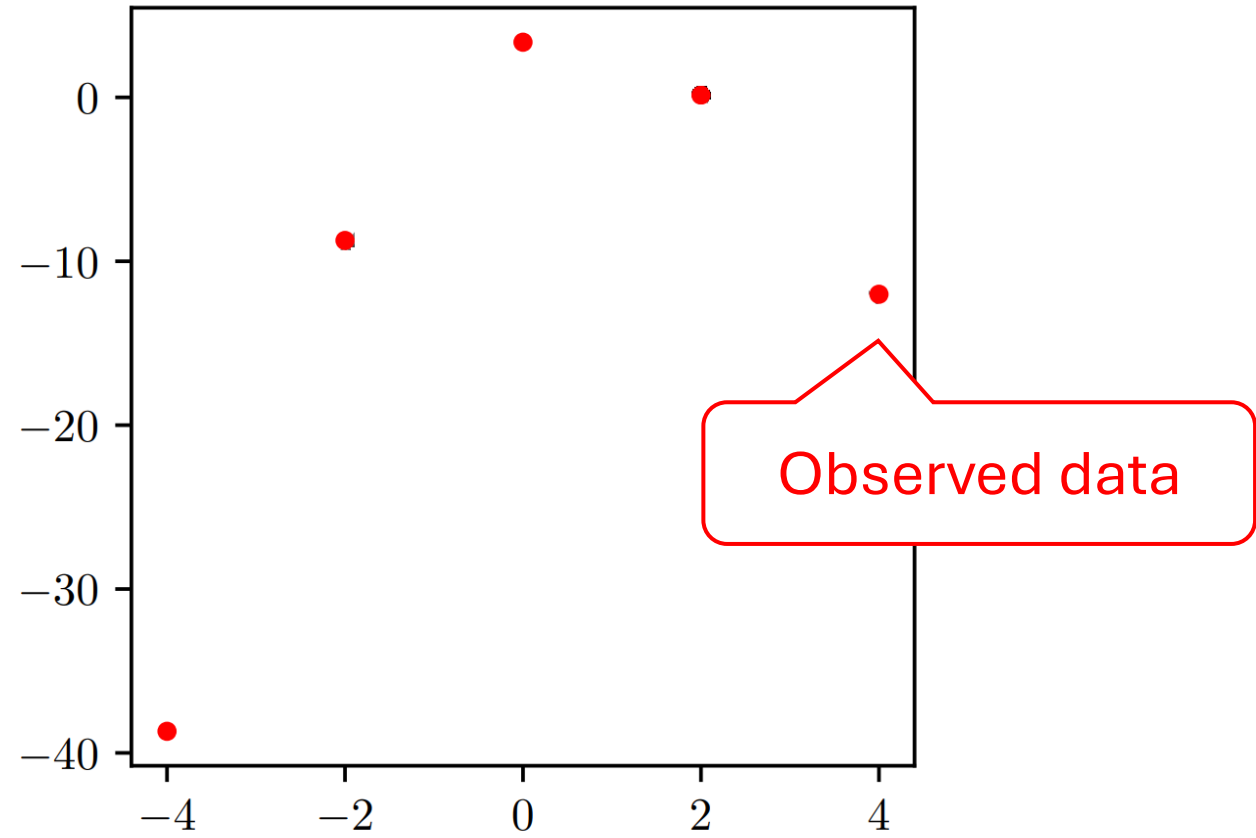
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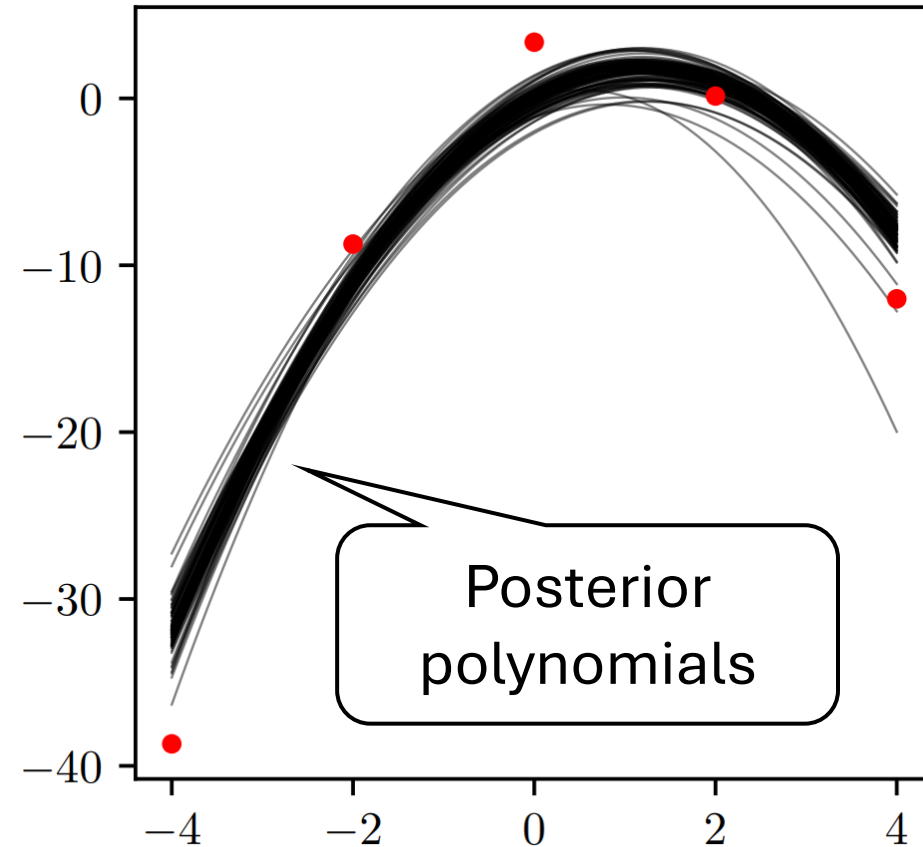
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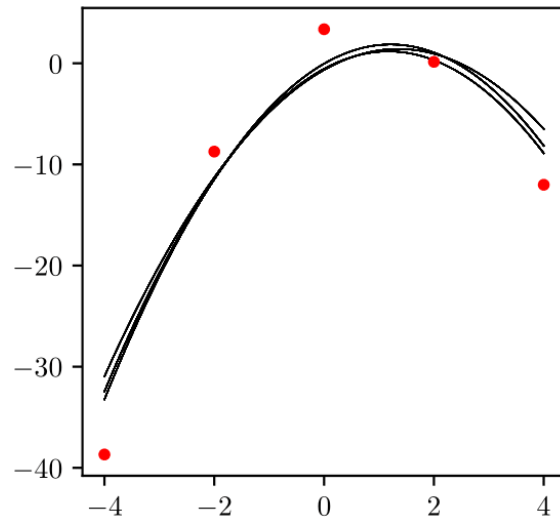
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Programmable Inference: Guide Programs

Traditional PPLs use a **generic** inference algorithm, but it may perform poorly



Little variance within
50 posterior samples

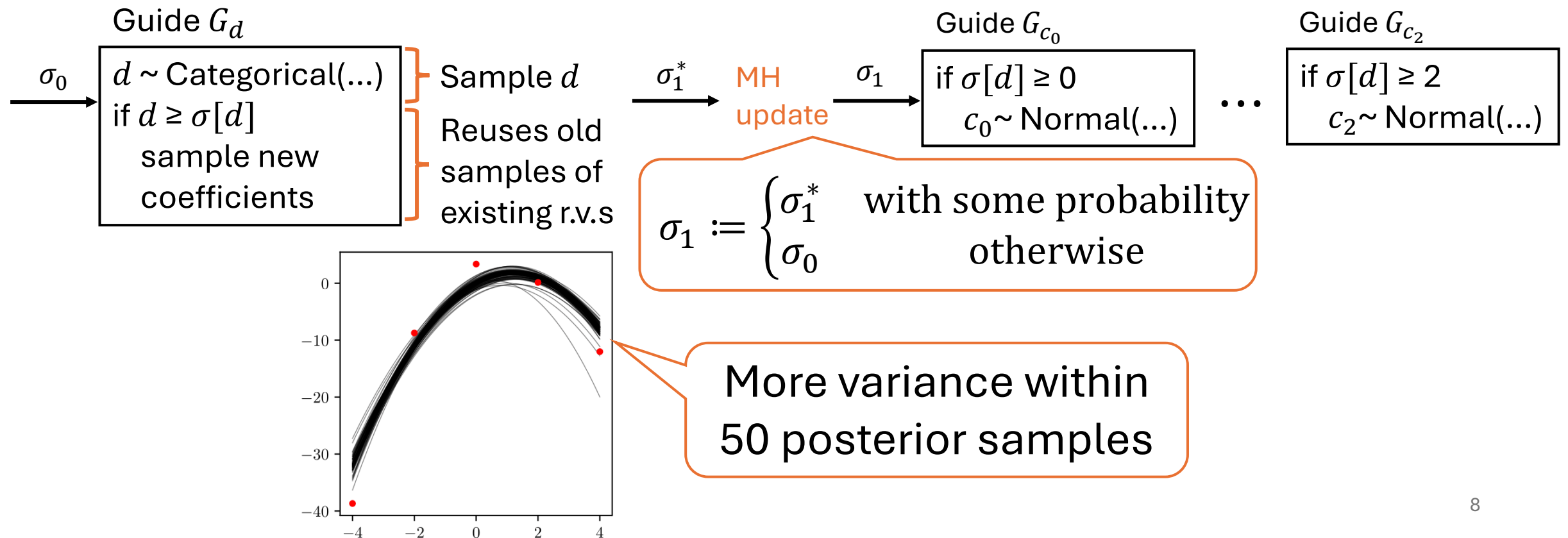
Modern PPLs (e.g., Gen and Pyro) support **programmable** inference where the inference algorithm is customized by a user-written guide program

It customizes a proposal distribution for MCMC

Programmable Inference: BMH

Multiple-Block Metropolis-Hastings (BMH) samples r.v.s **block by block**

- Split latent variables into four blocks: d, c_0, c_1, c_2
 - Specify proposal distributions for each block
- } Guide programs



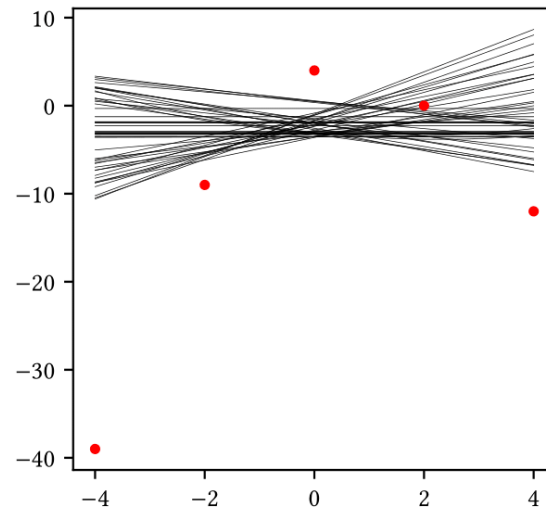
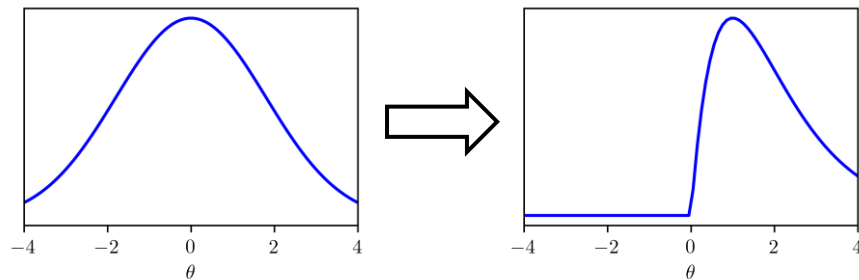
Contribution: Soundness Verification of BMH

Sufficient soundness condition of BMH

Model and guide must have the same support

Unsound BMH: replace a normal distribution with a gamma distribution

Normal distribution Gamma distribution



Guide is unable to sample negative values for coefficient c_2

Our contribution

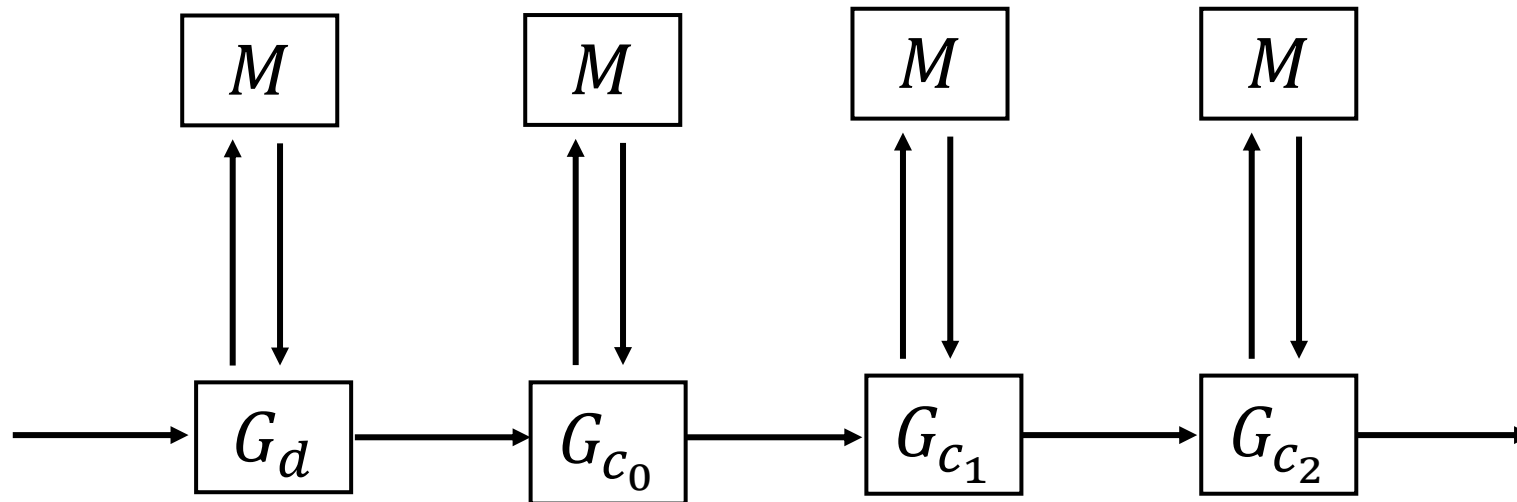
Develop a **type-based framework** for verifying that the model and guide have the same support in BMH

Outline

- Motivation for soundness verification of BMH
- Formulation of BMH
- Polynomial-time decidability of structural guide-type equality
- Coverage checking
- Implementation and evaluation

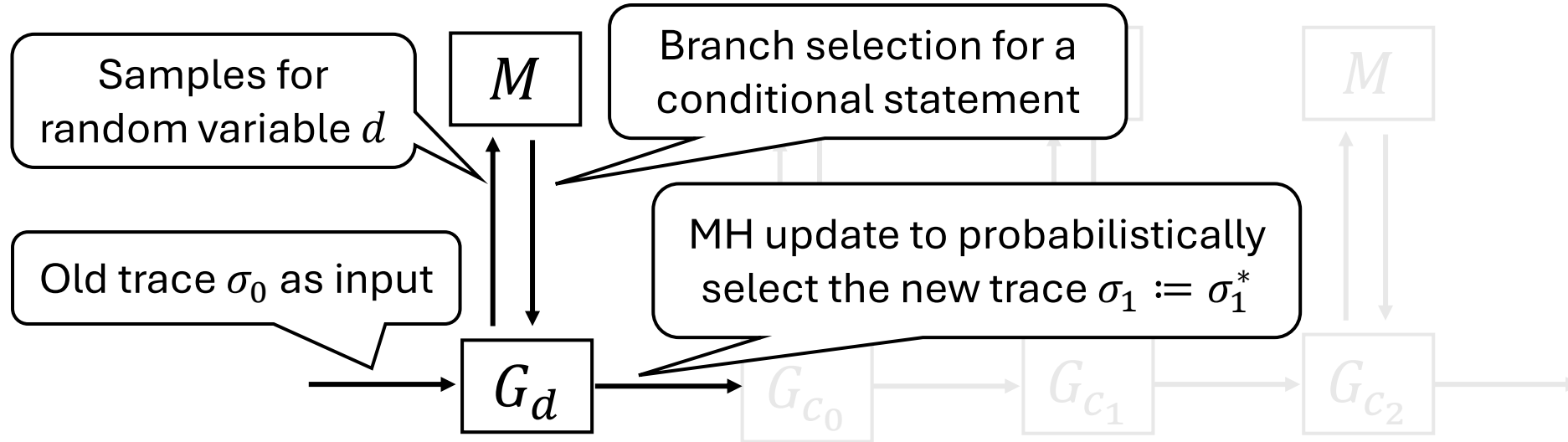
Contribution 1: Formulation of BMH

BMH is a sequential composition of guide programs



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BMH is a sequential composition of guide program



The communication protocol is described by a guide type:

$$\underbrace{\mathbb{N}_3 \wedge \mathbb{R} \wedge \&}_{\text{Distribution type of } d} \left\{ \begin{array}{l} \underbrace{\text{Branch selection}} \left\{ \begin{array}{l} 1, \\ \mathbb{R} \wedge \& \end{array} \right\} \text{Termination} \\ \mathbb{R} \wedge \& \left\{ \begin{array}{l} 1, \\ \mathbb{R} \wedge 1 \end{array} \right\} \end{array} \right.$$

Contribution 2: Structural Guide-Type Equality

Soundness ingredient

The model and guide must have an equal guide type

Context-free guide types

Types indexed by type parameters

$$\underbrace{T[X]} := \mathbb{R} \wedge \left(X \ \& \ \underbrace{T[T[X]]} \right)$$

Run a protocol T , followed
by a continuation X

The recursive call is a
sequential composition of T

Guide types may have infinite state spaces: $T[X]$, $T[T[X]]$, $T[T[T[X]]]$, ...

Contribution 2: Structural Guide-Type Equality

Structural guide-type equality is decidable in polynomial time

1. Translate guide types to context-free processes

Context-free grammars
viewed as processes

2. Require the norms (i.e., #of steps to reach termination) to be finite

3. Theorem (Hirshfeld et al., 1994). Bisimilarity of two context-free processes (i.e., structural guide-type equality) with finite norms is decidable in polynomial time.

Outline

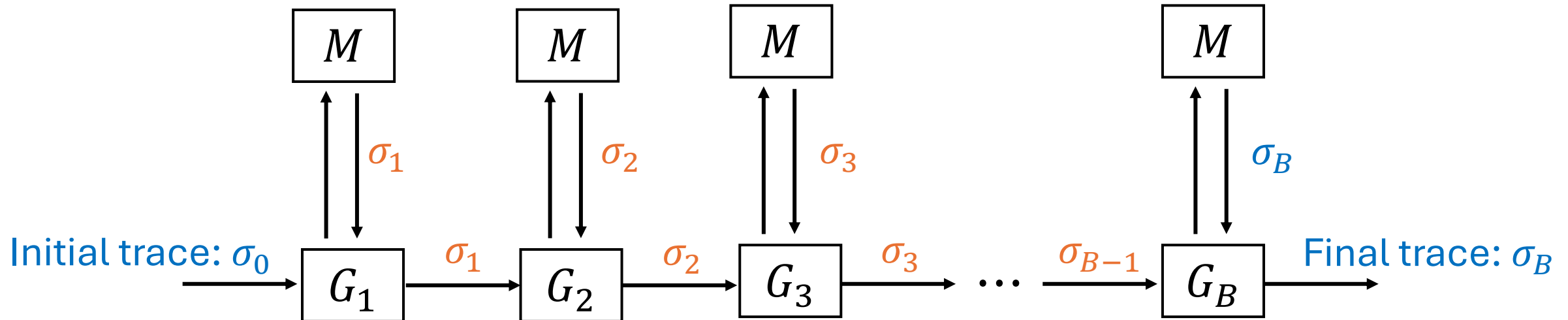
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Contribution 3: Coverage Checking of Composed Guides

Every random variable must be covered (i.e., sampled freshly by at least one guide) on any execution path

Problem statement of coverage checking

For an arbitrary initial trace σ_0 and final trace σ_B , do there exist intermediate traces $\sigma_1, \dots, \sigma_{B-1}$ such that σ_i can be generated from σ_{i-1} ($i = 1, \dots, B$)?



Contribution 3: Coverage Checking of Composed Guides

Coverage-annotated guide types

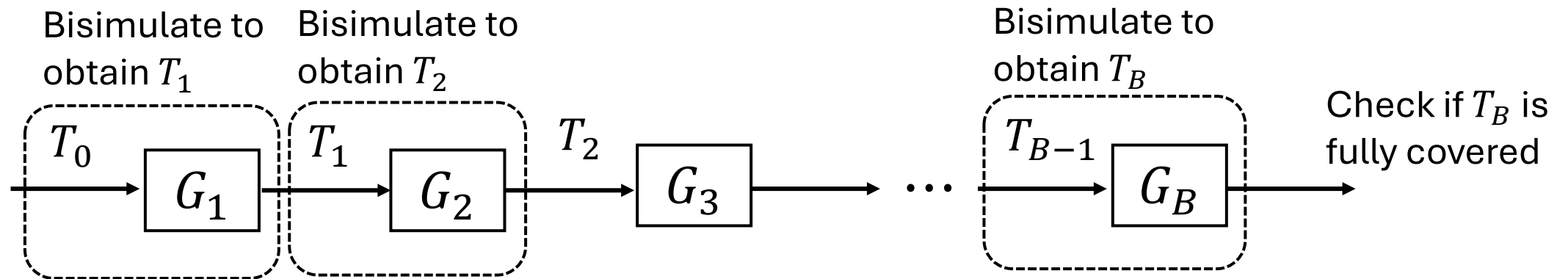
Subscript u means the sample is reused

$$\mathbb{R}_u \wedge \begin{cases} 1, \\ \mathbb{R}_c \wedge 1 \end{cases}$$

Subscript c means the sample is freshly sampled

Coverage-checking algorithm

Starting with a fully uncovered guide type T_0 , we bisimulate a guide type T_{i-1} alongside G_i to update coverage annotations



Implementation and Evaluation

Table 1. Experiment results of guide-type inference and coverage checking of 28 benchmark programs.

Program	Description	Guide Type	LOC	Type Inference		Coverage Check		
				Time (ms)	Prior Work	Match	Mismatch	Time (ms)
branching	Random control flow [Anglican]	Finite	46	1.33	✓	True Pos.	True Neg.	0.46
coordination	Coordination game [Anglican]	Finite	24	0.19	✓	True Pos.	True Neg.	0.34
drill	Oil wildcatter problem [Anglican]	Finite	56	0.17	✓	True Pos.	True Neg.	0.37
ex-1	Ex. 1 [52]	Finite	42	1.31	✓	True Pos.	True Neg.	0.46
gaussian	Gaussian with unknown means [Anglican]	Finite	20	0.16	✓	True Pos.	True Neg.	0.46
gbm	Geometric Brownian motion [Anglican]	Finite	35	0.25	✓	True Pos.	True Neg.	0.52
gda	Gaussian discriminant analysis [Anglican]	Finite	40	1.86	✓	True Pos.	True Neg.	3.17
gmm	Gaussian mixture model [Anglican]	Finite	75	4.73	✓	True Pos.	True Neg.	7.71
grw	Gaussian random walk [Anglican]	Finite	24	0.17	✓	True Pos.	True Neg.	0.74
hmm	Hidden Markov model [Anglican]	Finite	76	2.56	✓	True Pos.	True Neg.	7.21
kalman	Kalman filter [Anglican]	Finite	72	4.44	✓	True Pos.	True Neg.	7.54
kalman-chaos	Kalman for chaotic attractors [Anglican]	Finite	114	5.86	✓	True Pos.	True Neg.	5.68
lr	Bayesian linear regression [Anglican]	Finite	36	0.19	✓	True Pos.	True Neg.	1.15
run-factory	Beta-binomial model [Anglican]	Finite	20	0.13	✓	True Pos.	True Neg.	0.61
scientists	Posterior estimation with Gaussians [54]	Finite	40	0.27	✓	True Pos.	True Neg.	0.52
seq	Non-recursive sequence [52]	Finite	22	0.23	✓	True Pos.	True Neg.	0.46
sprinkler	Bayesian network [Anglican]	Finite	26	0.14	✓	True Pos.	True Neg.	0.43
user-behavior	Dishonest form filling [Anglican]	Finite	64	1.22	✓	True Pos.	True Neg.	3.17
vae	Variational autoencoder [Pyro]	Finite	48	4.20	✓	True Pos.	True Neg.	22.39
weight	Unreliable weight [Pyro]	Finite	18	0.26	✓	True Pos.	True Neg.	0.70
aircraft	Aircraft detection [Anglican]	Regular	117	6.19	✗	True Pos.	True Neg.	5.96
iter	Regular iteration [52]	Regular	47	2.01	✗	True Pos.	True Neg.	0.54
marsaglia	Marsaglia algorithm [Anglican]	Regular	76	3.51	✗	True Pos.	True Neg.	5.13
ptrace	Poisson trace [Anglican]	Regular	47	1.49	✗	True Pos.	True Neg.	0.40
ex-2	Ex. 2 [52]	Context-Free	78	4.77	✗	True Pos.	True Neg.	4.70
			93	15.48	✗	False Neg.	True Neg.	3.26
diter	Double iteration [52]	Context-Free	52	1.48	✗	True Pos.	True Neg.	0.57
			62	2.09	✗	False Neg.	True Neg.	0.49
gp-dsl	Gaussian process DSL [52]	Context-Free	242	879.53	✗	True Pos.	True Neg.	4.71
			261	2487.91	✗	False Neg.	True Neg.	4.59
recur	Context-free recursion [52]	Context-Free	71	11.53	✗	True Pos.	True Neg.	16.32
			83	15.55	✗	False Neg.	True Neg.	6.35

Type equality and inference

Type inference of some BMH benchmarks require structural guide-type equality

Coverage checking

- Effective in benchmarks with regular guide types (i.e., finite state spaces).
- It can return **False Negative** for benchmarks with infinite-state guide types

Conclusion

1. Formulated Multiple-Block Metropolis-Hastings (BMH) in guide-based programmable inference
2. Proved polynomial-time decidability of structural guide-type equality
3. Developed a coverage-checking algorithm for verifying that every random variable is freshly sampled at least once