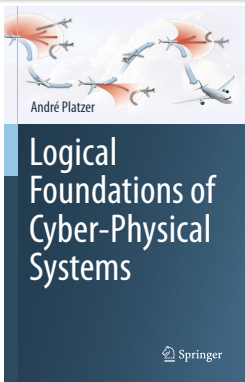
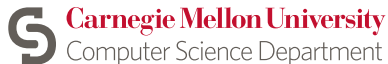


07: Control Loops & Invariants

Logical Foundations of Cyber-Physical Systems



Stefan Mitsch



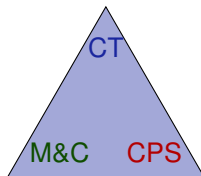
- 1 Learning Objectives
- 2 Induction for Loops
 - Iteration Axiom
 - Induction Axiom
 - Induction Rule for Loops
 - Loop Invariants
 - Simple Example
 - Contextual Soundness Requirements
- 3 Operationalize Invariant Construction
 - Bouncing Ball
 - Rescuing Misplaced Constants
 - Safe Quantum
- 4 Summary

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Learning Objectives

Control Loops & Invariants

rigorous reasoning for repetitions
identifying and expressing invariants
global vs. local reasoning
relating iterations to invariants
finitely accessible infinities
operationalize invariant construction
splitting & generalizations



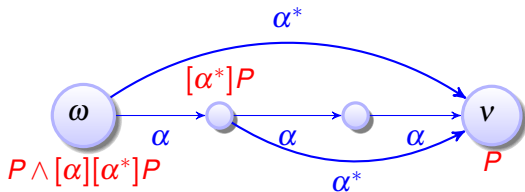
control loops
feedback mechanisms
dynamics of iteration

semantics of control loops
operational effects of control

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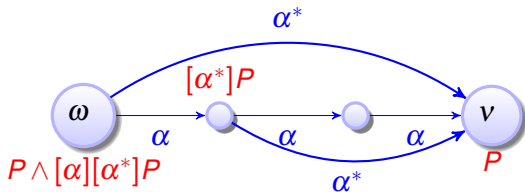
Iteration Axiom

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$



Iteration Axiom

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

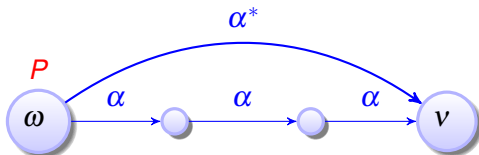


Problem: Proof for $[\alpha^*]P$ needs proof of $[\alpha][\alpha^*]P$

Induction Axiom

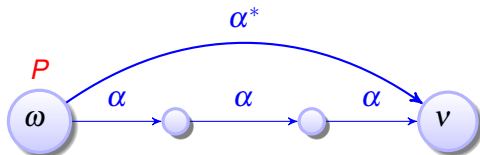
Lemma ()

$$\models [\alpha^*]P \leftrightarrow P \wedge$$



Lemma ()

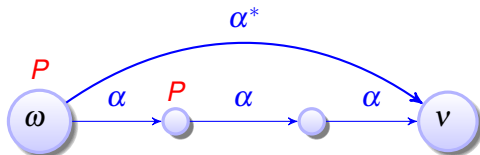
$$\vdash [\alpha^*]P \leftrightarrow P \wedge (P \rightarrow [\alpha]P)$$



$$P \rightarrow [\alpha]P$$

Lemma ()

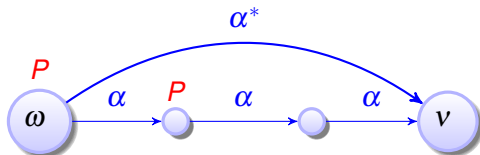
$$\vdash [\alpha^*]P \leftrightarrow P \wedge (P \rightarrow [\alpha]P)$$



$$P \rightarrow [\alpha]P$$

Lemma (I is sound)

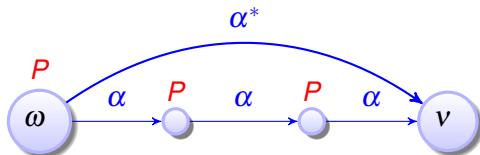
$$\vdash [\alpha^*]P \leftrightarrow P \wedge (P \rightarrow [\alpha]P)$$



$$P \rightarrow [\alpha]P \quad P \rightarrow [\alpha]P$$

Lemma (I is sound)

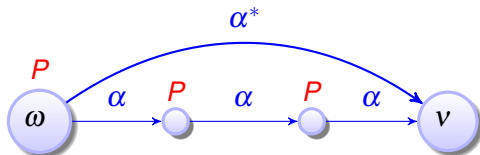
$$\vdash [\alpha^*]P \leftrightarrow P \wedge (P \rightarrow [\alpha]P)$$



$$P \rightarrow [\alpha]P \quad P \rightarrow [\alpha]P$$

Lemma (I is sound)

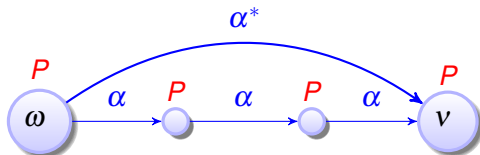
$$\vdash [\alpha^*]P \leftrightarrow P \wedge (P \rightarrow [\alpha]P)$$



$$P \rightarrow [\alpha]P \quad P \rightarrow [\alpha]P \quad P \rightarrow [\alpha]P$$

Lemma (I is sound)

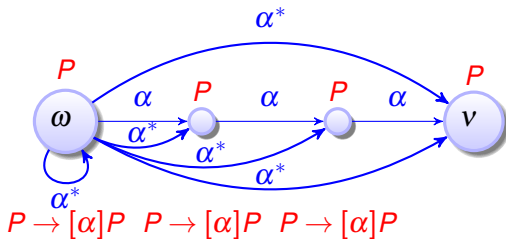
$$\vdash [\alpha^*]P \leftrightarrow P \wedge (P \rightarrow [\alpha]P)$$



$$P \rightarrow [\alpha]P \quad P \rightarrow [\alpha]P \quad P \rightarrow [\alpha]P$$

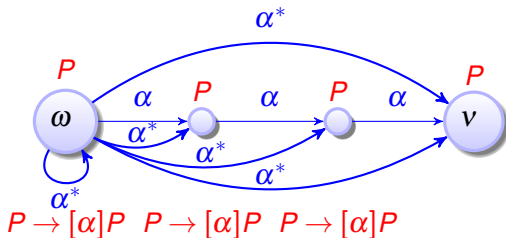
Lemma (I is sound)

$$\vdash [\alpha^*]P \leftrightarrow P \wedge [\alpha^*](P \rightarrow [\alpha]P)$$



Lemma (I is sound)

$$\vdash [\alpha^*]P \leftrightarrow P \wedge [\alpha^*](P \rightarrow [\alpha]P)$$



Problem: Inductive proof for $[\alpha^*]P$ needs proof of $[\alpha^*](P \rightarrow [\alpha]P)$

Induction Rule for Loops

Generalize induction step $[\alpha^*](P \rightarrow [\alpha]P)$

from $[\alpha^*]P \leftrightarrow P \wedge [\alpha^*](P \rightarrow [\alpha]P)$ by Gödel

$$\text{G} \frac{P}{[\alpha]P}$$

Lemma (Loop induction rule *ind* is sound)

$$\textit{ind} \frac{P \vdash [\alpha]P}{P \vdash [\alpha^*]P}$$

Induction Rule for Loops

Generalize induction step $[\alpha^*](P \rightarrow [\alpha]P)$

from $[\alpha^*]P \leftrightarrow P \wedge [\alpha^*](P \rightarrow [\alpha]P)$ by Gödel

$$\text{G} \frac{P}{[\alpha]P}$$

Lemma (Loop induction rule ind is sound)

$$\text{ind} \frac{P \vdash [\alpha]P}{P \vdash [\alpha^*]P}$$

Proof (Derived rule).

$$\frac{\frac{\text{id} \frac{*}{P \vdash P} \quad \frac{\text{G} \frac{P \vdash [\alpha]P}{\vdash P \rightarrow [\alpha]P}}{P \vdash [\alpha^*](P \rightarrow [\alpha]P)}}{\wedge R \frac{P \vdash P \wedge [\alpha^*](P \rightarrow [\alpha]P)}}{\text{I} \frac{P \vdash [\alpha^*]P}}$$

□

Induction Rule for Loops

Generalize induction step $[\alpha^*](P \rightarrow [\alpha]P)$

from $[\alpha^*]P \leftrightarrow P \wedge [\alpha^*](P \rightarrow [\alpha]P)$ by Gödel

$$G \frac{P}{[\alpha]P}$$

Lemma (Loop induction rule ind is sound)

$$ind \frac{P \vdash [\alpha]P}{P \vdash [\alpha^*]P}$$

Proof (Derived rule).

$$\frac{\frac{id \frac{*}{P \vdash P} \quad \frac{\frac{P \vdash [\alpha]P}{\vdash P \rightarrow [\alpha]P} \rightarrow R \quad G \frac{P \vdash [\alpha^*](P \rightarrow [\alpha]P)}{P \vdash [\alpha^*](P \rightarrow [\alpha]P)}}{\vdash P \wedge [\alpha^*](P \rightarrow [\alpha]P)} \wedge R}{P \vdash [\alpha^*]P} I$$

Problem: Rule ind is no equivalence. Its use of G may lose information: $[\alpha^*](P \rightarrow [\alpha]P)$ true but $P \vdash [\alpha]P$ is not valid. □

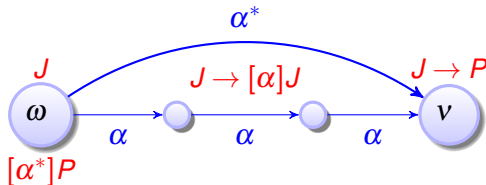
Loop Invariants

Generalize postcondition to strong loop invariant J by

$$M[\cdot] \frac{P \rightarrow Q}{[\alpha]P \rightarrow [\alpha]Q}$$

Lemma (Loop invariant rule loop is sound)

$$\text{loop} \frac{\Gamma \vdash J, \Delta \quad J \vdash [\alpha]J \quad J \vdash P}{\Gamma \vdash [\alpha^*]P, \Delta}$$



Loop Invariants

Generalize postcondition to strong loop invariant J by

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$$\text{loop} \frac{\Gamma \vdash J, \Delta \quad J \vdash [\alpha]J \quad J \vdash P}{\Gamma \vdash [\alpha^*]P, \Delta}$$

Proof (Derived rule).

$$\text{cut} \frac{\begin{array}{c} \text{ind} \frac{J \vdash [\alpha]J}{J \vdash [\alpha^*]J} \\ \rightarrow R \frac{J \vdash [\alpha^*]J, \Delta}{\Gamma \vdash J \rightarrow [\alpha^*]J, \Delta} \end{array} \quad \begin{array}{c} \frac{J \vdash P}{M[\cdot] \frac{J \vdash P}{[\alpha^*]J \vdash [\alpha^*]P}} \\ \rightarrow L \frac{\Gamma \vdash J, \Delta \quad M[\cdot] \frac{J \vdash P}{[\alpha^*]J \vdash [\alpha^*]P}}{\Gamma, J \rightarrow [\alpha^*]J \vdash [\alpha^*]P, \Delta} \end{array}}{\Gamma \vdash [\alpha^*]P, \Delta}$$

□

Loop Invariants

Generalize postcondition to strong loop invariant J by

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Lemma (Loop invariant rule loop is sound)

$$\text{loop} \frac{\Gamma \vdash J, \Delta \quad J \vdash [\alpha]J \quad J \vdash P}{\Gamma \vdash [\alpha^*]P, \Delta}$$

Proof (Derived rule).

$$\text{cut} \frac{\begin{array}{c} \text{ind} \frac{J \vdash [\alpha]J}{J \vdash [\alpha^*]J} \\ \rightarrow R \frac{J \vdash [\alpha^*]J}{\Gamma \vdash J \rightarrow [\alpha^*]J, \Delta} \end{array} \quad \begin{array}{c} \text{ind} \frac{J \vdash P}{[\alpha^*]J \vdash [\alpha^*]P} \\ \rightarrow L \frac{\Gamma \vdash J, \Delta \quad M[\cdot] \frac{J \vdash P}{[\alpha^*]J \vdash [\alpha^*]P}}{\Gamma, J \rightarrow [\alpha^*]J \vdash [\alpha^*]P, \Delta} \end{array}}{\Gamma \vdash [\alpha^*]P, \Delta}$$

Problem: Finding invariant J can be a challenge.

Misplaced $[\alpha^*]$ suggests that J needs to carry along info about α^* history.



A Simple Discrete Loop Example

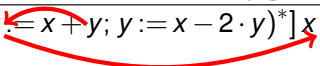
$$\text{loop} \frac{\Gamma \vdash J, \Delta \quad J \vdash [\alpha]J \quad J \vdash P}{\Gamma \vdash [\alpha^*]P, \Delta}$$

$$\begin{array}{c} \text{loop} \frac{x \geq 8 \wedge 5 \geq y \wedge y \geq 0 \vdash J \quad J \vdash [x := x + y; y := x - 2 \cdot y]J \quad J \vdash x \geq 0}{x \geq 8 \wedge 5 \geq y \wedge y \geq 0 \vdash [(x := x + y; y := x - 2 \cdot y)^*]x \geq 0} \\ \rightarrow R \frac{}{\vdash x \geq 8 \wedge 5 \geq y \wedge y \geq 0 \rightarrow [(x := x + y; y := x - 2 \cdot y)^*]x \geq 0} \end{array}$$

① $J \equiv x \geq 0$

A Simple Discrete Loop Example

$$\text{loop} \frac{\Gamma \vdash J, \Delta \quad J \vdash [\alpha]J \quad J \vdash P}{\Gamma \vdash [\alpha^*]P, \Delta}$$

$$\begin{array}{c} \text{loop} \frac{x \geq 8 \wedge 5 \geq y \wedge y \geq 0 \vdash J \quad J \vdash [x := x + y; y := x - 2 \cdot y]J \quad J \vdash x \geq 0}{x \geq 8 \wedge 5 \geq y \wedge y \geq 0 \vdash [(x := x + y; y := x - 2 \cdot y)^*]x \geq 0} \\ \rightarrow R \frac{}{\vdash x \geq 8 \wedge 5 \geq y \wedge y \geq 0 \rightarrow [(x := x + y; y := x - 2 \cdot y)^*]x \geq 0} \end{array}$$


1 $J \equiv x \geq 0$

stronger: Lacks info about y

A Simple Discrete Loop Example

$$\text{loop} \frac{\Gamma \vdash J, \Delta \quad J \vdash [\alpha]J \quad J \vdash P}{\Gamma \vdash [\alpha^*]P, \Delta}$$

$$\begin{array}{c} \text{loop} \frac{x \geq 8 \wedge 5 \geq y \wedge y \geq 0 \vdash J \quad J \vdash [x := x + y; y := x - 2 \cdot y]J \quad J \vdash x \geq 0}{x \geq 8 \wedge 5 \geq y \wedge y \geq 0 \vdash [(x := x + y; y := x - 2 \cdot y)^*]x \geq 0} \\ \rightarrow R \frac{}{\vdash x \geq 8 \wedge 5 \geq y \wedge y \geq 0 \rightarrow [(x := x + y; y := x - 2 \cdot y)^*]x \geq 0} \end{array}$$

1 $J \equiv x \geq 0$

stronger: Lacks info about y

2 $J \equiv x \geq 8 \wedge 5 \geq y \wedge y \geq 0$

A Simple Discrete Loop Example

$$\text{loop} \frac{\Gamma \vdash J, \Delta \quad J \vdash [\alpha]J \quad J \vdash P}{\Gamma \vdash [\alpha^*]P, \Delta}$$

$$\begin{array}{c} \text{loop} \frac{x \geq 8 \wedge 5 \geq y \wedge y \geq 0 \vdash J \quad J \vdash [x := x + y; y := x - 2 \cdot y]J \quad J \vdash x \geq 0}{x \geq 8 \wedge 5 \geq y \wedge y \geq 0 \vdash [(x := x + y; y := x - 2 \cdot y)^*]x \geq 0} \\ \rightarrow R \frac{}{\vdash x \geq 8 \wedge 5 \geq y \wedge y \geq 0 \rightarrow [(x := x + y; y := x - 2 \cdot y)^*]x \geq 0} \end{array}$$

① $J \equiv x \geq 0$

stronger: Lacks info about y

② $J \equiv x \geq 8 \wedge 5 \geq y \wedge y \geq 0$

weaker: Changes immediately

A Simple Discrete Loop Example

$$\text{loop} \frac{\Gamma \vdash J, \Delta \quad J \vdash [\alpha]J \quad J \vdash P}{\Gamma \vdash [\alpha^*]P, \Delta}$$

$$\begin{array}{c} \text{loop} \frac{x \geq 8 \wedge 5 \geq y \wedge y \geq 0 \vdash J \quad J \vdash [x := x + y; y := x - 2 \cdot y]J \quad J \vdash x \geq 0}{x \geq 8 \wedge 5 \geq y \wedge y \geq 0 \vdash [(x := x + y; y := x - 2 \cdot y)^*]x \geq 0} \\ \rightarrow R \frac{}{\vdash x \geq 8 \wedge 5 \geq y \wedge y \geq 0 \rightarrow [(x := x + y; y := x - 2 \cdot y)^*]x \geq 0} \end{array}$$

1 $J \equiv x \geq 0$

stronger: Lacks info about y

2 $J \equiv x \geq 8 \wedge 5 \geq y \wedge y \geq 0$

weaker: Changes immediately

3 $J \equiv x \geq 0 \wedge y \geq 0$

A Simple Discrete Loop Example

$$\text{loop} \frac{\Gamma \vdash J, \Delta \quad J \vdash [\alpha]J \quad J \vdash P}{\Gamma \vdash [\alpha^*]P, \Delta}$$

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❶ $J \equiv x \geq 0$

stronger: Lacks info about y

❷ $J \equiv x \geq 8 \wedge 5 \geq y \wedge y \geq 0$

weaker: Changes immediately

❸ $J \equiv x \geq 0 \wedge y \geq 0$

no: y may become negative if $x < y$

A Simple Discrete Loop Example

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① $J \equiv x \geq 0$

stronger: Lacks info about y

② $J \equiv x \geq 8 \wedge 5 \geq y \wedge y \geq 0$

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③ $J \equiv x \geq 0 \wedge y \geq 0$

no: y may become negative if $x < y$

④ $J \equiv x \geq y \wedge y \geq 0$

A Simple Discrete Loop Example

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❶ $J \equiv x \geq 0$

stronger: Lacks info about y

❷ $J \equiv x \geq 8 \wedge 5 \geq y \wedge y \geq 0$

weaker: Changes immediately

❸ $J \equiv x \geq 0 \wedge y \geq 0$

no: y may become negative if $x < y$

❹ $J \equiv x \geq y \wedge y \geq 0$

correct loop invariant

Forgot to Add Sequent Context Γ, Δ to Premises?

$$\frac{\Gamma \vdash J, \Delta \quad \Gamma??, J \vdash [\alpha]J, \Delta?? \quad \Gamma??, J \vdash P, \Delta??}{\Gamma \vdash [\alpha^*]P, \Delta}$$

Forgot to Add Sequent Context Γ, Δ to Premises?

$$\frac{\Gamma \vdash J, \Delta \quad \Gamma??, J \vdash [\alpha]J, \Delta?? \quad \Gamma??, J \vdash P, \Delta??}{\Gamma \vdash [\alpha^*]P, \Delta}$$

$$\begin{array}{l} \text{\color{red}⚡} \\ \hline x = 0 \vdash x \leq 1 \quad x = 0, x \leq 1 \vdash [x := x + 1]x \leq 1 \quad x \leq 1 \vdash x \leq 1 \\ \hline x = 0, x \leq 1 \vdash [(x := x + 1)^*]x \leq 1 \end{array}$$

Forgot to Add Sequent Context Γ, Δ to Premises?

$$\frac{\Gamma \vdash J, \Delta \quad \Gamma??, J \vdash [\alpha]J, \Delta?? \quad \Gamma??, J \vdash P, \Delta??}{\Gamma \vdash [\alpha^*]P, \Delta}$$

$$\begin{array}{l} \text{⚡} \\ \frac{x = 0 \vdash x \leq 1 \quad x = 0, x \leq 1 \vdash [x := x + 1]x \leq 1 \quad x \leq 1 \vdash x \leq 1}{x = 0, x \leq 1 \vdash [(x := x + 1)^*]x \leq 1} \end{array}$$

$$\begin{array}{l} \text{⚡} \\ \frac{x = 0 \vdash x \geq 0 \quad x \geq 0 \vdash [x := x + 1]x \geq 0 \quad x = 0, x \geq 0 \vdash x = 0}{x = 0 \vdash [(x := x + 1)^*]x = 0} \end{array}$$

Forgot to Add Sequent Context Γ, Δ to Premises?

$$\frac{\Gamma \vdash J, \Delta \quad \Gamma??, J \vdash [\alpha]J, \Delta?? \quad \Gamma??, J \vdash P, \Delta??}{\Gamma \vdash [\alpha^*]P, \Delta}$$

$$\begin{array}{l} \text{⚡} \\ \frac{x = 0 \vdash x \leq 1 \quad x = 0, x \leq 1 \vdash [x := x + 1]x \leq 1 \quad x \leq 1 \vdash x \leq 1}{x = 0, x \leq 1 \vdash [(x := x + 1)^*]x \leq 1} \end{array}$$

$$\begin{array}{l} \text{⚡} \\ \frac{x = 0 \vdash x \geq 0 \quad x \geq 0 \vdash [x := x + 1]x \geq 0 \quad x = 0, x \geq 0 \vdash x = 0}{x = 0 \vdash [(x := x + 1)^*]x = 0} \end{array}$$

Unsound! Be careful where assumptions go,
or CPS might go where it shouldn't.

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$$A \vdash [(\text{grav}; (?x=0; v:=-cv \cup ?x \neq 0))^*] B_{(x,v)}$$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B_{(x,v)} \equiv 0 \leq x \wedge x \leq H$$

$$\text{grav} \equiv \{x' = v, v' = -g \& x \geq 0\}$$

$$\text{loop} \frac{A \vdash j(x,v) \quad \frac{}{j(x,v) \vdash [\text{grav}; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)} \quad j(x,v) \vdash B(x,v)}{A \vdash [(\text{grav}; (?x=0; v:=-cv \cup ?x \neq 0))^*]B(x,v)}$$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$\text{grav} \equiv \{x' = v, v' = -g \& x \geq 0\}$$

Proving Quantum the Bouncing Ball

$$\text{loop} \frac{A \vdash j(x,v) \quad \frac{j(x,v) \vdash [\text{grav}; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)}{j(x,v) \vdash [\text{grav}; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)} \quad j(x,v) \vdash B(x,v)}{A \vdash [(\text{grav}; (?x=0; v:=-cv \cup ?x \neq 0))^*]B(x,v)}$$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$\text{grav} \equiv \{x' = v, v' = -g \& x \geq 0\}$$

Proving Quantum the Bouncing Ball

$$\frac{\frac{A \vdash j(x,v) \quad \frac{j(x,v) \vdash [\text{grav}][?x=0; v:=-cv \cup ?x \neq 0]j(x,v)}{j(x,v) \vdash [\text{grav}; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)}}{j(x,v) \vdash [\text{grav}; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)}}{A \vdash [(\text{grav}; (?x=0; v:=-cv \cup ?x \neq 0))^*]B(x,v)}$$

loop

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$\text{grav} \equiv \{x' = v, v' = -g \& x \geq 0\}$$

Proving Quantum the Bouncing Ball

$$\begin{array}{c}
 \text{MR} \frac{j(x,v) \vdash [\text{grav}]j(x,v)}{j(x,v) \vdash [\text{?}x=0; v:=-cv \cup \text{?}x \neq 0]j(x,v)} \\
 \text{[;]} \frac{j(x,v) \vdash [\text{grav}][\text{?}x=0; v:=-cv \cup \text{?}x \neq 0]j(x,v)}{j(x,v) \vdash [\text{grav}; (\text{?}x=0; v:=-cv \cup \text{?}x \neq 0)]j(x,v)} \\
 \text{loop} \frac{A \vdash j(x,v) \quad j(x,v) \vdash [\text{grav}; (\text{?}x=0; v:=-cv \cup \text{?}x \neq 0)]j(x,v) \quad j(x,v) \vdash B(x,v)}{A \vdash [(\text{grav}; (\text{?}x=0; v:=-cv \cup \text{?}x \neq 0))^*]B(x,v)}
 \end{array}$$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$\text{grav} \equiv \{x' = v, v' = -g \& x \geq 0\}$$

Proving Quantum the Bouncing Ball

$$\begin{array}{c}
 \text{MR} \\
 \text{[;]} \\
 \text{loop}
 \end{array}
 \frac{
 \frac{
 \frac{
 j(x,v) \vdash [\text{grav}]j(x,v) \text{ [}\cup\text{]}
 }{
 j(x,v) \vdash [?x=0; v:=-cv]j(x,v) \wedge [?x\neq 0]j(x,v)
 }{
 j(x,v) \vdash [?x=0; v:=-cv \cup ?x\neq 0]j(x,v)
 }{
 j(x,v) \vdash [\text{grav}][?x=0; v:=-cv \cup ?x\neq 0]j(x,v)
 }
 }{
 \frac{
 A \vdash j(x,v) \quad j(x,v) \vdash [\text{grav}; (?x=0; v:=-cv \cup ?x\neq 0)]j(x,v)
 }{
 j(x,v) \vdash [\text{grav}; (?x=0; v:=-cv \cup ?x\neq 0)]j(x,v)
 }
 \quad j(x,v) \vdash B(x,v)
 }{
 A \vdash [(\text{grav}; (?x=0; v:=-cv \cup ?x\neq 0))^*]B(x,v)
 }
 }$$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$\text{grav} \equiv \{x' = v, v' = -g \& x \geq 0\}$$

Proving Quantum the Bouncing Ball

$$\begin{array}{c}
 \text{MR} \\
 \text{[;]} \\
 \text{loop}
 \end{array}
 \frac{
 \frac{
 \frac{
 \frac{
 \overline{j(x,v) \vdash [?x=0; v := -cv]j(x,v)}{\wedge R}
 \quad
 \overline{j(x,v) \vdash [?x \neq 0]j(x,v)}
 }{
 j(x,v) \vdash [?x=0; v := -cv]j(x,v) \wedge [?x \neq 0]j(x,v)
 }
 }{
 j(x,v) \vdash [grav]j(x,v) \text{ [U]}
 }
 }{
 j(x,v) \vdash [grav][?x=0; v := -cv \cup ?x \neq 0]j(x,v)
 }
 }{
 \frac{
 \frac{
 \frac{
 \overline{j(x,v) \vdash [grav; (?x=0; v := -cv \cup ?x \neq 0)]j(x,v)}{\text{[;]}
 }
 \quad
 \overline{j(x,v) \vdash B(x,v)}
 }{
 j(x,v) \vdash [grav; (?x=0; v := -cv \cup ?x \neq 0)]j(x,v)
 }
 }{
 \frac{
 \overline{A \vdash j(x,v)}
 }{\text{loop}}
 }{
 A \vdash [(grav; (?x=0; v := -cv \cup ?x \neq 0))^*]B(x,v)
 }
 }
 }$$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$\text{grav} \equiv \{x' = v, v' = -g \& x \geq 0\}$$

Proving Quantum the Bouncing Ball

$$\begin{array}{c}
 \frac{j(x,v) \vdash [?x=0][v:=-cv]j(x,v)}{j(x,v) \vdash [?x=0; v:=-cv]j(x,v)} \text{ [;} \\
 \frac{j(x,v) \vdash [?x=0; v:=-cv]j(x,v) \quad j(x,v) \vdash [?x \neq 0]j(x,v)}{j(x,v) \vdash [?x=0; v:=-cv]j(x,v) \wedge [?x \neq 0]j(x,v)} \wedge R \\
 \frac{j(x,v) \vdash [grav]j(x,v) \text{ [U]} \quad j(x,v) \vdash [?x=0; v:=-cv]j(x,v) \wedge [?x \neq 0]j(x,v)}{j(x,v) \vdash [?x=0; v:=-cv \cup ?x \neq 0]j(x,v)} \\
 \text{MR} \frac{j(x,v) \vdash [grav][?x=0; v:=-cv \cup ?x \neq 0]j(x,v)}{j(x,v) \vdash [grav; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)} \\
 \text{[;} \frac{A \vdash j(x,v) \quad j(x,v) \vdash [grav; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)}{j(x,v) \vdash [grav; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)} \quad j(x,v) \vdash B(x,v) \\
 \text{loop} \frac{j(x,v) \vdash [grav; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v) \quad j(x,v) \vdash B(x,v)}{A \vdash [(grav; (?x=0; v:=-cv \cup ?x \neq 0))^*]B(x,v)}
 \end{array}$$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$\text{grav} \equiv \{x' = v, v' = -g \& x \geq 0\}$$

Proving Quantum the Bouncing Ball

$$\begin{array}{c}
 \frac{\frac{\frac{[?], \rightarrow R \quad \overline{j(x,v), x=0 \vdash [v := -cv]j(x,v)}}{j(x,v) \vdash [?x=0][v := -cv]j(x,v)}}{j(x,v) \vdash [?x=0; v := -cv]j(x,v)} \quad \frac{}{j(x,v) \vdash [?x \neq 0]j(x,v)}}{\wedge R \quad \frac{}{j(x,v) \vdash [?x=0; v := -cv]j(x,v) \wedge [?x \neq 0]j(x,v)}}}{j(x,v) \vdash [grav]j(x,v) \cup \frac{}{j(x,v) \vdash [?x=0; v := -cv \cup ?x \neq 0]j(x,v)}}} \\
 \frac{}{MR \quad \frac{}{j(x,v) \vdash [grav][?x=0; v := -cv \cup ?x \neq 0]j(x,v)}}} \\
 \frac{A \vdash j(x,v) \quad \frac{j(x,v) \vdash [grav; (?x=0; v := -cv \cup ?x \neq 0)]j(x,v)}{j(x,v) \vdash [grav; (?x=0; v := -cv \cup ?x \neq 0)]j(x,v)} \quad j(x,v) \vdash B(x,v)}{\text{loop} \quad \frac{}{A \vdash [(grav; (?x=0; v := -cv \cup ?x \neq 0))^*]B(x,v)}}}
 \end{array}$$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$\text{grav} \equiv \{x' = v, v' = -g \& x \geq 0\}$$

Proving Quantum the Bouncing Ball

$$\begin{array}{c}
 \frac{j(x,v), x=0 \vdash j(x,-cv)}{[:=]} \\
 \frac{[:=]}{[?], \rightarrow R} \frac{j(x,v), x=0 \vdash [v := -cv]j(x,v)}{j(x,v) \vdash [?x=0][v := -cv]j(x,v)} \\
 \frac{[?], \rightarrow R}{\wedge R} \frac{j(x,v) \vdash [?x=0; v := -cv]j(x,v)}{j(x,v) \vdash [?x=0; v := -cv]j(x,v) \wedge [?x \neq 0]j(x,v)} \quad \frac{}{j(x,v) \vdash [?x \neq 0]j(x,v)} \\
 \frac{j(x,v) \vdash [grav]j(x,v) \quad \wedge R}{[?], \rightarrow R} \frac{j(x,v) \vdash [?x=0; v := -cv]j(x,v) \wedge [?x \neq 0]j(x,v)}{j(x,v) \vdash [?x=0; v := -cv \cup ?x \neq 0]j(x,v)} \\
 \frac{MR}{[?]} \frac{j(x,v) \vdash [grav][?x=0; v := -cv \cup ?x \neq 0]j(x,v)}{j(x,v) \vdash [grav; (?x=0; v := -cv \cup ?x \neq 0)]j(x,v)} \\
 \frac{A \vdash j(x,v) \quad j(x,v) \vdash [grav; (?x=0; v := -cv \cup ?x \neq 0)]j(x,v)}{loop} \quad j(x,v) \vdash B(x,v) \\
 \hline
 A \vdash [(grav; (?x=0; v := -cv \cup ?x \neq 0))^*]B(x,v)
 \end{array}$$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$grav \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Bouncing Ball

$$\begin{array}{c}
 \frac{j(x,v), x=0 \vdash j(x,-cv)}{[:=]} \\
 \frac{[?], \rightarrow R \quad \frac{j(x,v), x=0 \vdash [v:=-cv]j(x,v)}{[?]} \quad \frac{j(x,v), x \neq 0 \vdash j(x,v)}{[?]}}{[?]} \\
 \frac{[?]}{\wedge R} \quad \frac{j(x,v) \vdash [?x=0; v:=-cv]j(x,v) \quad j(x,v) \vdash [?x \neq 0]j(x,v)}{j(x,v) \vdash [?x=0; v:=-cv]j(x,v) \wedge [?x \neq 0]j(x,v)} \\
 \frac{j(x,v) \vdash [\text{grav}]j(x,v) \quad [?]}{j(x,v) \vdash [?x=0; v:=-cv \cup ?x \neq 0]j(x,v)} \\
 \frac{MR \quad j(x,v) \vdash [\text{grav}][?x=0; v:=-cv \cup ?x \neq 0]j(x,v)}{[?]} \\
 \frac{A \vdash j(x,v) \quad \frac{j(x,v) \vdash [\text{grav}; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)}{j(x,v) \vdash [\text{grav}; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)} \quad j(x,v) \vdash B(x,v)}{\text{loop}} \\
 A \vdash ([\text{grav}; (?x=0; v:=-cv \cup ?x \neq 0)]^*)B(x,v)
 \end{array}$$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$\text{grav} \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Bouncing Ball

$A \vdash j(x, v)$

$j(x, v) \vdash [\text{grav}](j(x, v))$

$j(x, v), x=0 \vdash j(x, -cv)$

$j(x, v), x \neq 0 \vdash j(x, v)$

$j(x, v) \vdash B(x, v)$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x, v) \equiv 0 \leq x \wedge x \leq H$$

$$\text{grav} \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Bouncing Ball

$A \vdash j(x, v)$

$j(x, v) \vdash [\text{grav}\{x' = v, v' = -g \ \& \ x \geq 0\}](j(x, v))$

$j(x, v), x = 0 \vdash j(x, -cv)$

$j(x, v), x \neq 0 \vdash j(x, v)$

$j(x, v) \vdash B(x, v)$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x, v) \equiv 0 \leq x \wedge x \leq H$$

$$\text{grav} \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Bouncing Ball

$A \vdash j(x,v)$

$j(x,v) \vdash [\text{grav}\{x'=v, v'=-g \ \& \ x \geq 0\}](j(x,v))$

$j(x,v), x=0 \vdash j(x,-cv)$

$j(x,v), x \neq 0 \vdash j(x,v)$

$j(x,v) \vdash B(x,v)$

① $j(x,v) \equiv 0 \leq x \wedge x \leq H$

$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$

$B(x,v) \equiv 0 \leq x \wedge x \leq H$

$\text{grav} \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$

Proving Quantum the Bouncing Ball

$A \vdash j(x, v)$

$j(x, v) \vdash [\text{grav}\{x' = v, v' = -g \ \& \ x \geq 0\}](j(x, v))$

$j(x, v), x = 0 \vdash j(x, -cv)$

$j(x, v), x \neq 0 \vdash j(x, v)$

$j(x, v) \vdash B(x, v)$

① $j(x, v) \equiv 0 \leq x \wedge x \leq H$

weak: fails ODE if $v \gg 0$

$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$

$B(x, v) \equiv 0 \leq x \wedge x \leq H$

$\text{grav} \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$

Proving Quantum the Bouncing Ball

$A \vdash j(x, v)$

$j(x, v) \vdash [\text{grav}\{x' = v, v' = -g \ \& \ x \geq 0\}](j(x, v))$

$j(x, v), x = 0 \vdash j(x, -cv)$

$j(x, v), x \neq 0 \vdash j(x, v)$

$j(x, v) \vdash B(x, v)$

① $j(x, v) \equiv 0 \leq x \wedge x \leq H$

weak: fails ODE if $v \gg 0$

② $j(x, v) \equiv 2gx = 2gH - v^2 \wedge x \geq 0$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x, v) \equiv 0 \leq x \wedge x \leq H$$

$$\text{grav} \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Bouncing Ball

$A \vdash j(x, v)$

$j(x, v) \vdash [\text{grav}\{x' = v, v' = -g \ \& \ x \geq 0\}](j(x, v))$

$j(x, v), x = 0 \vdash j(x, -cv)$

$j(x, v), x \neq 0 \vdash j(x, v)$

$j(x, v) \vdash B(x, v)$

① $j(x, v) \equiv 0 \leq x \wedge x \leq H$

weak: fails ODE if $v \gg 0$

② $j(x, v) \equiv 2gx = 2gH - v^2 \wedge x \geq 0$

works: implicitly links v and x

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x, v) \equiv 0 \leq x \wedge x \leq H$$

$$\text{grav} \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

$$['] \text{-----} j(x,v) \vdash [x'=v, v'=-g \& x \geq 0]j(x,v)$$

$$\begin{array}{l} \text{[:=]} \\ \text{[:] } \\ \text{[']} \end{array} \frac{}{j(x,v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2][v := -gt](x \geq 0 \rightarrow j(x,v))}$$
$$\frac{}{j(x,v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2; v := -gt](x \geq 0 \rightarrow j(x,v))}$$
$$\frac{}{j(x,v) \vdash [x' = v, v' = -g \& x \geq 0]j(x,v)}$$

Proving Quantum the Bouncing Ball

$$\begin{array}{l} \text{[:=]} \\ \text{[:=]} \\ \text{[:]} \\ \text{[']} \end{array} \frac{}{\begin{array}{l} j(x, v) \vdash \forall t \geq 0 [x := H - \frac{g}{2} t^2] (x \geq 0 \rightarrow j(x, -gt)) \\ j(x, v) \vdash \forall t \geq 0 [x := H - \frac{g}{2} t^2] [v := -gt] (x \geq 0 \rightarrow j(x, v)) \\ j(x, v) \vdash \forall t \geq 0 [x := H - \frac{g}{2} t^2; v := -gt] (x \geq 0 \rightarrow j(x, v)) \\ j(x, v) \vdash [x' = v, v' = -g \& x \geq 0] j(x, v) \end{array}}$$

Proving Quantum the Bouncing Ball

$\forall R$	$j(x, v) \vdash \forall t \geq 0 (H - \frac{g}{2}t^2 \geq 0 \rightarrow j(H - \frac{g}{2}t^2, -gt))$
$[:=]$	$j(x, v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2](x \geq 0 \rightarrow j(x, -gt))$
$[:=]$	$j(x, v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2][v := -gt](x \geq 0 \rightarrow j(x, v))$
$[:]$	$j(x, v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2; v := -gt](x \geq 0 \rightarrow j(x, v))$
$[\prime]$	$j(x, v) \vdash [x' = v, v' = -g \ \& \ x \geq 0]j(x, v)$

Proving Quantum the Bouncing Ball

$\rightarrow R$	$j(x, v) \vdash t \geq 0 \rightarrow H - \frac{g}{2}t^2 \geq 0 \rightarrow j(H - \frac{g}{2}t^2, -gt)$
$\forall R$	$j(x, v) \vdash \forall t \geq 0 (H - \frac{g}{2}t^2 \geq 0 \rightarrow j(H - \frac{g}{2}t^2, -gt))$
$[:=]$	$j(x, v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2](x \geq 0 \rightarrow j(x, -gt))$
$[:=]$	$j(x, v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2][v := -gt](x \geq 0 \rightarrow j(x, v))$
$[:]$	$j(x, v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2; v := -gt](x \geq 0 \rightarrow j(x, v))$
$[\prime]$	$j(x, v) \vdash [x' = v, v' = -g \ \& \ x \geq 0]j(x, v)$

Proving Quantum the Bouncing Ball

$$\begin{array}{l} \text{j}(x, v), t \geq 0, H - \frac{g}{2}t^2 \geq 0 \vdash \text{j}(H - \frac{g}{2}t^2, -gt) \\ \hline \rightarrow R \quad \text{j}(x, v) \vdash t \geq 0 \rightarrow H - \frac{g}{2}t^2 \geq 0 \rightarrow \text{j}(H - \frac{g}{2}t^2, -gt) \\ \hline \forall R \quad \text{j}(x, v) \vdash \forall t \geq 0 (H - \frac{g}{2}t^2 \geq 0 \rightarrow \text{j}(H - \frac{g}{2}t^2, -gt)) \\ \hline [:=] \quad \text{j}(x, v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2] (x \geq 0 \rightarrow \text{j}(x, -gt)) \\ \hline [:=] \quad \text{j}(x, v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2] [v := -gt] (x \geq 0 \rightarrow \text{j}(x, v)) \\ \hline [;] \quad \text{j}(x, v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2; v := -gt] (x \geq 0 \rightarrow \text{j}(x, v)) \\ \hline ['] \quad \text{j}(x, v) \vdash [x' = v, v' = -g \& x \geq 0] \text{j}(x, v) \end{array}$$

$$j(x,v) \equiv 2gx=2gH-v^2 \wedge x \geq 0$$

$$2gx=2gH-v^2 \wedge x \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2 \wedge (H-\frac{g}{2}t^2) \geq 0$$

$$j(x,v), t \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash j(H-\frac{g}{2}t^2, -gt)$$

$\rightarrow R$	$j(x,v) \vdash t \geq 0 \rightarrow H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt)$
$\forall R$	$j(x,v) \vdash \forall t \geq 0 (H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt))$
$[:=]$	$j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2](x \geq 0 \rightarrow j(x, -gt))$
$[:=]$	$j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2][v:=-gt](x \geq 0 \rightarrow j(x,v))$
$[:]$	$j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2; v:=-gt](x \geq 0 \rightarrow j(x,v))$
$[']$	$j(x,v) \vdash [x'=v, v'=-g \& x \geq 0]j(x,v)$

Proving Quantum the Bouncing Ball

$$\begin{array}{l}
 \text{AR} \frac{\overline{2gx=2gH-v^2 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2} \quad \overline{H-\frac{g}{2}t^2 \geq 0 \vdash H-\frac{g}{2}t^2 \geq 0}}{2gx=2gH-v^2 \wedge x \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2 \wedge (H-\frac{g}{2}t^2) \geq 0} \\
 \frac{j(x,v), t \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash j(H-\frac{g}{2}t^2, -gt)}{\rightarrow R} \\
 \frac{\rightarrow R}{\forall R} \frac{j(x,v) \vdash t \geq 0 \rightarrow H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt)}{j(x,v) \vdash \forall t \geq 0 (H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt))} \\
 \frac{\forall R}{[:=]} \frac{j(x,v) \vdash \forall t \geq 0 [x := H-\frac{g}{2}t^2] (x \geq 0 \rightarrow j(x, -gt))}{j(x,v) \vdash \forall t \geq 0 [x := H-\frac{g}{2}t^2] [v := -gt] (x \geq 0 \rightarrow j(x, v))} \\
 \frac{[:=]}{[:]} \frac{j(x,v) \vdash \forall t \geq 0 [x := H-\frac{g}{2}t^2] [v := -gt] (x \geq 0 \rightarrow j(x, v))}{j(x,v) \vdash \forall t \geq 0 [x := H-\frac{g}{2}t^2; v := -gt] (x \geq 0 \rightarrow j(x, v))} \\
 \frac{[:]}{[']} \frac{j(x,v) \vdash \forall t \geq 0 [x := H-\frac{g}{2}t^2; v := -gt] (x \geq 0 \rightarrow j(x, v))}{j(x,v) \vdash [x' = v, v' = -g \ \& \ x \geq 0] j(x, v)}
 \end{array}$$

Proving Quantum the Bouncing Ball

$$\begin{array}{c}
 \text{*} \\
 \frac{\mathbb{R} \text{---} 2gx=2gH-v^2 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2 \quad H-\frac{g}{2}t^2 \geq 0 \vdash H-\frac{g}{2}t^2 \geq 0}{\wedge \text{R} \text{---} 2gx=2gH-v^2 \wedge x \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2 \wedge (H-\frac{g}{2}t^2) \geq 0} \\
 \frac{j(x,v), t \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash j(H-\frac{g}{2}t^2, -gt)}{\rightarrow \text{R} \text{---} j(x,v) \vdash t \geq 0 \rightarrow H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt)} \\
 \frac{\forall \text{R} \text{---} j(x,v) \vdash \forall t \geq 0 (H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt))}{[:=] \text{---} j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2](x \geq 0 \rightarrow j(x, -gt))} \\
 \frac{[:=] \text{---} j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2][v:=-gt](x \geq 0 \rightarrow j(x,v))}{[:] \text{---} j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2; v:=-gt](x \geq 0 \rightarrow j(x,v))} \\
 \frac{['] \text{---} j(x,v) \vdash [x'=v, v'=-g \& x \geq 0]j(x,v)}{['] \text{---} j(x,v) \vdash [x'=v, v'=-g \& x \geq 0]j(x,v)}
 \end{array}$$

Proving Quantum the Bouncing Ball

$$\begin{array}{c}
 \begin{array}{c}
 \text{*} \\
 \mathbb{R} \text{---} \\
 2gx=2gH-v^2 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2
 \end{array}
 \quad
 \begin{array}{c}
 \text{*} \\
 \text{id} \text{---} \\
 H-\frac{g}{2}t^2 \geq 0 \vdash H-\frac{g}{2}t^2 \geq 0
 \end{array} \\
 \hline
 \wedge \mathbb{R} \text{---} \\
 2gx=2gH-v^2 \wedge x \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2 \wedge (H-\frac{g}{2}t^2) \geq 0 \\
 \hline
 \begin{array}{c}
 \text{j}(x,v), t \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash \text{j}(H-\frac{g}{2}t^2, -gt) \\
 \rightarrow \mathbb{R} \text{---} \\
 \text{j}(x,v) \vdash t \geq 0 \rightarrow H-\frac{g}{2}t^2 \geq 0 \rightarrow \text{j}(H-\frac{g}{2}t^2, -gt) \\
 \forall \mathbb{R} \text{---} \\
 \text{j}(x,v) \vdash \forall t \geq 0 (H-\frac{g}{2}t^2 \geq 0 \rightarrow \text{j}(H-\frac{g}{2}t^2, -gt)) \\
 [:=] \text{---} \\
 \text{j}(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2](x \geq 0 \rightarrow \text{j}(x, -gt)) \\
 [:=] \text{---} \\
 \text{j}(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2][v:=-gt](x \geq 0 \rightarrow \text{j}(x,v)) \\
 [:] \text{---} \\
 \text{j}(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2; v:=-gt](x \geq 0 \rightarrow \text{j}(x,v)) \\
 ['] \text{---} \\
 \text{j}(x,v) \vdash [x'=v, v'=-g \& x \geq 0] \text{j}(x,v)
 \end{array}
 \end{array}$$

Proving Quantum the Bouncing Ball

$$\begin{array}{c}
 \text{IR} \frac{\text{---} * \text{---}}{2gx=2gH-v^2 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2} \text{id} \frac{\text{---} * \text{---}}{H-\frac{g}{2}t^2 \geq 0 \vdash H-\frac{g}{2}t^2 \geq 0} \\
 \wedge R \frac{\text{---}}{2gx=2gH-v^2 \wedge x \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2 \wedge (H-\frac{g}{2}t^2) \geq 0} \\
 \text{---} \\
 \text{---} \\
 \rightarrow R \frac{\text{---}}{j(x,v), t \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash j(H-\frac{g}{2}t^2, -gt)} \\
 \text{---} \\
 \forall R \frac{\text{---}}{j(x,v) \vdash t \geq 0 \rightarrow H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt)} \\
 \text{---} \\
 \forall R \frac{\text{---}}{j(x,v) \vdash \forall t \geq 0 (H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt))} \\
 \text{---} \\
 [:=] \frac{\text{---}}{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2](x \geq 0 \rightarrow j(x, -gt))} \\
 \text{---} \\
 [:=] \frac{\text{---}}{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2][v:=-gt](x \geq 0 \rightarrow j(x,v))} \\
 \text{---} \\
 [i] \frac{\text{---}}{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2; v:=-gt](x \geq 0 \rightarrow j(x,v))} \\
 \text{---} \\
 ['] \frac{\text{---}}{j(x,v) \vdash [x'=v, v'=-g \& x \geq 0]j(x,v)}
 \end{array}$$

- Wait, was this actually a safety proof for Quantum?

Proving Quantum the Bouncing Ball

$$\begin{array}{c}
 \text{IR} \frac{\text{---} * \text{---}}{2gx=2gH-v^2 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2} \text{id} \frac{\text{---} * \text{---}}{H-\frac{g}{2}t^2 \geq 0 \vdash H-\frac{g}{2}t^2 \geq 0} \\
 \wedge R \frac{\text{---}}{2gx=2gH-v^2 \wedge x \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2 \wedge (H-\frac{g}{2}t^2) \geq 0} \\
 \frac{j(x,v), t \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash j(H-\frac{g}{2}t^2, -gt)}{\rightarrow R} \\
 \frac{j(x,v) \vdash t \geq 0 \rightarrow H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt)}{\forall R} \\
 \frac{j(x,v) \vdash \forall t \geq 0 (H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt))}{[:=]} \\
 \frac{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2](x \geq 0 \rightarrow j(x, -gt))}{[:=]} \\
 \frac{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2][v:=-gt](x \geq 0 \rightarrow j(x,v))}{[:]} \\
 \frac{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2; v:=-gt](x \geq 0 \rightarrow j(x,v))}{[:]} \\
 \frac{[:]}{j(x,v) \vdash [x'=v, v'=-g \& x \geq 0]j(x,v)}
 \end{array}$$

- Wait, was this actually a safety proof for Quantum?
- Oh no! The solutions we sneaked into $[']$ only solve the ODE/IVP if $x = H, v = 0$ which assumption $j(x,v)$ can't guarantee!

Proving Quantum the Bouncing Ball

$$\begin{array}{c}
 \text{IR} \frac{\text{---} * \text{---}}{2gx=2gH-v^2 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2} \text{id} \frac{\text{---} * \text{---}}{H-\frac{g}{2}t^2 \geq 0 \vdash H-\frac{g}{2}t^2 \geq 0} \\
 \wedge \text{R} \frac{\text{---}}{2gx=2gH-v^2 \wedge x \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2 \wedge (H-\frac{g}{2}t^2) \geq 0} \\
 \frac{j(x,v), t \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash j(H-\frac{g}{2}t^2, -gt)}{\rightarrow \text{R} \frac{\text{---}}{j(x,v) \vdash t \geq 0 \rightarrow H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt)}} \\
 \frac{\text{---}}{\forall \text{R} \frac{\text{---}}{j(x,v) \vdash \forall t \geq 0 (H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt))}} \\
 \frac{\text{---}}{[:=] \frac{\text{---}}{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2](x \geq 0 \rightarrow j(x, -gt))}} \\
 \frac{\text{---}}{[:=] \frac{\text{---}}{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2][v:=-gt](x \geq 0 \rightarrow j(x,v))}} \\
 \frac{\text{---}}{[:] \frac{\text{---}}{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2; v:=-gt](x \geq 0 \rightarrow j(x,v))}} \\
 \frac{\text{---}}{['] \frac{\text{---}}{j(x,v) \vdash [x'=v, v'=-g \& x \geq 0]j(x,v)}}
 \end{array}$$

- Wait, was this actually a safety proof for Quantum?
- Oh no! The solutions we sneaked into ['] only solve the ODE/IVP if $x = H, v = 0$ which assumption $j(x,v)$ can't guarantee!
- **Never use solutions without proof!** ▶ Todo redo proof with true solution

Clumsy Quantum Misplaced the Constants

loop $A \vdash [\alpha^*]B(x,v)$

1 $j(x,v) \equiv 2gx = 2gH - v^2 \wedge x \geq 0$

2 $p \equiv c=1 \wedge g > 0$

loop $A \vdash [\alpha^*]B(x,v)$

- 1 $j(x,v) \equiv 2gx = 2gH - v^2 \wedge x \geq 0$
- 2 $p \equiv c=1 \wedge g > 0$
- 3 $J \equiv j(x,v) \wedge p$ as loop invariant

Clumsy Quantum Misplaced the Constants

$$\text{loop} \frac{\frac{\mathbb{R} \overline{A \vdash j(x,v) \wedge p} \quad \square \wedge \overline{j(x,v) \wedge p \vdash [\alpha](j(x,v) \wedge p)} \quad \mathbb{R} \overline{j(x,v) \wedge p \vdash B(x,v)}}{A \vdash [\alpha^*] B(x,v)}}{*}$$

- 1 $j(x,v) \equiv 2gx = 2gH - v^2 \wedge x \geq 0$
- 2 $p \equiv c = 1 \wedge g > 0$
- 3 $J \equiv j(x,v) \wedge p$ as loop invariant

Clumsy Quantum Misplaced the Constants

$$\Box \wedge [\alpha](P \wedge Q) \leftrightarrow [\alpha]P \wedge [\alpha]Q$$

$$\frac{\text{loop} \quad \frac{\mathbb{R} \quad A \vdash j(x,v) \wedge p \quad * \quad \frac{\text{above} \quad \frac{j(x,v) \wedge p \vdash [\alpha]j(x,v) \quad \vee \quad j(x,v) \wedge p \vdash [\alpha]p}{j(x,v) \wedge p \vdash [\alpha]j(x,v) \wedge [\alpha]p}}{\Box \wedge \quad j(x,v) \wedge p \vdash [\alpha](j(x,v) \wedge p)} \quad \mathbb{R} \quad j(x,v) \wedge p \vdash B(x,v)}}{A \vdash [\alpha^*]B(x,v)}}$$

- 1 $j(x,v) \equiv 2gx = 2gH - v^2 \wedge x \geq 0$
- 2 $p \equiv c=1 \wedge g > 0$
- 3 $J \equiv j(x,v) \wedge p$ as loop invariant

Clumsy Quantum Misplaced the Constants

$$\Box \wedge [\alpha](P \wedge Q) \leftrightarrow [\alpha]P \wedge [\alpha]Q$$

$$\forall p \rightarrow [\alpha]p \quad (FV(p) \cap BV(\alpha) = \emptyset)$$

$$\frac{\frac{\frac{\text{loop}}{\mathbb{R}} \overline{A \vdash j(x,v) \wedge p} \quad *}{\mathbb{R}} \overline{A \vdash j(x,v) \wedge p} \quad \Box \wedge \frac{\frac{\text{above} \quad \frac{\overline{j(x,v) \wedge p \vdash [\alpha]j(x,v)} \quad \vee \quad \frac{\overline{j(x,v) \wedge p \vdash [\alpha]p}}{*}}{\overline{j(x,v) \wedge p \vdash [\alpha]j(x,v) \wedge [\alpha]p}}}{\overline{j(x,v) \wedge p \vdash [\alpha](j(x,v) \wedge p)}} \quad \mathbb{R} \overline{j(x,v) \wedge p \vdash B(x,v)}}{\overline{A \vdash [\alpha^*]B(x,v)}}$$

- 1 $j(x,v) \equiv 2gx = 2gH - v^2 \wedge x \geq 0$
- 2 $p \equiv c = 1 \wedge g > 0$
- 3 $J \equiv j(x,v) \wedge p$ as loop invariant

Clumsy Quantum Misplaced the Constants

$$\Box \wedge [\alpha](P \wedge Q) \leftrightarrow [\alpha]P \wedge [\alpha]Q$$

$$\forall p \rightarrow [\alpha]p \quad (FV(p) \cap BV(\alpha) = \emptyset)$$

$$\frac{\frac{\frac{\text{loop}}{\mathbb{R}} \overline{A \vdash j(x,v) \wedge p} \quad *}{\wedge R} \quad \frac{\frac{\text{above}}{j(x,v) \wedge p \vdash [\alpha]j(x,v)} \quad \frac{\text{*}}{j(x,v) \wedge p \vdash [\alpha]p}}{j(x,v) \wedge p \vdash [\alpha]j(x,v) \wedge [\alpha]p}}{\wedge \wedge} \quad \frac{\mathbb{R}}{j(x,v) \wedge p \vdash B(x,v)} \quad *}{A \vdash [\alpha^*]B(x,v)}$$

- 1 $j(x,v) \equiv 2gx = 2gH - v^2 \wedge x \geq 0$
- 2 $p \equiv c = 1 \wedge g > 0$
- 3 $J \equiv j(x,v) \wedge p$ as loop invariant

Note: constants $c = 1 \wedge g > 0$ that never change are usually elided from J

Proposition (Quantum can bounce around safely)

$0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 = c \rightarrow$

$[(\{x' = v, v' = -g \& x \geq 0\}; (?x = 0; v := -cv \cup ?x \neq 0))^*](0 \leq x \wedge x \leq H)$

requires $(0 \leq x \wedge x = H \wedge v = 0)$

requires $(g > 0 \wedge 1 = c)$

ensures $(0 \leq x \wedge x \leq H)$

$\{\{x' = v, v' = -g \& x \geq 0\};$

$(?x = 0; v := -cv \cup ?x \neq 0)\}^* @invariant(2gx = 2gH - v^2 \wedge x \geq 0)$

Invariant Contracts

Invariants play a crucial role in CPS design. Capture them if you can.
Use **@invariant()** contracts in your hybrid programs.

- 1 Learning Objectives
- 2 Induction for Loops
 - Iteration Axiom
 - Induction Axiom
 - Induction Rule for Loops
 - Loop Invariants
 - Simple Example
 - Contextual Soundness Requirements
- 3 Operationalize Invariant Construction
 - Bouncing Ball
 - Rescuing Misplaced Constants
 - Safe Quantum
- 4 Summary

The lion's share of understanding comes from understanding what does change (variants/progress measures) and what doesn't change (invariants).

Invariants are a fundamental force of CS

Variants are another fundamental force of CS

Summary: Loops, Generalizations, Splittings

$$I \quad [\alpha^*]P \leftrightarrow P \wedge [\alpha^*](P \rightarrow [\alpha]P)$$

$$G \quad \frac{P}{[\alpha]P}$$

$$M[\cdot] \quad \frac{P \rightarrow Q}{[\alpha]P \rightarrow [\alpha]Q}$$

$$\text{loop} \quad \frac{\Gamma \vdash J, \Delta \quad J \vdash [\alpha]J \quad J \vdash P}{\Gamma \vdash [\alpha^*]P, \Delta}$$

$$\text{MR} \quad \frac{\Gamma \vdash [\alpha]Q, \Delta \quad Q \vdash P}{\Gamma \vdash [\alpha]P, \Delta}$$

$$[\wedge] \quad [\alpha](P \wedge Q) \leftrightarrow [\alpha]P \wedge [\alpha]Q$$

$$\forall p \rightarrow [\alpha]p \quad (FV(p) \cap BV(\alpha) = \emptyset)$$

5 Appendix

- Iteration Axiom
- Iterations & Splitting the Box
- Iteration & Generalizations

compositional semantics \Rightarrow compositional rules!

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$A \vdash [\alpha^*]B$$

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$[*] \frac{A \vdash B \wedge [\alpha][\alpha^*]B}{A \vdash [\alpha^*]B}$$

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$\frac{\frac{[*] \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash B \wedge [\alpha][\alpha^*]B}}{[*] \frac{A \vdash B \wedge [\alpha][\alpha^*]B}}{A \vdash [\alpha^*]B}}$$

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$\begin{array}{c}
 \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 \frac{[*]}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \frac{[*]}{A \vdash B \wedge [\alpha][\alpha^*]B} \\
 \frac{[*]}{A \vdash [\alpha^*]B}
 \end{array}$$

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$[] \wedge [\alpha](P \wedge Q) \leftrightarrow [\alpha]P \wedge [\alpha]Q$$

$$\begin{array}{l}
 \frac{}{A \vdash B \wedge [\alpha]B \wedge [\alpha][\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \frac{}{A \vdash B \wedge [\alpha][\alpha^*]B} \\
 \frac{}{A \vdash [\alpha^*]B}
 \end{array}$$

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$[\wedge] [\alpha](P \wedge Q) \leftrightarrow [\alpha]P \wedge [\alpha]Q$$

$$\begin{array}{c}
 \frac{}{A \vdash B \wedge [\alpha]B \wedge [\alpha]([\alpha]B \wedge [\alpha][\alpha^*]B)} \\
 \frac{[\wedge]}{A \vdash B \wedge [\alpha]B \wedge [\alpha][\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \frac{[\wedge]}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 \frac{[*]}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \frac{[*]}{A \vdash B \wedge [\alpha][\alpha^*]B} \\
 \frac{[*]}{A \vdash [\alpha^*]B}
 \end{array}$$

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$[] \wedge [\alpha](P \wedge Q) \leftrightarrow [\alpha]P \wedge [\alpha]Q$$

$$\begin{array}{c}
 \frac{}{A \vdash B \wedge [\alpha]B \wedge [\alpha][\alpha]B \wedge [\alpha][\alpha][\alpha][\alpha^*]B} \\
 \frac{}{A \vdash B \wedge [\alpha]B \wedge [\alpha]([\alpha]B \wedge [\alpha][\alpha][\alpha^*]B)} \\
 \frac{}{A \vdash B \wedge [\alpha]B \wedge [\alpha][\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \frac{}{A \vdash B \wedge [\alpha][\alpha^*]B} \\
 \frac{}{A \vdash [\alpha^*]B}
 \end{array}$$

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$[\Box] \wedge [\alpha](P \wedge Q) \leftrightarrow [\alpha]P \wedge [\alpha]Q$$

$$\begin{array}{c}
 \frac{A \vdash B \quad A \vdash [\alpha]B \quad A \vdash [\alpha][\alpha]B \quad A \vdash [\alpha][\alpha][\alpha][\alpha^*]B}{\wedge R} \\
 \frac{A \vdash B \wedge [\alpha]B \wedge [\alpha][\alpha]B \wedge [\alpha][\alpha][\alpha][\alpha^*]B}{[\Box] \wedge} \\
 \frac{A \vdash B \wedge [\alpha]B \wedge [\alpha]([\alpha]B \wedge [\alpha][\alpha][\alpha^*]B)}{[\Box] \wedge} \\
 \frac{A \vdash B \wedge [\alpha]B \wedge [\alpha][\alpha](B \wedge [\alpha][\alpha^*]B)}{[\Box] \wedge} \\
 \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}{[*]} \\
 \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{[*]} \\
 \frac{A \vdash B \wedge [\alpha][\alpha^*]B}{[*]} \\
 \frac{A \vdash [\alpha^*]B}{[*]}
 \end{array}$$

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$[] \wedge [\alpha](P \wedge Q) \leftrightarrow [\alpha]P \wedge [\alpha]Q$$

$$\begin{array}{c}
 A \vdash B \quad A \vdash [\alpha]B \quad A \vdash [\alpha][\alpha]B \quad A \vdash [\alpha][\alpha][\alpha][\alpha^*]B \\
 \wedge R \frac{}{A \vdash B \wedge [\alpha]B \wedge [\alpha][\alpha]B \wedge [\alpha][\alpha][\alpha][\alpha^*]B} \\
 [] \wedge \frac{}{A \vdash B \wedge [\alpha]B \wedge [\alpha]([\alpha]B \wedge [\alpha][\alpha][\alpha^*]B)} \\
 [] \wedge \frac{}{A \vdash B \wedge [\alpha]B \wedge [\alpha][\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 [] \wedge \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 [*] \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 [*] \frac{}{A \vdash B \wedge [\alpha][\alpha^*]B} \\
 [*] \frac{}{A \vdash [\alpha^*]B}
 \end{array}$$

- 1 Simple approach ... if we don't mind unrolling until the end of time
- 2 Useful for finding counterexamples

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$\begin{array}{c}
 \hline
 A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)) \\
 \hline
 [*] \hline
 A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B) \\
 \hline
 [*] \hline
 A \vdash B \wedge [\alpha][\alpha^*]B \\
 \hline
 [*] \hline
 A \vdash [\alpha^*]B
 \end{array}$$

Loops of Proofs: Iterations & Generalizations

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$\text{MR} \frac{\Gamma \vdash [\alpha]Q, \Delta \quad Q \vdash P}{\Gamma \vdash [\alpha]P, \Delta}$$

$$\begin{array}{c} A \vdash B \\ \hline \wedge\text{R} \frac{A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\ \hline [*] \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash B \wedge [\alpha][\alpha^*]B} \\ \hline [*] \frac{A \vdash B \wedge [\alpha][\alpha^*]B}{A \vdash [\alpha^*]B} \end{array}$$

Loops of Proofs: Iterations & Generalizations

$$[*] \quad [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$\text{MR} \quad \frac{\Gamma \vdash [\alpha]Q, \Delta \quad Q \vdash P}{\Gamma \vdash [\alpha]P, \Delta}$$

$$\begin{array}{c}
 A \vdash [\alpha]J_1 \quad \frac{}{J_1 \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 A \vdash B_{\text{MR}} \quad \frac{}{A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 \wedge\text{R} \quad \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 [*] \quad \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 [*] \quad \frac{}{A \vdash B \wedge [\alpha][\alpha^*]B} \\
 [*] \quad \frac{}{A \vdash [\alpha^*]B}
 \end{array}$$

Loops of Proofs: Iterations & Generalizations

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$\text{MR} \frac{\Gamma \vdash [\alpha]Q, \Delta \quad Q \vdash P}{\Gamma \vdash [\alpha]P, \Delta}$$

$$\begin{array}{c}
 \begin{array}{c}
 J_2 \vdash B \quad \frac{J_2 \vdash [\alpha]J_3 \quad \dots}{J_2 \vdash [\alpha][\alpha^*]B} \\
 J_1 \vdash [\alpha]J_2 \quad \wedge R \frac{\quad}{J_2 \vdash B \wedge [\alpha][\alpha^*]B} \\
 J_1 \vdash B \quad \text{MR} \frac{\quad}{J_1 \vdash [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 A \vdash [\alpha]J_1 \quad \wedge R \frac{\quad}{J_1 \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 A \vdash B \quad \text{MR} \frac{\quad}{A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}
 \end{array} \\
 \wedge R \frac{\quad}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 [*] \frac{\quad}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 [*] \frac{\quad}{A \vdash B \wedge [\alpha][\alpha^*]B} \\
 [*] \frac{\quad}{A \vdash [\alpha^*]B}
 \end{array}$$

Loops of Proofs: Common Generalizations

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$\text{MR} \frac{\Gamma \vdash [\alpha]Q, \Delta \quad Q \vdash P}{\Gamma \vdash [\alpha]P, \Delta}$$

$$\begin{array}{c}
 J \vdash B \quad \frac{J \vdash [\alpha]J \quad \dots}{J \vdash [\alpha][\alpha^*]B} \\
 J \vdash [\alpha]J \quad \wedge R \frac{J \vdash B \quad J \vdash [\alpha][\alpha^*]B}{J \vdash B \wedge [\alpha][\alpha^*]B} \\
 J \vdash B \text{MR} \frac{J \vdash [\alpha]J \quad J \vdash B \wedge [\alpha][\alpha^*]B}{J \vdash [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 A \vdash [\alpha]J \quad \wedge R \frac{A \vdash [\alpha]J \quad J \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 A \vdash B \text{MR} \frac{A \vdash B \text{MR} \quad J \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 \wedge R \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 [*] \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 [*] \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash B \wedge [\alpha][\alpha^*]B} \\
 [*] \frac{A \vdash B \wedge [\alpha][\alpha^*]B}{A \vdash [\alpha^*]B}
 \end{array}$$

Loops of Proofs: Extracting a Proof Rule

$$\begin{array}{c}
 \frac{J \vdash [\alpha]J \quad J \vdash B}{A \vdash [\alpha^*]B} \qquad \qquad \qquad [*] \quad [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P \\
 \\
 \text{MR} \quad \frac{\Gamma \vdash [\alpha]Q, \Delta \quad Q \vdash P}{\Gamma \vdash [\alpha]P, \Delta} \\
 \\
 \begin{array}{c}
 J \vdash B \quad \frac{J \vdash [\alpha]J \quad \dots}{J \vdash [\alpha][\alpha^*]B} \\
 \wedge R \quad \frac{J \vdash B \quad J \vdash [\alpha][\alpha^*]B}{J \vdash B \wedge [\alpha][\alpha^*]B} \\
 \\
 J \vdash B_{\text{MR}} \quad \frac{J \vdash B \quad J \vdash [\alpha][\alpha^*]B}{J \vdash B \wedge [\alpha][\alpha^*]B} \\
 \\
 A \vdash [\alpha]J \quad \wedge R \quad \frac{A \vdash [\alpha]J \quad J \vdash B \wedge [\alpha][\alpha^*]B}{A \vdash [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \\
 A \vdash B_{\text{MR}} \quad \frac{A \vdash B \quad A \vdash [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \\
 \wedge R \quad \frac{A \vdash B \quad A \vdash [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \\
 [*] \quad \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \\
 [*] \quad \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash B \wedge [\alpha][\alpha^*]B} \\
 \\
 [*] \quad \frac{A \vdash B \wedge [\alpha][\alpha^*]B}{A \vdash [\alpha^*]B}
 \end{array}
 \end{array}$$

Loops of Proofs: Extracting a Proof Rule

$$\frac{A \vdash J \quad J \vdash [\alpha]J \quad J \vdash B}{A \vdash [\alpha^*]B}$$

$$[*] \quad [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$\text{MR} \quad \frac{\Gamma \vdash [\alpha]Q, \Delta \quad Q \vdash P}{\Gamma \vdash [\alpha]P, \Delta}$$

$$J \vdash B \quad \frac{J \vdash [\alpha]J \quad \dots}{J \vdash [\alpha][\alpha^*]B}$$

$$J \vdash [\alpha]J \quad \wedge\text{R} \quad \frac{J \vdash B \quad \text{MR} \quad \frac{J \vdash [\alpha]J \quad \dots}{J \vdash [\alpha][\alpha^*]B}}{J \vdash B \wedge [\alpha][\alpha^*]B}$$

$$J \vdash B \quad \text{MR} \quad \frac{J \vdash B \wedge [\alpha][\alpha^*]B}{J \vdash [\alpha](B \wedge [\alpha][\alpha^*]B)}$$

$$A \vdash [\alpha]J \quad \wedge\text{R} \quad \frac{A \vdash [\alpha]J \quad \text{MR} \quad \frac{J \vdash B \wedge [\alpha][\alpha^*]B}{J \vdash [\alpha](B \wedge [\alpha][\alpha^*]B)}}{A \vdash [\alpha]J \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}$$

$$A \vdash B \quad \text{MR} \quad \frac{A \vdash B \quad \wedge\text{R} \quad \frac{A \vdash [\alpha]J \quad \text{MR} \quad \frac{J \vdash B \wedge [\alpha][\alpha^*]B}{J \vdash [\alpha](B \wedge [\alpha][\alpha^*]B)}}{A \vdash [\alpha]J \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}}{A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}$$

$$\wedge\text{R} \quad \frac{A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}$$

$$[*] \quad \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}$$

$$[*] \quad \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash B \wedge [\alpha][\alpha^*]B}$$

$$[*] \quad \frac{A \vdash B \wedge [\alpha][\alpha^*]B}{A \vdash [\alpha^*]B}$$

Loops of Proofs: Loop Invariants

$$\text{loop} \frac{A \vdash J \quad J \vdash [\alpha]J \quad J \vdash B}{A \vdash [\alpha^*]B}$$

Invariant J generalized
intermediate condition

$$[*] \quad [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$\text{MR} \quad \frac{\Gamma \vdash [\alpha]Q, \Delta \quad Q \vdash P}{\Gamma \vdash [\alpha]P, \Delta}$$

$$\begin{array}{c}
 \begin{array}{c}
 J \vdash B \quad J \vdash [\alpha]J \quad \dots \\
 \hline
 J \vdash [\alpha][\alpha^*]B \\
 \hline
 J \vdash [\alpha]J \quad \wedge\text{R} \quad \hline
 J \vdash B \wedge [\alpha][\alpha^*]B \\
 \hline
 J \vdash B \quad \text{MR} \quad \hline
 J \vdash [\alpha](B \wedge [\alpha][\alpha^*]B) \\
 \hline
 A \vdash [\alpha]J \quad \wedge\text{R} \quad \hline
 J \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B) \\
 \hline
 A \vdash B \quad \text{MR} \quad \hline
 A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)) \\
 \hline
 \wedge\text{R} \quad \hline
 A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)) \\
 \hline
 [*] \quad \hline
 A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B) \\
 \hline
 [*] \quad \hline
 A \vdash B \wedge [\alpha][\alpha^*]B \\
 \hline
 [*] \quad \hline
 A \vdash [\alpha^*]B
 \end{array}
 \end{array}$$



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