

Sorting 7A

Quadratic Sorts



Selection Sort

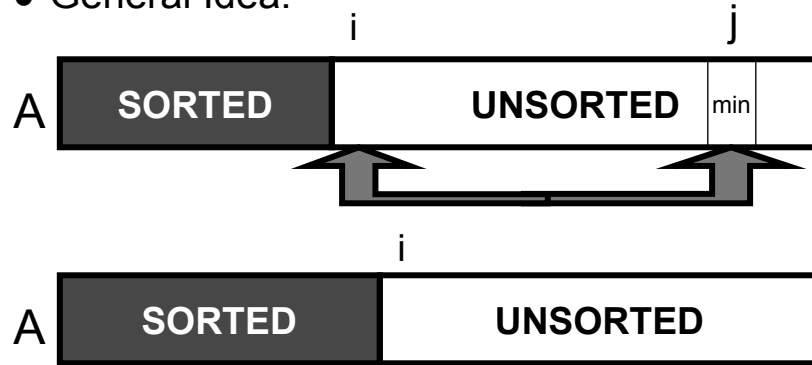


- Let A be an array of n elements, and we wish to sort these elements in non-decreasing order.
- Basic Algorithm:
 - Set $i = 0$.
 - While $i < n$ do the following:
 - Find j , $i \leq j \leq n-1$, such that $A[j] \leq A[k]$, $\forall k, i \leq k \leq n-1$.
 - Swap $A[j]$ with $A[i]$
 - Add 1 to i
- This algorithm works in place, meaning it uses its own storage to perform the sort.



Selection Sort

- General Idea:



Loop invariant: $A[0..i-1]$ are sorted in non-decreasing order.



Selection Sort Example

66	44	99	55	11	88	22	77	33
<u>11</u>	44	99	55	66	88	22	77	33
11	<u>22</u>	99	55	66	88	44	77	33
11	22	<u>33</u>	55	66	88	44	77	99
11	22	33	<u>44</u>	66	88	55	77	99
11	22	33	44	<u>55</u>	88	66	77	99
11	22	33	44	55	<u>66</u>	88	77	99
11	22	33	44	55	66	<u>77</u>	88	99
11	22	33	44	55	66	77	<u>88</u>	99

Run time analysis



- Worst Case:
 - Search for 1st min: n-1 comparisons
 - Search for 2nd min: n-2 comparisons
 - ...
 - Search for 2nd-to-last min: 1 comparison
 - Total comparisons:
(n-1) + (n-2) + ... + 2 + 1 = O(_____)
- Average Case: = O(_____)
- Best Case: = O(_____)

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Insertion Sort



- Let A be an array of n elements, and we wish to sort these elements in non-decreasing order.
- Basic Algorithm:
 - Set $i = 1$
 - While $i < n$ do the following:
 - Set $item = A[i]$.
 - Shift $A[j]$ to $A[j+1]$, $\forall j, j < i$, where $A[j] > item$.
 - Let k be the smallest j above, or $k=i$ if no shifts.
 - Set $A[k] = item$.
 - Add 1 to i .
- This algorithm also works in place.

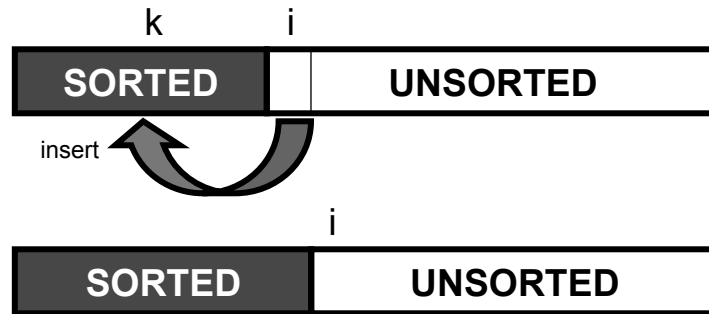
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Insertion Sort

- General Idea:



Loop invariant: $A[0..i-1]$ are sorted in non-decreasing order.



Insertion Sort Example

<u>66</u>	44	99	55	11	88	22	77	33
44	<u>66</u>	99	55	11	88	22	77	33
44	66	<u>99</u>	55	11	88	22	77	33
44	55	66	<u>99</u>	11	88	22	77	33
11	44	55	66	<u>99</u>	88	22	77	33
11	44	55	66	88	<u>99</u>	22	77	33
11	22	44	55	66	88	<u>99</u>	77	33
11	22	44	55	66	77	88	<u>99</u>	33
11	22	33	44	55	66	77	88	<u>99</u>



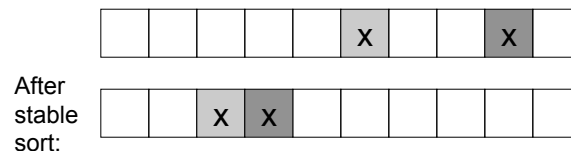
Run time analysis

- Worst Case (when does this occur?):
 - Insert 2nd element: 1 comparison
 - Insert 3rd element: 2 comparisons
 - ...
 - Insert last element: n-1 comparisons
 - Total comparisons:
 - $1 + 2 + \dots + (n-1) = O(\underline{\hspace{2cm}})$
- Average Case: $= O(\underline{\hspace{2cm}})$
- Best Case: $= O(\underline{\hspace{2cm}})$



Stable Sorts

- A sort is stable if two elements with the same value maintain their same relative order before and after the sort is performed.



- Is selection sort stable?
- Is insertion sort stable?

Quadratic Sorts

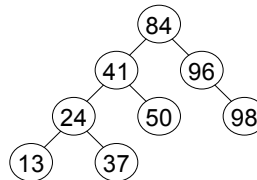


- Quadratic sorts have a worst-case order of complexity of $O(n^2)$
- Selection sort always performs poorly, even on a sequence of sorted elements!
- Insertion sort performs much better if the elements are sorted or nearly sorted.
- Another famous quadratic sort: “Bubble sort”

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Tree Sort



- Build a binary search tree out of the elements.
- Traverse the tree using an inorder traversal to get the elements in increasing order.
- Worst case order of complexity:
 - $O(n^2)$ to build the binary search tree (Why?)
 - $O(n)$ to traverse the binary tree. (Why?)
 - Total: $O(n^2) + O(n) = \underline{\hspace{2cm}}$

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