

Lecture 24: Spectral Graph Theory I

[Questions motivating us now...]: Given undirected graph G ...

- how to "cluster" vertices into 2+ groups?
- how to "visualize" in 2 or 3 dims?
- how long does random walk take to get "close" to stationary dist π ?

Henceforth for simplicity: G is d -regular.

[Everything has appropriate analogues for non-regular graphs, but math & formulas get a little annoying, let's keep things simple!]

$$K = \begin{matrix} \text{adj. matrix} \\ \vdots \\ \vdots \\ \vdots \end{matrix} = \frac{1}{d} A$$

symmetric, $K = K^T$

[Remember I told you linear algebra is 100x easier with symmetric matrices? This is one example.]

Stationary distribution: $\pi[u]$ proportional to $\text{deg}[u]$... $\therefore \pi = \begin{bmatrix} \frac{1}{n} \\ \vdots \\ \frac{1}{n} \end{bmatrix}$ = uniform dist. 😊

$$K\pi = \pi \Rightarrow \pi \text{ is eigenvector of } K, \text{ eigenvalue } 1.$$

[we'll talk about eigenvectors more soon]

Say for each vertex $v \in V$ we have a real # data value $f(v)$.

[Could be anything - temperature, color, height, x-coord. in a plot...]

$f: V \rightarrow \mathbb{R}$

define: $E[f] = \frac{1}{2} \sum_{(u,v) \text{ edge}} [(f(u) - f(v))^2]$

↑ average

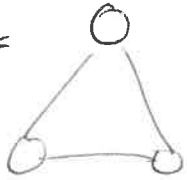
↑ (u,v) edge

"average variance on edges"

"energy"
"quadratic form"

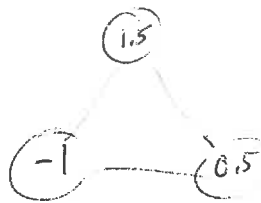
[don't worry about too much, just convenient to have this factor]

eg: $G =$

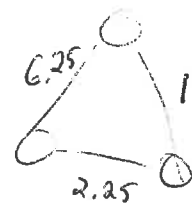


$d=2$

labeled by f



edges labeled by squared diff



(2)

$$\text{avg} = \frac{9.5}{3} \approx 3.166$$

$$\Rightarrow E[f] \approx 1.584$$

Whatever, just an example

**
eg.
**

Let $S \subseteq V$. Let $f(v) = \begin{cases} 1 & \text{if } v \in S \\ 0 & \text{if } v \notin S \end{cases}$

What's $E[f]$?

$$E[f] = \frac{1}{2} \sum_{(u,v) \text{ edge}} [f(u) - f(v)]^2$$

0 if $u, v \in S$ or $u, v \notin S$

$(\pm 1)^2 = 1$ if u, v on "opp. sides": (u, v) crosses "cut" (S, \bar{S})

$$= \frac{1}{2} \sum_{(u,v) \text{ edge}} [u \in S, v \notin S \text{ or } u \notin S, v \in S]$$

$$= \Pr_{(u,v) \text{ edge (directed)}} [u \in S, v \notin S]$$

" $u \rightarrow v$ is stepping out of S "

"fractional size of boundary of S "



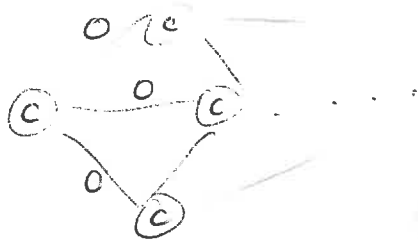
[Has something to do with how long it takes random walk to "escape out of S "]

Fact: $E[f] \geq 0$ always. [Why? Well, it's avg. of squares, ≥ 0 .]

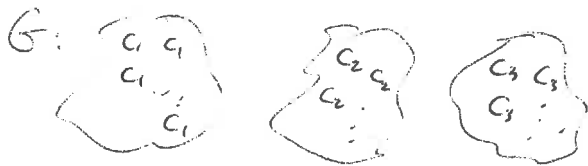
Q: For what f is $E[f] = 0$?

A: $f(v) = 0 \forall v$, $f(v) = 1 \forall v$, $f(v) = c \forall v$ [Anything else?]

That's it if G is connected.



Say now G might be disconnected.



We'll think of $f: V \rightarrow \mathbb{R}$ as a vector: $\begin{bmatrix} f(v_1) \\ f(v_2) \\ \vdots \\ f(v_n) \end{bmatrix} \in \mathbb{R}^n$.

Say G has 3 connected comps, S_1, S_2, S_3 .

Then \exists 3 linearly independent f 's with $E[f] = 0$:

$$s_1 \begin{bmatrix} 1 \\ \vdots \\ 0 \\ \vdots \\ 0 \end{bmatrix}, s_2 \begin{bmatrix} 0 \\ \vdots \\ 1 \\ \vdots \\ 0 \end{bmatrix}, s_3 \begin{bmatrix} 0 \\ \vdots \\ 0 \\ \vdots \\ 1 \end{bmatrix}$$

[and every f with $E[f] = 0$ is a lin. combo of them, coeffs c_1, c_2, c_3 .

Connection between linear algebra & graph theory!

What's largest $E[f]$ can be? [Well... ∞ . We need to fix some "scaling".]

\rightarrow Multiplying f by k multiplies $E[f]$ by k^2 . [So the natural way to scale is also quadratic.]

Say f "normalized" if $E_{v \text{ vertex}} [f(v)^2] = 1$.

[Can always mult by a const. to achieve this.]



[Trying to map vertices $v \in V$ into real line - avg. squared dist to 0 = 1 - so that edges "stretched" as much as possible.]

What graph might ~~allow~~ ~~this~~ for a very large E ?

Say G is bipartite: vertices U, V , all edges between them

$$\text{Let } f(w) = \begin{cases} +1 & \text{if } w \in U \\ -1 & \text{if } w \in V. \end{cases}$$

Normalized? \checkmark

$$E[f(w)^2] = E[(\pm 1)^2] = 1.$$

$$E[f] = \frac{1}{2} E_{(u,v) \text{ edge}} [(f(u) - f(v))^2] = 2.$$

Prop: For normalized f , $E[f] \leq 2$ always; equality possible iff G is bipartite.

Proof: $E[f] = \frac{1}{2} \left(\sum_{(u,v) \text{ edge}} [f(u)^2] - 2 \sum_{(u,v) \text{ edge}} [f(u)f(v)] + \sum_{(u,v) \text{ edge}} [f(v)^2] \right)$

normalized \downarrow
 u is a unif. rand. vertex (regularity \uparrow)

$= \frac{1}{2} (1) - \frac{1}{2} \cdot 2 \sum_{(u,v) \text{ edge}} [f(u)f(v)] + \frac{1}{2} (1) = 1 - E[f(u)f(v)] \leq 2$ provided...

Claim: $|E[f(u)f(v)]| \leq 1$. Proof: ~~Cauchy-Schwarz~~
 "Cauchy-Schwarz": $|E[XY]| \leq \sqrt{E[X^2]} \sqrt{E[Y^2]}$
 $\leq \sqrt{E[f(u)^2]} \sqrt{E[f(v)^2]} = \sqrt{1} \sqrt{1} = 1$

Spectral Graph Theory: G has k "almost disconnected" clusters

$\Leftrightarrow E[f_1], \dots, E[f_k] \approx 0$ for k lin. independent f_1, \dots, f_k



G is almost bipartite $\Leftrightarrow E[f] \approx 2$ for some f
 (= huge max. cut)



[Where do eigenvalues come in?] [Almost there...] [Ex. if G not bip, how to find f maximizing $E[f]$?] 2

What is

$$u \rightarrow \begin{bmatrix} K \\ -\frac{1}{d} & -\frac{1}{d} & \dots & -\frac{1}{d} \end{bmatrix} \begin{bmatrix} f \\ \vdots \\ f \\ \vdots \\ f \end{bmatrix} = h$$

at any $(u, v_1), (u, v_2), \dots, (u, v_d)$

$h(u) = \frac{1}{d} f(v_1) + \dots + \frac{1}{d} f(v_d)$, v_i 's are neighbors of u .
 $= E[f(v)]$
 v : (u, v) edge

[Applying K to f : replace each data value with average over neighbors.]

def: Inner product: $\langle f, g \rangle = \frac{1}{n} (f \cdot g) = \frac{1}{n} (f(v_1)g(v_1) + \dots + f(v_n)g(v_n))$ (5)
 ("non-standard") $= \mathbb{E}_{\text{vertex}} [f(v)g(v)]$

["correlation between f & g "]

$\therefore \langle f, f \rangle = \mathbb{E}[f(v)^2] = 1$ if f "normalized"

(sets like "unit vec")

Q: What is $\langle f, Kf \rangle$?

A: $\mathbb{E}_{\text{edge}} [f(u) \cdot (Kf)(u)] = \mathbb{E}_{\text{edge}} [f(u) \cdot \mathbb{E}_{\text{edge}} [f(v)]] = \mathbb{E}_{\text{edge}} [f(u)f(v)]$
 (Probably natural?)

Q: What is $\mathbb{E}[f^2]$?

A: $= \frac{1}{2} \mathbb{E}_{\text{edge}} [(f(u) - f(v))^2] = \frac{1}{2} (\mathbb{E}_{\text{edge}} [f(u)^2] - 2 \mathbb{E}_{\text{edge}} [f(u)f(v)] + \mathbb{E}_{\text{edge}} [f(v)^2])$
 $= \frac{1}{2} \langle f, f \rangle - \langle f, Kf \rangle + \frac{1}{2} \langle f, f \rangle$
 $= \langle f, f \rangle - \langle f, Kf \rangle$
 $= \langle f, (I - K)f \rangle = \langle f, Lf \rangle$
 (ident mtr)

def: Laplacian matrix of G : $L := I - K = I - \frac{1}{2}A = \frac{1}{2}(2I - A)$
 ("symmetric")

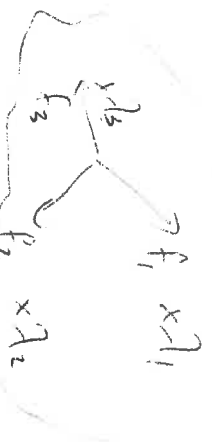
The reason to care about it is not that $I - K$ is inherently meaningful, just that... $\mathbb{E}[f^2] = \langle f, Lf \rangle = \frac{1}{n} \mathbb{E}[\sum_i (L_i f_i)^2]$

Recall: A symmetric matrix like L stretches \mathbb{R}^n by

stretch factors (eigenvalues) $\lambda_1, \dots, \lambda_n$ along axes L
 (orthonormal eigenvectors) f_1, \dots, f_n

Say f_i normalized. $L f_i = \lambda_i \cdot f_i$

$\Rightarrow \langle f_i, L f_i \rangle = \langle f_i, \lambda_i f_i \rangle = \lambda_i \mathbb{E}[f_i(v)^2] = \lambda_i$



$\therefore \mathbb{E}[f_i^2] = \lambda_i$. \therefore all λ_i are between 0 & 2.

(well known)

G has k connected components $\Leftrightarrow k$ of the \mathbb{Z}_2 are 0

G is bipartite \Leftrightarrow one of the \mathbb{Z}_2 is 2.



⑥